

On Implementation Security and ISAP v2.0

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NIST Lightweight Cryptography Workshop 2020

- Prime applications of LWC are:
 - Constraint devices (low computing power and memory)
 - Industry use cases (require protection from physical attacks)
- Many NIST LWC submissions use lightweight building blocks
 - Reduce cost of masking
- But the story doesn't quite end here ...

- Passive Implementation Attacks
 - DPA, **SPA**
- Fast & Compact Co-Processor for *ASCON* and *ISAP*
 - Decent implementation security from unprotected building blocks
- **Active Implementation Attacks**
 - DFA, SFA, SIFA
- Summary

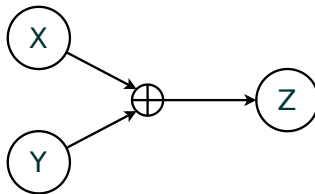
Passive Implementation Attacks

- Requires varying inputs or outputs [KJJ99]
- Exploits differences in power consumption
- Countermeasures:
 - Implementation-level: (Higher-order) masking [GP99]
 - Mode-level: GGM-like constructions [GGM84]

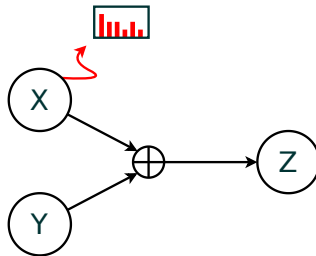
- DPA attack on plaintext during authenticated decryption
 - Without knowledge of the key
- Keys are not the only asset → plaintext could contain:
 - Keys
 - Firmware
 - Intellectual Property
- Countermeasures:
 - Implementation-level: Masking (careful with leveled implementations)
 - Mode-level: Verify authenticity of ciphertext/nonce before decryption

- Similar: Profiled Attacks, Template Attacks
- Can be super powerful!
- How they work:
 1. Characterize target device → build templates
 2. Look at few/many different intermediate steps of a computation
 3. Match templates and get noisy leakages
 4. Combine leakage using multivariate PDFs/belief propagation

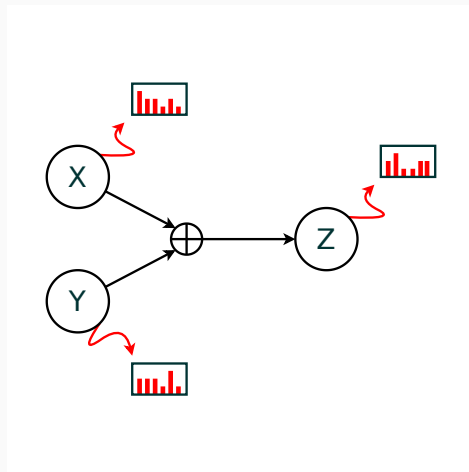
- Two types of nodes
 - Variable nodes
 - Factor nodes (here: XOR)



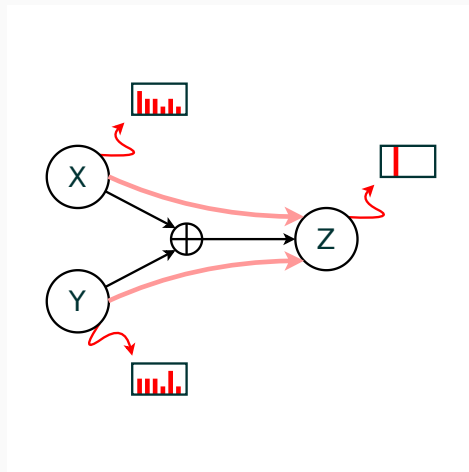
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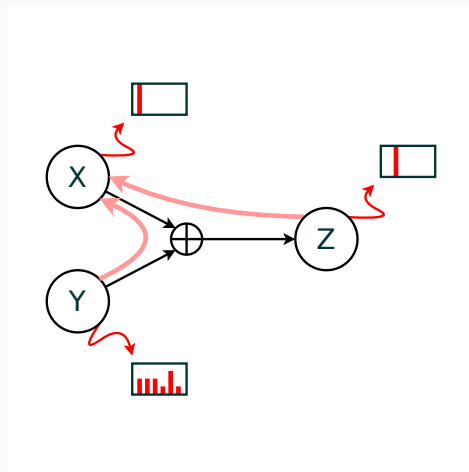
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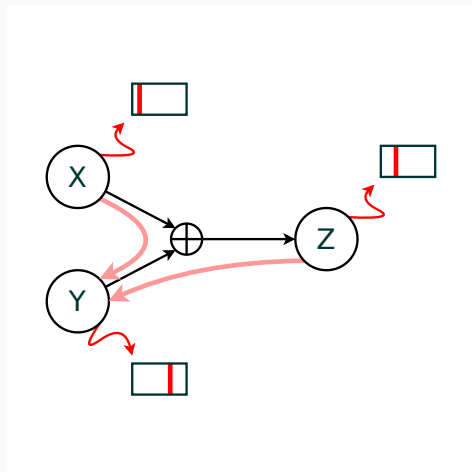
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 - Passes information between both types of nodes
- ⇒ Combine leakage information



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- Presented at CHES 2020 [KPP20]
- Accumulate and combine **hundreds/thousands** of leakage points
 - State recoveries possible for 8, 16, (32)-bit devices
- Similar results in other publications [VGS14; PP19; Bel+20]
 - Also work against masked implementations
- Downside?
 - Attacker requires precise knowledge about algorithm execution

- Masking alone is not very effective:
 - SPA attacks on 1st-order masking [OM07; HTM09]
 - SPA attacks on higher-order masking [PPM17]
- Hiding is more effective:
 - Uncertainty about when computations occur (belief propagation [Rav+20])
 - Increased SNR on low-end devices

Fast&Compact Co-Processor for Ascon and ISAP

- To be presented at CARDIS 2020 [SP20]
- Tight coupling of *ASCON-p* to processors register file
 - No additional registers needed
 - No interface needed

⇒ Half the area of dedicated co-processor designs
- Mode remains entirely in software → flexible:
 - *ASCON*, *ASCON-HASH*, *ASCON-XOF*
 - *ISAP-A-128A*, *ISAP-A-128*
- Speed-ups by a factor of about 50 to 80

Table 1: Runtime and code size comparison of ASCON and ISAP, with/without 1-round ASCON- p hardware acceleration (HW-A)

Implementations	Cycles/Byte			Binary Size (B)
	64 B	1536 B	long	
Ascon-C (-03)	164.3	110.6	108.3	11 716
Ascon-C (-0s)	269.7	187.1	183.5	2 104
Ascon-ASM + HW-A	4.2	2.2	2.1	888
AsconHash-ASM + HW-A	4.6	2.6	2.5	484
ISAP-A-128a-C (-03)	1 184.3	386.9	352.3	11 052
ISAP-A-128a-ASM + HW-A	29.1	5.2	4.2	1 844

Table 2: Area comparison of the RISC-V RI5CY core and various co-processor designs

Design	kGE	
	Standalone	Integration
RI5CY base design	45.6	-
This work	4.2	0.5
ASCON co-processor [Gro+15]	7.1	?
ASCON co-processor [Gro]	9.4	?
ISAP co-processor (estimated) [Dob+19]	≤ 12.8	?

- With respect to passive attacks, ISAP provides protection against ...
 - DPA protection (incl. plaintext recovery)
- Co-processor provides basic protection against SPA attacks ...
 - Attacker is essentially forced to perform localized EM analysis
 - Simple steps to further improve SPA protection are discussed [SP20]
- ISAP also offers protection against active attacks ...

Active Implementation Attacks

- Requires constant (unknown) inputs
- Exploits differences between correct/faulty computations [Bar+06]
- Generally applicable to AEAD if either:
 - Nonce is repeated (Enc)
 - Unverified plaintext is released (Dec)
 - Tag verification before decryption (Dec)
- Countermeasures:
 - Implementation-level: Redundant computation (temporal/spatial)
 - Mode-level protection possible

- Exploits faulty computations only [Fuh+13]
- Requires varying (unknown) inputs
- Applicable to AEAD if:
 - Key addition occurs before output (e.g. tag or ciphertext)
- Countermeasures:
 - Implementation-level: Redundant computation (temporal/spatial)
 - Mode-level protection possible

- Exploits faulted but correct computations only [Dob+18]
- Requires varying (known¹) inputs
- Countermeasures:
 - Implementation-level: Specific combination of redundancy and masking [Dae+20]
 - Mode-level protection possible

¹When considering attacks on AEAD.

Summary

Dealing with typical implementation attacks purely on algorithmic level:

Attack Type	Primary Countermeasure	Comment
DPA	Masking	Low-end: Higher-order
SPA	Hiding	Shuffling, Dummy Ops.
DFA/SFA	Redundancy	Spatial/Temporal
SIFA	(Masking + Redundancy)	[Dae+20]

Least common multiple: Hiding + Masking + Redundancy

(Remark: Masking + Redundancy alone is not very effective against SPA [OM07; HTM09; PPM17])

What can be done on mode-level (ISAP):

Attack Type	Primary Countermeasure	Comment
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DFA/SFA	Redundancy	Spatial/Temporal
SIFA	(Masking + Redundancy)	[Dae+20]

Least common multiple: Hiding

(Remark: Shuffling + “cheap” masking can also give an effective hiding scheme)

- ISAP is fast and compact in applications that require implementation security
- Goes well with other schemes utilizing $\text{KECCAK-}p[400]$ or $\text{ASCON-}p$
 - Fast AEAD without implementation security
 - Hashing functionality
 - (See e.g. compact co-processor design)

Thank you!

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