

Dumbo, Jumbo, and Delirium: Parallel AEAD for the Lightweight Circus

Tim Beyne¹, Yu Long Chen¹, Christoph Dobraunig², Bart Mennink²

¹ KU Leuven (Belgium)

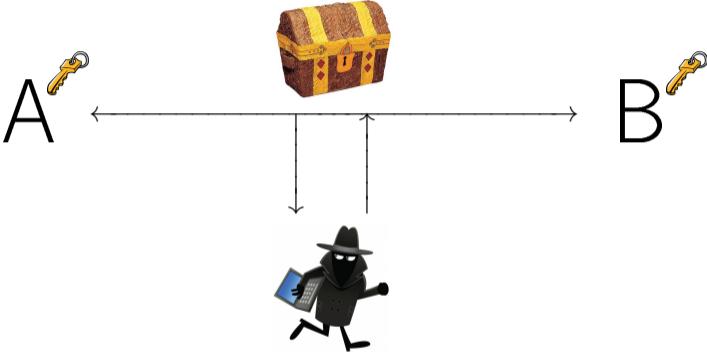
² Radboud University (The Netherlands)

NIST Lightweight Cryptography Workshop 2019
November 6, 2019

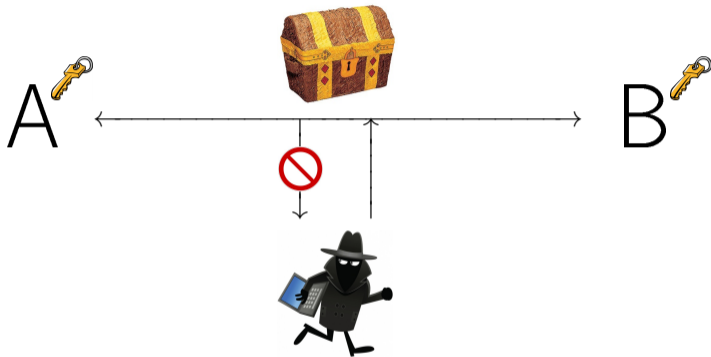
Authenticated Encryption



Authenticated Encryption



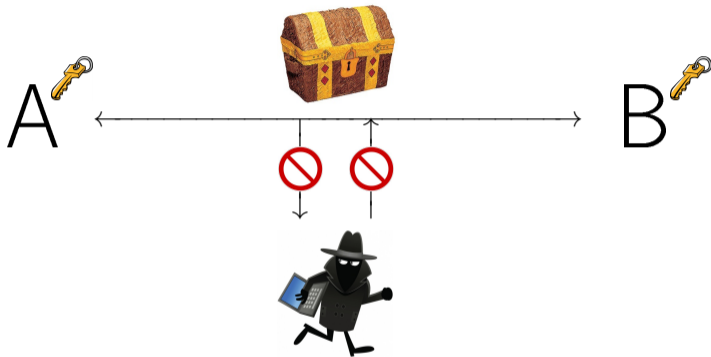
Authenticated Encryption



Encryption

- No outsider can learn anything about data

Authenticated Encryption



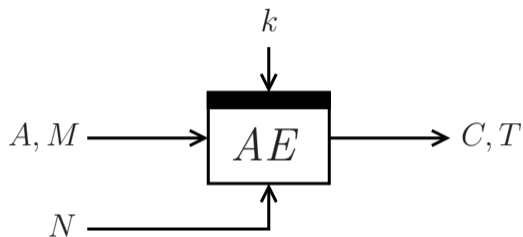
Encryption

- No outsider can learn anything about data

Authentication

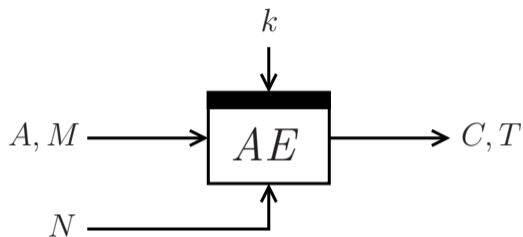
- No outsider can manipulate data

Authenticated Encryption



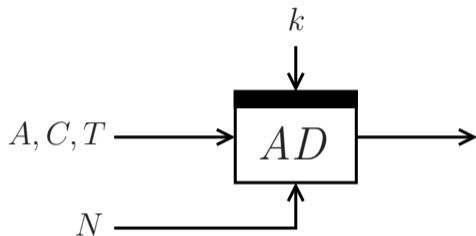
- Ciphertext C encryption of message M
- Tag T authenticates associated data A and message M

Authenticated Encryption



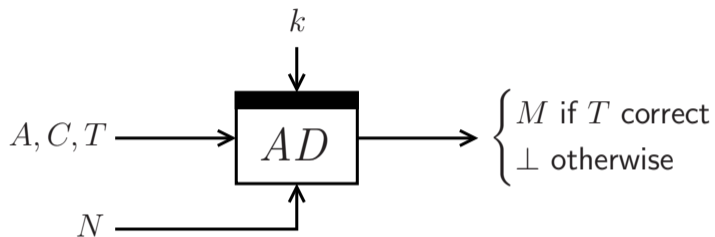
- Ciphertext C encryption of message M
- Tag T authenticates associated data A and message M
- Nonce N randomizes the scheme

Authenticated Decryption



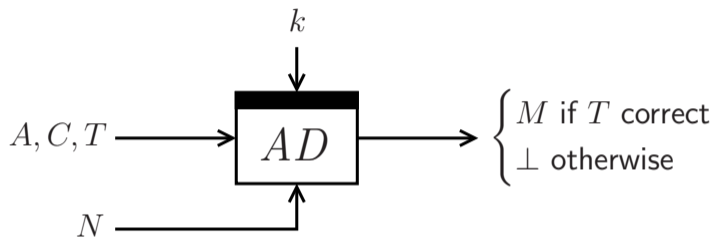
- Authenticated decryption needs to satisfy that
 - Message disclosed if tag is **correct**
 - Message is not leaked if tag is **incorrect**

Authenticated Decryption



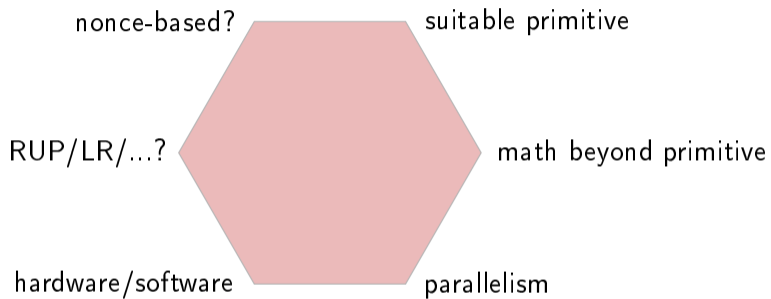
- Authenticated decryption needs to satisfy that
 - Message disclosed if tag is **correct**
 - Message is not leaked if tag is **incorrect**

Authenticated Decryption

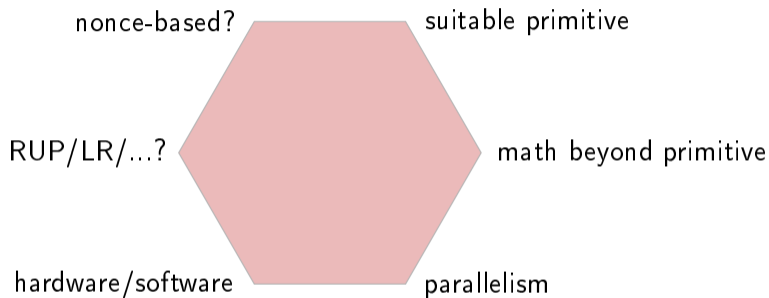


- Authenticated decryption needs to satisfy that
 - Message disclosed if tag is **correct**
 - Message is not leaked if tag is **incorrect**
- Correctness: $AD_k(N, A, AE_k(N, A, M)) = M$

Lightweight Authenticated Encryption



Lightweight Authenticated Encryption



Our goal: minimize state size and complexity of design while still meeting expected security strength 2^{112} and limit on online complexity 2^{50} bytes

What Primitive?

Tweakable Block Cipher



Block Cipher



Permutation



What Primitive?

Tweakable Block Cipher



Block Cipher



Permutation

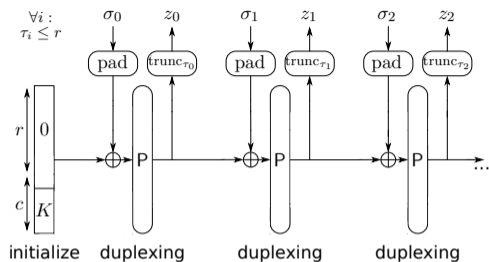


Permutation is the best suited choice

What Mode?

Established Approach

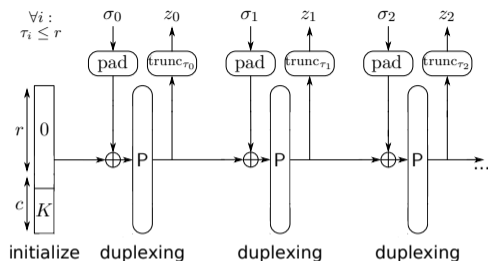
- Keyed duplex/sponge [BDPV11,MRV15,DMV17]
- Inherently sequential



What Mode?

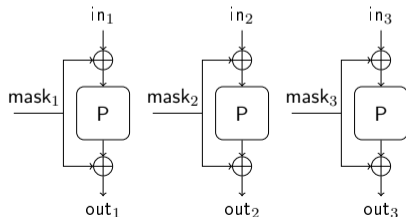
Established Approach

- Keyed duplex/sponge [BDPV11,MRV15,DMV17]
- Inherently sequential



Our Approach

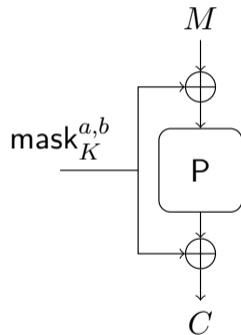
- Parallel evaluation of the permutation
→ requires proper masking
- Evaluating it in forward direction only
→ requires proper mode of use
- Goal: minimize permutation size



What Mask?

Simplified Version of MEM [GJMN16]

- φ_1 is fixed LFSR, $\varphi_2 = \varphi_1 \oplus \text{id}$
- $\text{mask}_K^{a,b} = \varphi_2^b \circ \varphi_1^a \circ P(K \| 0^{n-k})$



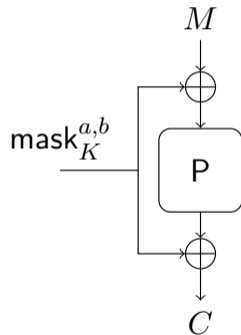
What Mask?

Simplified Version of MEM [GJMN16]

- φ_1 is fixed LFSR, $\varphi_2 = \varphi_1 \oplus \text{id}$
- $\text{mask}_K^{a,b} = \varphi_2^b \circ \varphi_1^a \circ P(K \| 0^{n-k})$

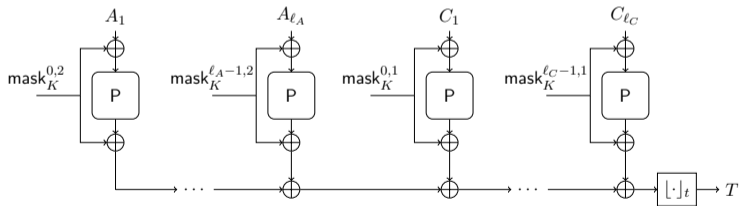
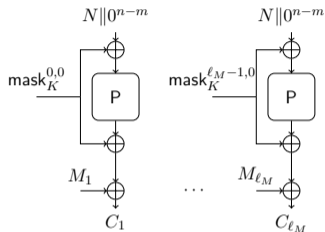
Features

- Constant-time
- Simple to implement
- More efficient than alternatives



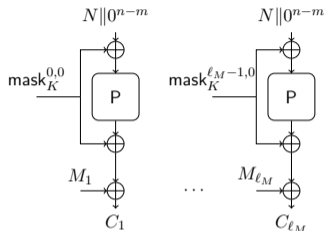
Elephant Authenticated Encryption Mode

$$\text{mask}_K^{a,b} = \varphi_2^b \circ \varphi_1^a \circ P(K \parallel 0^{n-k})$$



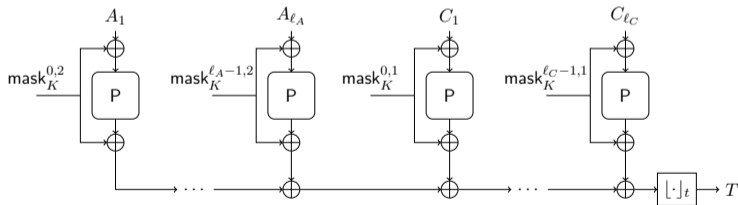
Elephant Authenticated Encryption Mode

$$\text{mask}_K^{a,b} = \varphi_2^b \circ \varphi_1^a \circ P(K \| 0^{n-k})$$



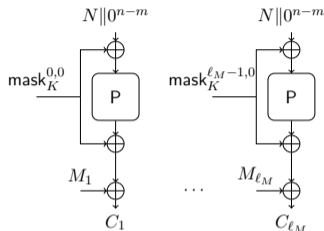
Encryption

- Nonce N input to all P calls
- K and counter in mask
- Padding $M_1 \dots M_{\ell_M} \xleftarrow{n} M$
- Ciphertext $C \leftarrow [C_1 \dots C_{\ell_M}]_{|M|}$



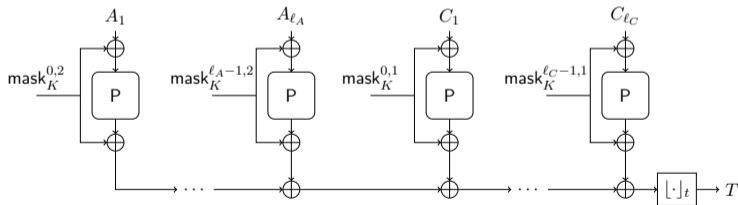
Elephant Authenticated Encryption Mode

$$\text{mask}_K^{a,b} = \varphi_2^b \circ \varphi_1^a \circ P(K \| 0^{n-k})$$



Encryption

- Nonce N input to all P calls
- K and counter in mask
- Padding $M_1 \dots M_{\ell_M} \stackrel{n}{\leftarrow} M$
- Ciphertext $C \leftarrow [C_1 \dots C_{\ell_M}]_{|M|}$

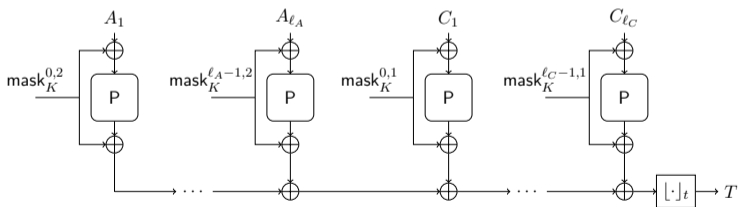
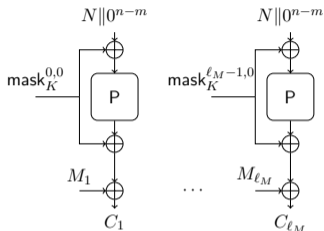


Authentication

- Padding $A_1 \dots A_{\ell_A} \stackrel{n}{\leftarrow} N \| A \| 1$
- Padding $C_1 \dots C_{\ell_C} \stackrel{n}{\leftarrow} C \| 1$
- K and counter in mask
- Tag T truncated to t bits

Elephant Authenticated Encryption Mode

$$\text{mask}_K^{a,b} = \varphi_2^b \circ \varphi_1^a \circ P(K \| 0^{n-k})$$

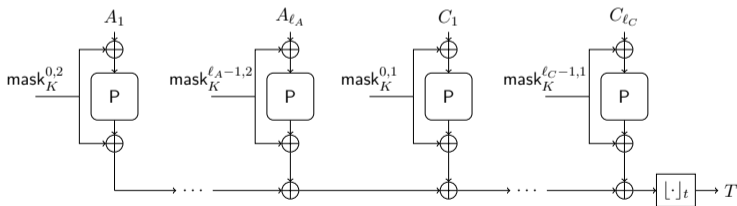
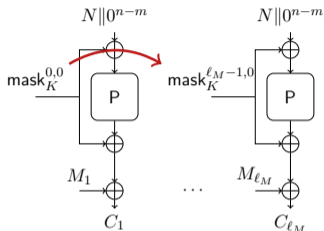


Mode Properties

- Encrypt-then-MAC
 - CTR encryption
 - Wegman-Carter-Shoup
- Fully parallelizable
- Uses single primitive P
- P in forward direction only

Elephant Authenticated Encryption Mode

$$\text{mask}_K^{a,b} = \varphi_2^b \circ \varphi_1^a \circ P(K \| 0^{n-k})$$



Mode Properties

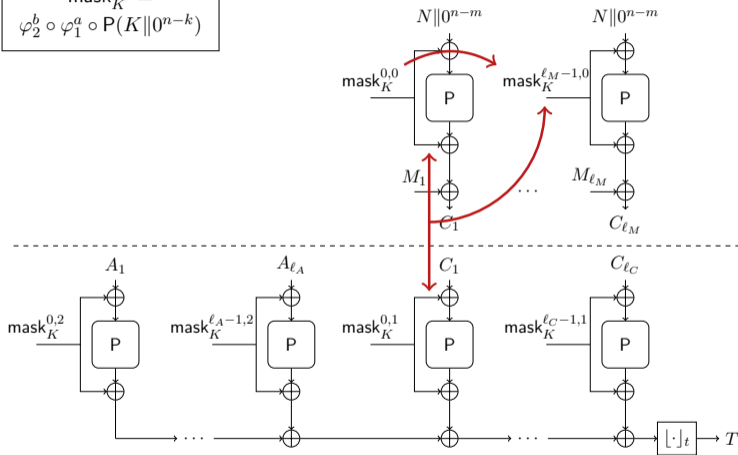
- Encrypt-then-MAC
 - CTR encryption
 - Wegman-Carter-Shoup
- Fully parallelizable
- Uses single primitive P
- P in forward direction only

Mask Properties

- Mask can be easily updated
- $\text{mask}_K^{i,0} = \varphi_1 \circ \text{mask}_K^{i-1,0}$

Elephant Authenticated Encryption Mode

$$\text{mask}_K^{a,b} = \varphi_2^b \circ \varphi_1^a \circ P(K \| 0^{n-k})$$



Mode Properties

- Encrypt-then-MAC
 - CTR encryption
 - Wegman-Carter-Shoup
- Fully parallelizable
- Uses single primitive P
- P in forward direction only

Mask Properties

- Mask can be easily updated
- $\text{mask}_K^{i,0} = \varphi_1 \circ \text{mask}_K^{i-1,0}$
- $\text{mask}_K^{i-1,0} \oplus \text{mask}_K^{i-1,1} = \text{mask}_K^{i,0}$

Security of Mode

$$\mathbf{Adv}_{\text{Elephant}}^{\text{ae}}(\mathcal{A}) \lesssim \frac{4\sigma p}{2^n}$$

- σ is online complexity, p is offline complexity
- Assumptions:
 - P is random permutation
 - φ_1 has maximal length and $\varphi_2^b \circ \varphi_1^a \neq \varphi_2^{b'} \circ \varphi_1^{a'}$ for $(a, b) \neq (a', b')$
 - \mathcal{A} is nonce-based adversary

Security of Mode

$$\mathbf{Adv}_{\text{Elephant}}^{\text{ae}}(\mathcal{A}) \lesssim \frac{4\sigma p}{2^n}$$

- σ is online complexity, p is offline complexity
- Assumptions:
 - P is random permutation
 - φ_1 has maximal length and $\varphi_2^b \circ \varphi_1^a \neq \varphi_2^{b'} \circ \varphi_1^{a'}$ for $(a, b) \neq (a', b')$
 - \mathcal{A} is nonce-based adversary

Parameters of NIST lightweight call
can be met with a 160-bit permutation!

Instantiation



Dumbo

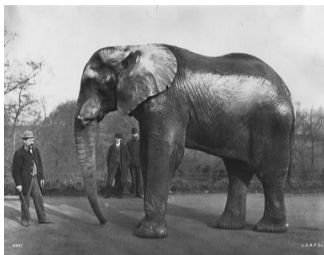
- Spongent- π [160]
- Minimalist design
 - Time complexity 2^{112}
 - Data complexity 2^{46}

Instantiation



Dumbo

- Spongnet- π [160]
- Minimalist design
 - Time complexity 2^{112}
 - Data complexity 2^{46}



Jumbo

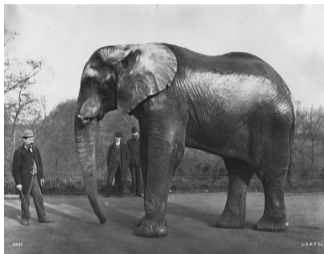
- Spongnet- π [176]
- Conservative design
 - Time complexity 2^{127}
 - Data complexity 2^{46}
- ISO/IEC standardized

Instantiation



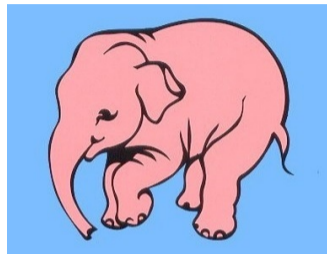
Dumbo

- Spongent- π [160]
- Minimalist design
 - Time complexity 2^{112}
 - Data complexity 2^{46}



Jumbo

- Spongent- π [176]
- Conservative design
 - Time complexity 2^{127}
 - Data complexity 2^{46}
- ISO/IEC standardized



Delirium

- Keccak- f [200]
- High security
 - Time complexity 2^{127}
 - Data complexity 2^{70}
- NIST standardized

Technical Specification of Instances

instance	k	m	n	t	P	φ_1	expected security strength	limit on online complexity
Dumbo	128	96	160	64	80-round Spongent- π [160]	φ_{Dumbo}	2^{112}	$2^{50}/(n/8)$
Jumbo	128	96	176	64	90-round Spongent- π [176]	φ_{Jumbo}	2^{127}	$2^{50}/(n/8)$
Delirium	128	96	200	128	18-round Keccak- f [200]	$\varphi_{\text{Delirium}}$	2^{127}	$2^{74}/(n/8)$

- All LFSRs operate on 8-bit words:

$$\varphi_{\text{Dumbo}} : (x_0, \dots, x_{19}) \mapsto (x_1, \dots, x_{19}, x_0 \lll 3 \oplus x_3 \ll 7 \oplus x_{13} \gg 7)$$

$$\varphi_{\text{Jumbo}} : (x_0, \dots, x_{21}) \mapsto (x_1, \dots, x_{21}, x_0 \lll 1 \oplus x_3 \ll 7 \oplus x_{19} \gg 7)$$

$$\varphi_{\text{Delirium}} : (x_0, \dots, x_{24}) \mapsto (x_1, \dots, x_{24}, x_0 \lll 1 \oplus x_2 \lll 1 \oplus x_{13} \ll 1)$$

- All have maximal length and $\varphi_2^b \circ \varphi_1^a \neq \varphi_2^{b'} \circ \varphi_1^{a'}$ for $(a, b) \neq (a', b')$

Conclusion

Elephant

- Parallel lightweight AE with small state
- Mode: provably secure in random permutation model
- Primitives: standardized and well-studied
- **Dumbo** and **Jumbo** for hardware
- **Delirium** for software

Thank you for your attention!