

High-Speed Implementation of the SHA-3 Candidate Shabal

2nd SHA-3 Conference

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Introduction

2 ways of studying SHA-3 candidates hardware cost on FPGA

- “Fair and comprehensive comparison” of all the candidates.
 - [Tillich *et al.*, ePrint 2009], [Kobayashi *et al.*, ePrint 2010], [Henzen *et al.*, CHES 2010], [Gaj *et al.*, CHES 2010], [Baldwin *et al.*, 2nd SHA-3 Conf.]
- Investigate **hardware optimizations** of some ones.

Why have we focused on Shabal and its throughput?

- It can be implemented in resource-constrained devices [Detray *et al.*, ePrint 2010].
- At first sight, its restricted parallelizability seems to **limit its throughput** (TP).
 - ⇒ Fortunately, **a neat method allows to improve its TP!**
 - ⇒ **Unfolding Method**

Outline

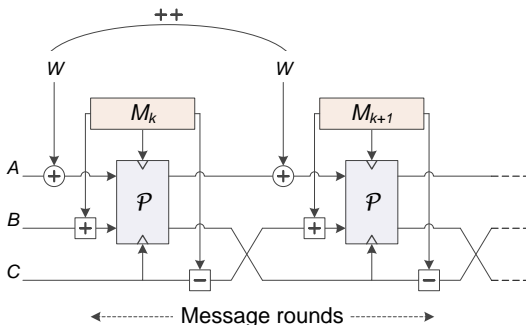
- ① Main Characteristics of Shabal
- ② Hardware Investigations
- ③ Implementation Results
- ④ Conclusion and Future Works

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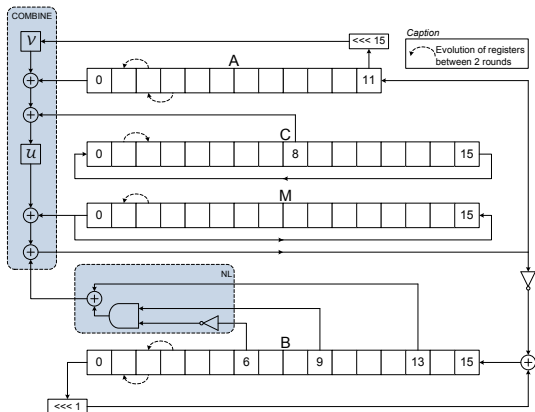
Parameters and Generic Construction

- Iteration of permutation \mathcal{P} on padded message blocks M_k (16×32 -bit words).
- Internal state (A, B, C) initialized by an Initial Vector.
- W : 2×32 -bit counter.
- \oplus : XOR, \boxplus (\boxminus): Wordwise addition (subtraction) modulo 2^{32} .



Permutation \mathcal{P}

- 48 iterations of a Non-Linear Feedback Shift Register (NLFSR)
- A : 12×32 -bit words, B and C : 16×32 -bit words.
- $\mathcal{U}(x) = 3 \times x \bmod 2^{32}$, $\mathcal{V}(x) = 5 \times x \bmod 2^{32}$, \lll : left rotation.



- (Final update of A omitted on the figure)

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Unfolding Method

- Unroll the Shabal core
- Method already used for hash functions [Hoare *et al.*, IASTED 2002] [Lien *et al.*, CT-RSA 2004] [Lee *et al.*, JSPF 2008].
- Factor 1

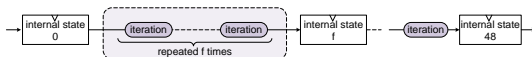


- Factor 2

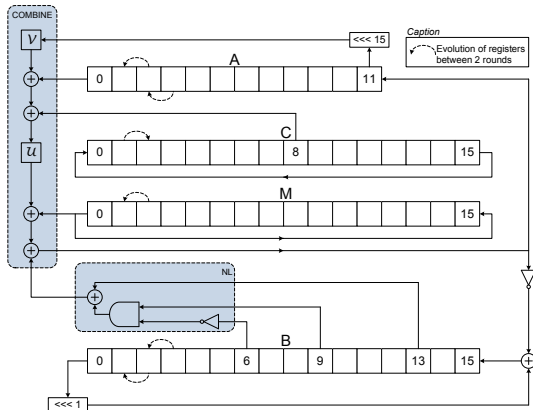


- ...

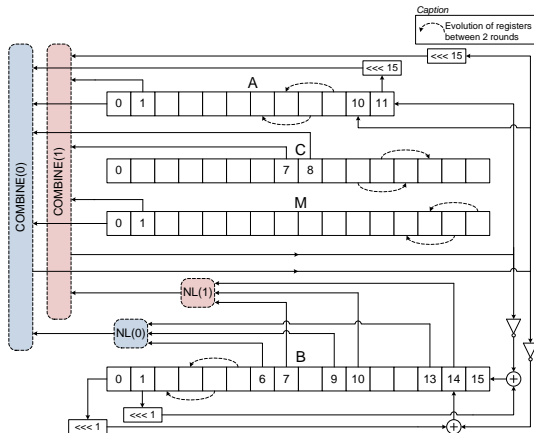
- Factor f



Let's Recall \mathcal{P} ...



...Unfolding Method on \mathcal{P} with Factor 2



- We let you imagine the complexity of circuits with growing f ...

Why Using Unfolding Method?

- ① Save register loading time.
⇒ Impact on Frequency.
- ② ↘ # Clock Cycles per message block
 - Yet, TP (for long messages) = $\frac{\text{Message block length} \times \text{Frequency}}{\# \text{ Clock Cycles per message block}}$.
 - How can we explain an hypothetical gain?
 - Ex.: when $f = 1 \rightarrow 2$, # Clock Cycles per message block/2.
→ To get TP ↗, Frequency ($f = 2$) > Frequency ($f = 1$) / 2.

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Implementation Results (1/2)

- VHDL
- Xilinx Virtex-5 and ISE Tools (v11.1i)
- Optimization: Speed High
- Software padding
- Post-Place and Route Static Timing Report

Implementation Results (2/2)

Unfolding factor f	Area (slices)	Max. Freq. (MHz)	TP (Gbps)	TP/Area (Mbps/slice)
1	1533	247	2.634	1.718
2	1556	144	3.072	1.974
3	1607	101	3.232	2.011
4	1715	76	3.242	1.890
6	1884	50	3.200	1.698
8	1958	37	3.157	1.612

- More efficient: factor 3.
- Highest TP: factor 4.
- Method limit > 4 .

Improvements on Previous Works

Reference	Area (slices)	Max. Freq. (MHz)	TP (Gbps)	TP/Area (Mbps/slice)
[Baldwin <i>et al.</i>]	2768	138	1.450	0.523
[Feron <i>et al.</i>]	1171	126	2.588	2.210
[Kobayashi <i>et al.</i>]	1251	214	2.282	1.390
[Detrey <i>et al.</i>]	153	256	2.051	13.407
Our work ($f = 4$)	1607	101	3.232	2.011

- Highest TP of the state-of-the-art

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Conclusion and Future Works

- At first sight, the restricted parallelibility of Shabal seems to be a factor of **TP limitation**
- \Rightarrow Fortunately, **unrolling Shabal improves its TP (+23%)!**
- **Future Works**
 - Unroll the **48** iterations and insert flip-flops (**retiming**)
 - **Manual Place and Route**
- Our work is available for free evaluation
 - shabal.com
- **Significant contribution** in the hardware benchmarking of the SHA-3 candidates

A Funny Conclusion...

Contrary to Chabal, Shabal can be fast and lightweight.



Thanks for your attention and your future questions, comments, suggestions, etc.!

SAPHIR2 Project

- SAPHIR2 project supported by the French Agence Nationale de la Recherche (ANR-08-VERS-014)
 - Security and Analysis of Primitives of Hashing Innovatory and Recent 2
- **Partners:** ANSSI, Cryptolog International, EADS Defence & Security (CSCSC), France Télécom, Gemalto, INRIA, LIENS, Sagem Sécurité, UVSQ.
- **Goals:**
 - follow and participate to the NIST SHA-3 Contest (cryptanalysis, implementations, optimizations, etc.),
 - to support the candidates proposed by the SAPHIR projects 1 and 2.

Partenaires SAPHIR2

