From: MEGE, Alexandre <alexandre.mege@airbus.com>

Sent: Monday, June 3, 2019 5:32 AM

To: lightweight-crypto
Cc: lwc-forum@list.nist.gov

Subject: OFFICIAL COMMENT: Fountain

Dear All,

I have found some tag collision with Fountain reference code.

It seems those collisions are produced when the middle byte of AD changes from 0x80 to 0x01, with same pattern as PT on left and right of this Byte.

Examples:

Key = 000102030405060708090A0B0C0D0E0F

Nonce = 000102030405060708090A0B

PT = 000001010202030304040505060607

AD = 00000101020203030404050506060780000001010202030304040505060607

CT = 5274801E6C0131A255F74245E8C0F3384B52EEF66E37734A0789C378FC0A26

Key = 000102030405060708090A0B0C0D0E0F

Nonce = 000102030405060708090A0B

PT = 000001010202030304040505060607

AD = 00000101020203030404050506060701000001010202030304040505060607 CT = 5274801E6C0131A255F74245E8C0F3384B52EEF66E37734A0789C378FC0A26

And

Key = 000102030405060708090A0B0C0D0E0F

Nonce = 000102030405060708090A0B

PT = 000001

AD = 00000180000001

CT = 62F8E554A1CADFA82D4BF736FDD03550DCABFC

Key = 000102030405060708090A0B0C0D0E0F

Nonce = 000102030405060708090A0B

PT = 000001

AD = 00000101000001

CT = 62F8E554A1CADFA82D4BF736FDD03550DCABFC

Best regards, Alexandre Mège **From:** Zhang Bin <martin_zhangbin@hotmail.com>

Sent: Monday, June 3, 2019 9:58 AM **To:** MEGE, Alexandre; lightweight-crypto

Cc: lwc-forum@list.nist.gov

Subject: 回复: OFFICIAL COMMENT: Fountain

Dear Alexandre,

Thanks for the interest on Fountain and for pointing out this implementation problem in Fountain's reference code, which is resulted from some typos on LFSR numbers inconsistent with the specification. Fountain is **invertible** by the design. The reference code will be revised soon and there is no change of Fountain's specification.

Best Regards,

Bin Zhang

发件人: MEGE, Alexandre <alexandre.mege@airbus.com>

发送时间: 2019 年 6 月 3 日 17:31 收件人: lightweight-crypto@nist.gov

抄送: lwc-forum@list.nist.gov

主题: [lwc-forum] OFFICIAL COMMENT: Fountain

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CT = 5274801E6C0131A255F74245E8C0F3384B52EEF66E37734A0789C378FC0A26

And

Key = 000102030405060708090A0B0C0D0E0F

From: raluca.e.posteuca@gmail.com
Sent: Thursday, June 6, 2019 3:33 AM

To: lightweight-crypto
Cc: lwc-forum@list.nist.gov

Subject: OFFICIAL COMMENT: Fountain

Dear author,

We have analyzed Fountain and we found the following results regarding its security.

For each (Key, IV) pair, with probability $2^{(-32)}$, after 48 steps of the initialization phase, there exists an internal state with the same properties as the output of the loading phase (the constants are in the same positions). Based on this observation, we have designed a slide attack, that leads to predictable ciphertexts (up to 40 bits long) with a time complexity of up to $12 \times (2 \times 80)$. Moreover, computing the ciphertexts' XOR, we can also recover up to 8 bits of the internal state and up to 4 equations involving bits of the internal state.

Best regards,

Raluca Posteuca

COSIC, KU Leuven



Virus-free. www.avast.com

From: Zhang Bin <martin_zhangbin@hotmail.com>

Sent: Thursday, June 6, 2019 10:21 PM

To: raluca.e.posteuca@gmail.com; lightweight-crypto

Cc: lwc-forum@list.nist.gov

Subject: 回复: [lwc-forum] OFFICIAL COMMENT: Fountain

Attachments: New reference code of Fountain v1.zip

Dear all,

Please find the new reference code for Fountain and the report file as attached, which corrected the previous mentioned implementation errors that are inconsistent with the specification, and may affect the reading of the report and cross-checking of the design. There is no change of Fountain's specification, only the test vectors in Chapter 7 of the report file are revised accordingly. There is also a Chapter 8 in the document to record the change log.

Please refer to the attached package for the details. Thanks.

Dear Raluca,

Thanks for the interest on Fountain and the analysis. I would appreciate if the details of the mentioned attack could be provided.

Best Regards,

Bin Zhang

发件人: raluca.e.posteuca@gmail.com <raluca.e.posteuca@gmail.com>

发送时间: 2019 年 6 月 6 日 15:32 收件人: lightweight-crypto@nist.gov

抄送: lwc-forum@list.nist.gov

主题: [lwc-forum] OFFICIAL COMMENT: Fountain

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For each (Key, IV) pair, with probability 2^{-32} , after 48 steps of the initialization phase, there exists an internal state with the same properties as the output of the loading phase (the constants are in the same positions). Based on this observation, we have designed a slide attack, that leads to predictable ciphertexts (up to 40 bits long) with a time complexity of up to $12 \times (2 \times 80)$. Moreover, computing the ciphertexts' XOR, we can also recover up to 8 bits of the internal state and up to 4 equations involving bits of the internal state.

Best regards, Raluca Posteuca COSIC, KU Leuven From: Zhang Bin <martin_zhangbin@hotmail.com>

Sent: Thursday, June 20, 2019 10:24 PM

To: raluca.e.posteuca@gmail.com; lightweight-crypto

Cc: lwc-forum@list.nist.gov

Subject: 回复: [lwc-forum] OFFICIAL COMMENT: Fountain

Dear Raluca,

Thanks again for the interest on Fountain. Since the details of the designed slide attack have not been revealed so far, I just present some analysis on the feasibility of launching a slide attack on Fountain.

The first concern is how to identify the slide pairs when the key is unknown and what is the complexity for doing so. Precisely, in the single key model (in which Fountain has claimed 112-bit security), after 48 steps of initialization, the nonlinear feedbacks has affected 192 state bits, and the least significant two bytes of each register are shifted results from itself. If the attacker manipulates the three IV byte positions, IV_0, IV_4 and IV_8 by choosing the appropriate IV's such that IV_0=IV'_3, IV_4=IV'_7 and IV_8=IV'_11 to meet these three bytes sliding conditions, he still has to deal with the other state bits to meet the sliding conditions, that is the two states are slidable if and only if the full 256-bit state conditions are satisfied, not only the 32-bit constants conditions. But now all the other 256-3*8=232 bit positions are either affected by the unknown key or are the shifted constant values and the 192 feedback positions correspond to a non-linear system of algebraic equations. Thus, the probability that (key, IV) and (key, IV') are slide pairs is 2^(-232) after the manipulating the IV values, not 2^(-32) when the key is unknown. Thus to detect such a slide pair with the appropriate IVs, one has to randomly change the value of IVs to meet the 232 bit conditions, which will have a complexity much higher than 2^(112). Besides, if we take into account the data/memory complexities for fulfilling this task, the complexity will definitely exceed the 112-bit bound.

In the relate-key model (in which Fountain has made no claims), similar analysis also holds and it seems difficult to generate the slide pairs and the complexity for doing so is well above the 112-bit security bound.

If we assume the adversary is able to solve the non-linear algebraic system which describes the inherent relation between the sliding state pairs, we cannot assure that the captured keystream is just related to the specific value and the complexity to find the exact sliding partner is well above the claimed security bound.

Thus if I am not wrong, based on the above analysis, Fountain is immune to the slide attacks.

Best Regards,

Bin Zhang

发件人: raluca.e.posteuca@gmail.com <raluca.e.posteuca@gmail.com>

发送时间: 2019 年 6 月 6 日 15:32 收件人: lightweight-crypto@nist.gov

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