

A Keyed Sponge Construction with Pseudorandomness in the Standard Model

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Joint work with

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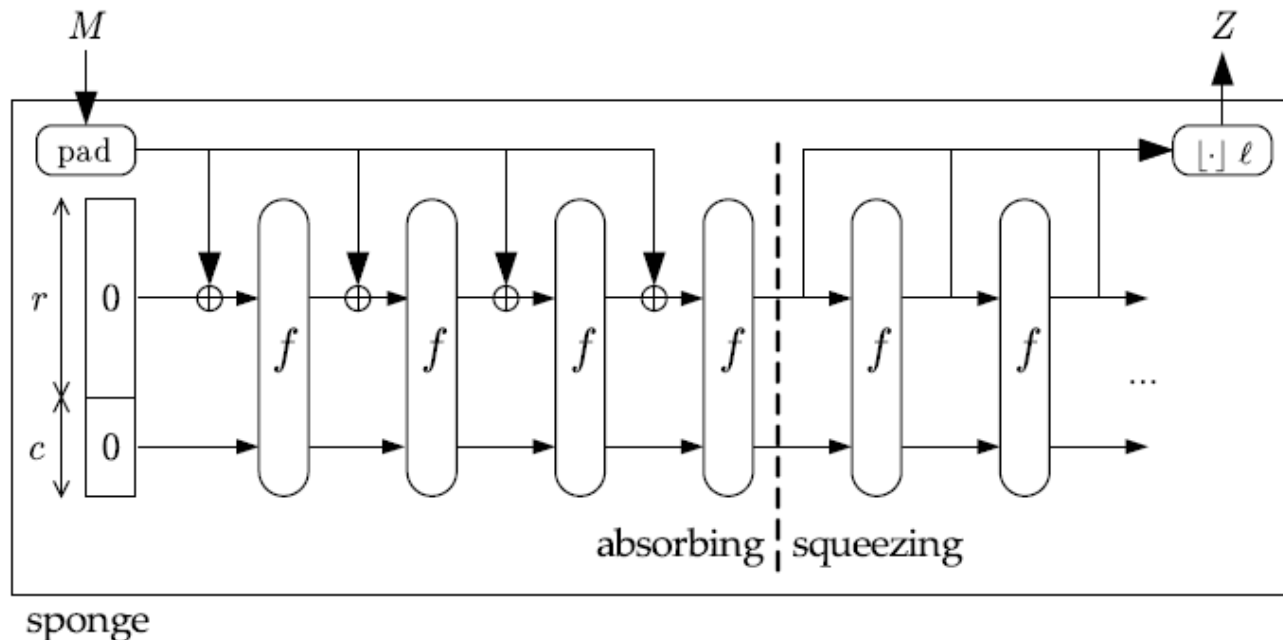
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Sponge Construction

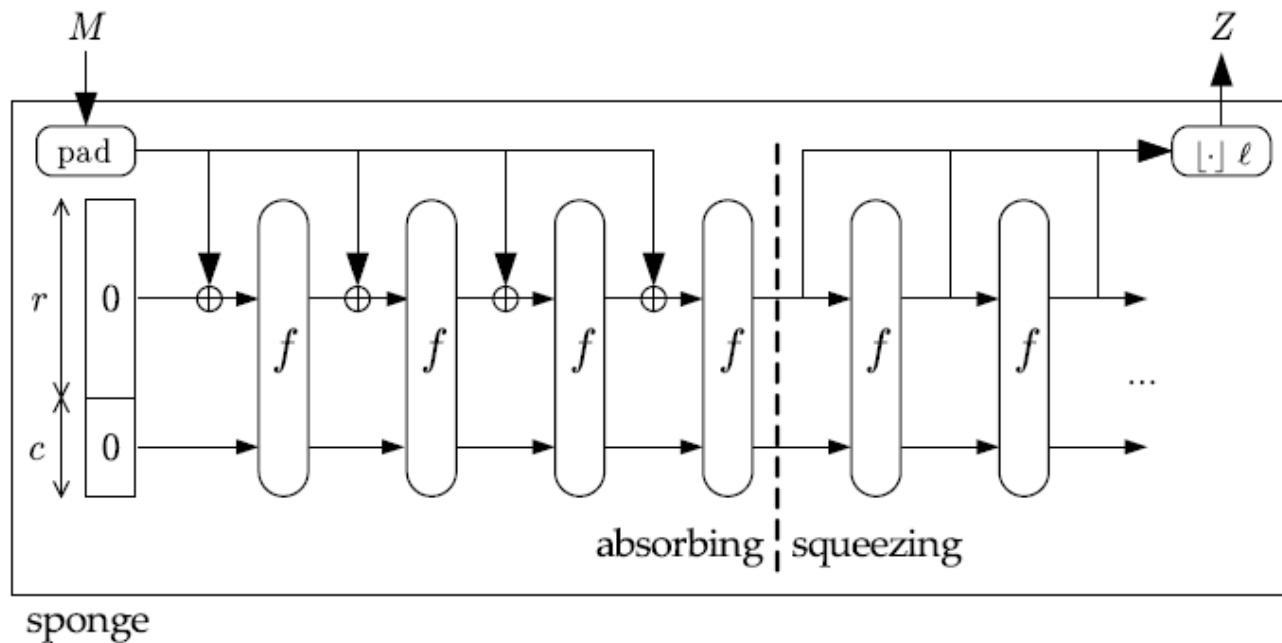
- Designed by Bertoni, Daemen, Peeters, and Van Assche (Eurocrypt '08);
- Influenced concrete hash designs such as Keccak, PHOTON, Quark, and Spongent.



f : a b -bit permutation with $b = r + c$

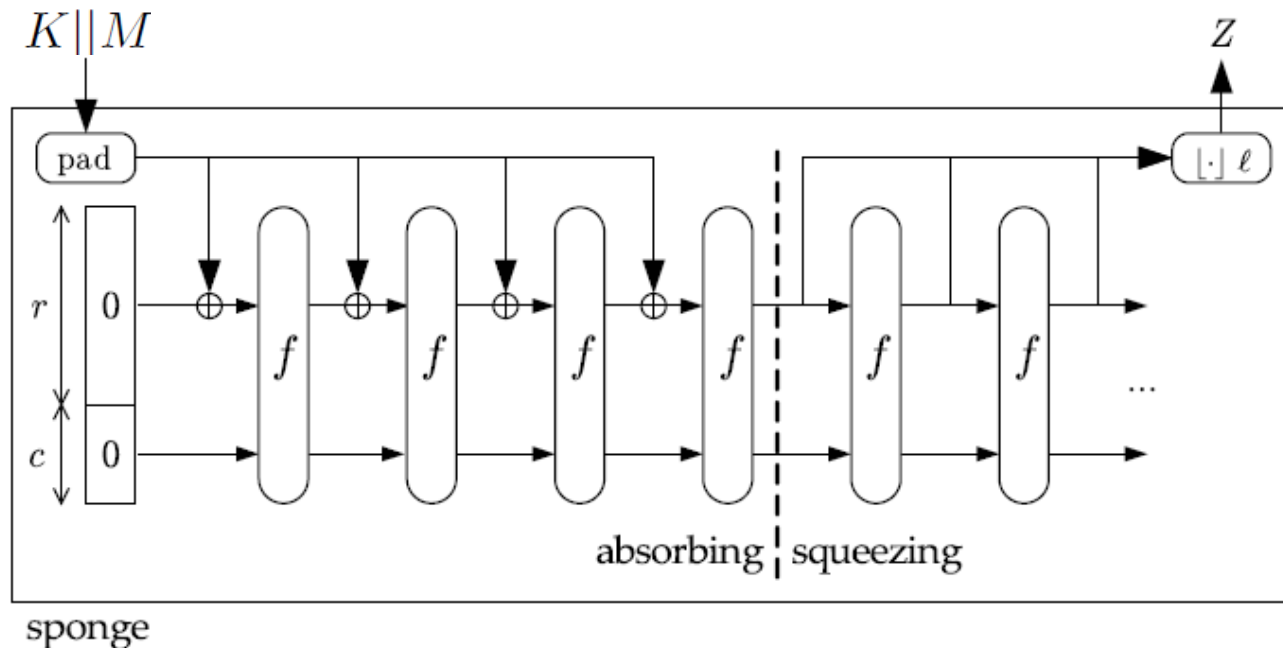
Security of Sponge Construction (Eurocrypt '08)

- Sponge is indifferentiable from a random oracle when f is an ideal permutation or an ideal function (in other words, a fixed-input-length random oracle).



A Keyed Sponge Construction

- Defined by Bertoni, Daemen, Peeters, and Van Assche (SKEW'11).



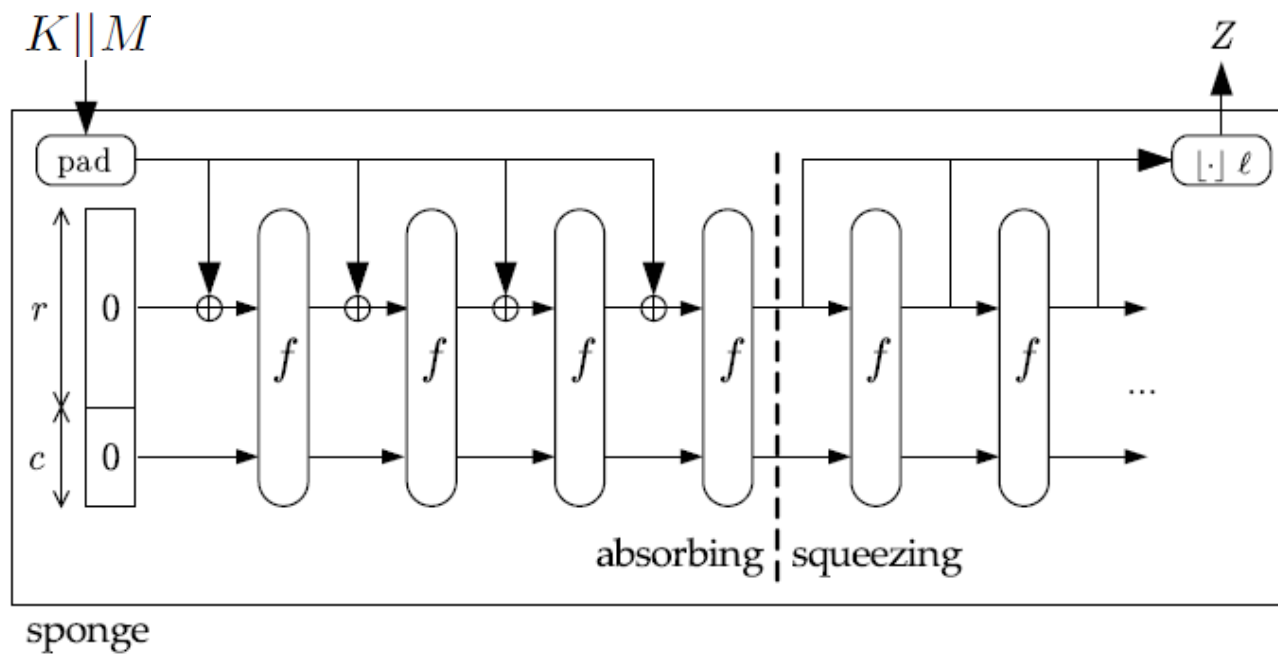
f : a b -bit permutation with $b = r + c$

Applications of A Keyed Sponge Construction (SKEW '11)

- Encryption as a stream cipher
 - Squeezing $\text{sponge}(K||IV)$, or
 - Random-access key stream block $k_i = \text{sponge}(K||IV||i)$.
- Authentication: $\text{Sponge}(K||M)=\text{MAC}$.

Security of Keyed Sponge Construction (SKEW '11)

- Pseudorandomness is proved in the ideal permutation model.

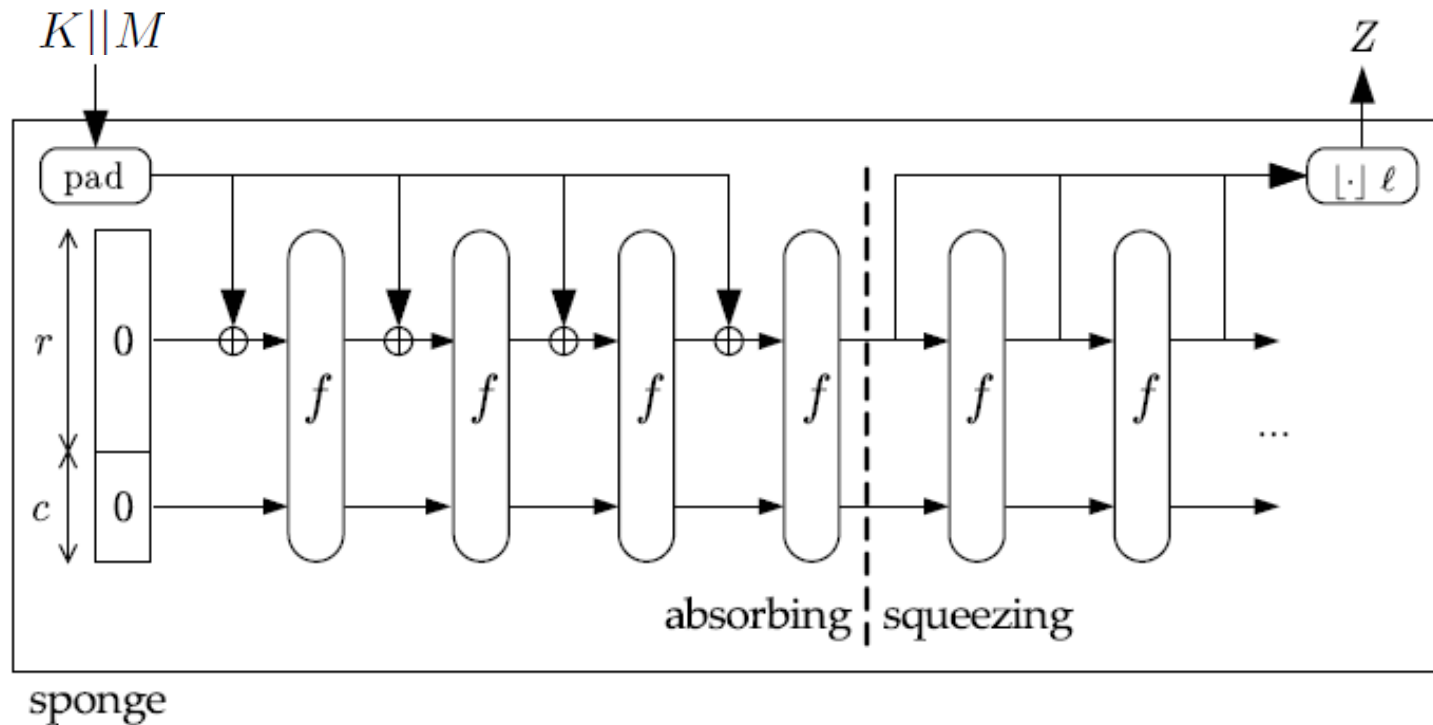


Proof assumes f is an **ideal** permutation.

Our work

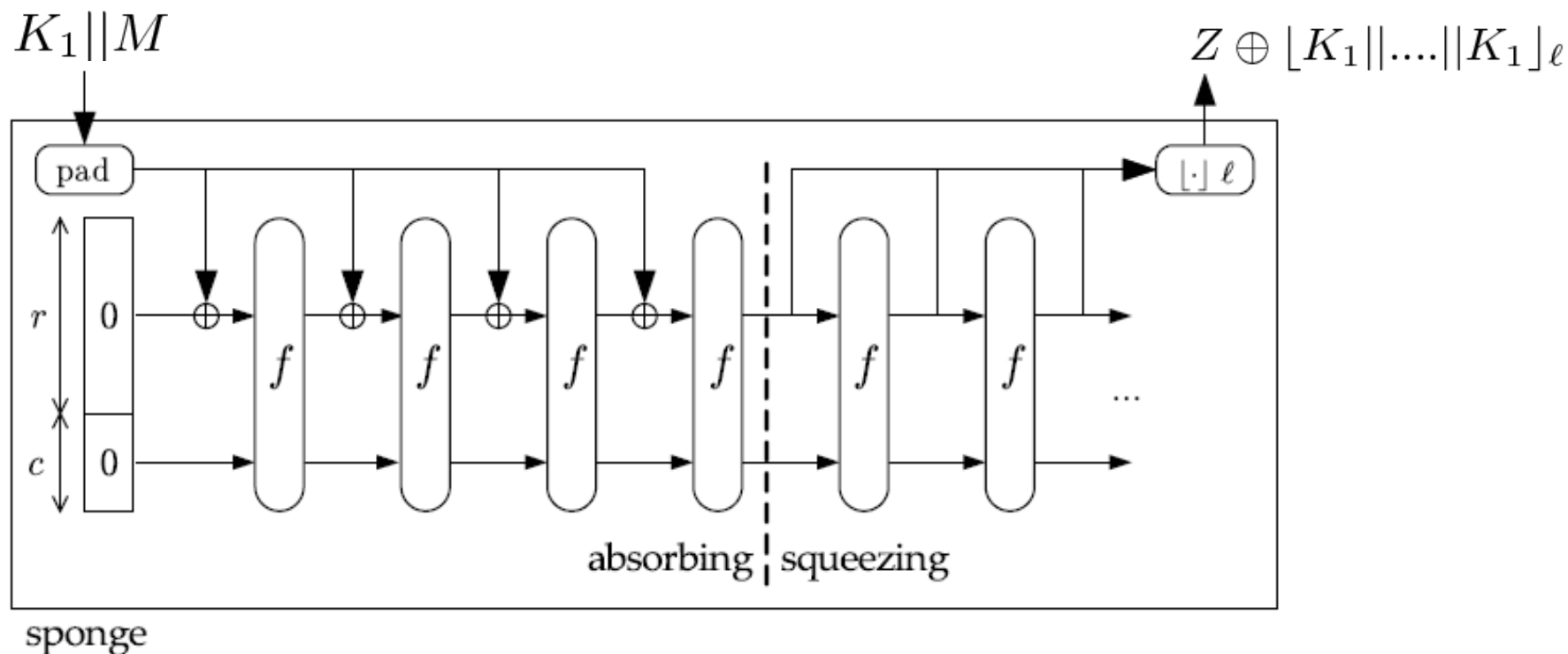
- We give a new keyed sponge construction based on the Even-Mansour permutation.
- We give variants for three key sizes.
- The security of the construction doesn't depend on the ideal model, but on the standard model with a practical assumption.

A Keyed Sponge Construction (SKEW '11)



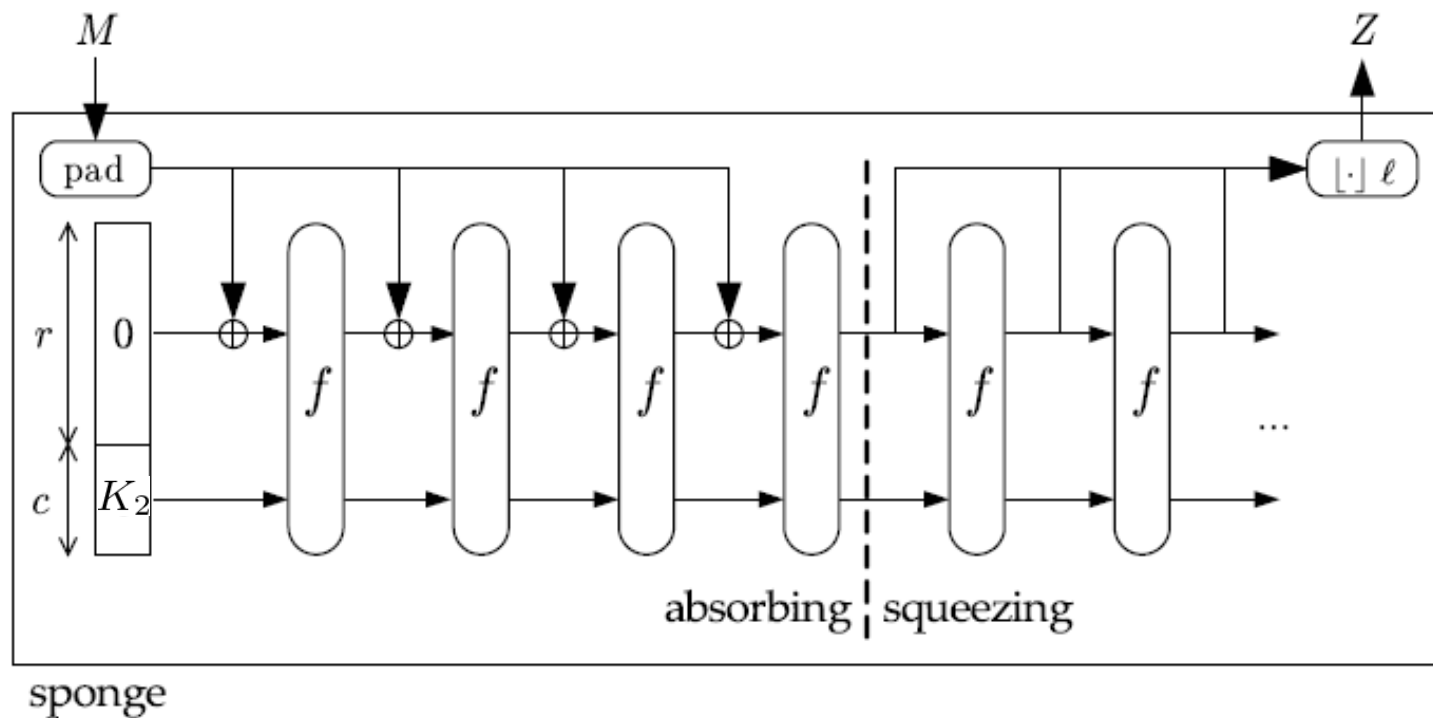
Our Keyed Sponge Construction #1 (No modification to Sponge)

- K_1 is an r -bit secret key.



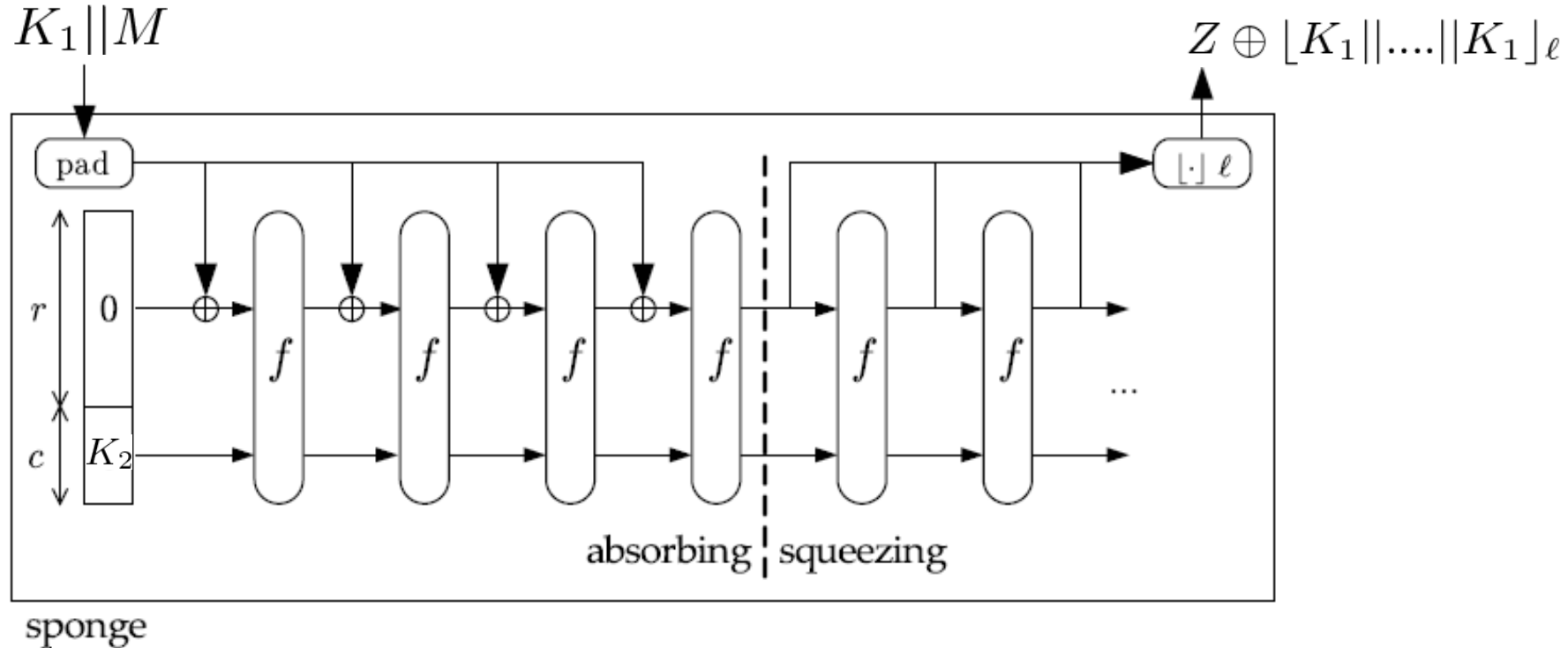
Our Keyed Sponge Construction #2 (Modify the Initial Value)

- K_2 is a c -bit secret key.



Our Keyed Sponge Construction #3 (Combination of #1 and #2)

- K_1 and K_2 are r -bit and c -bit secret keys.

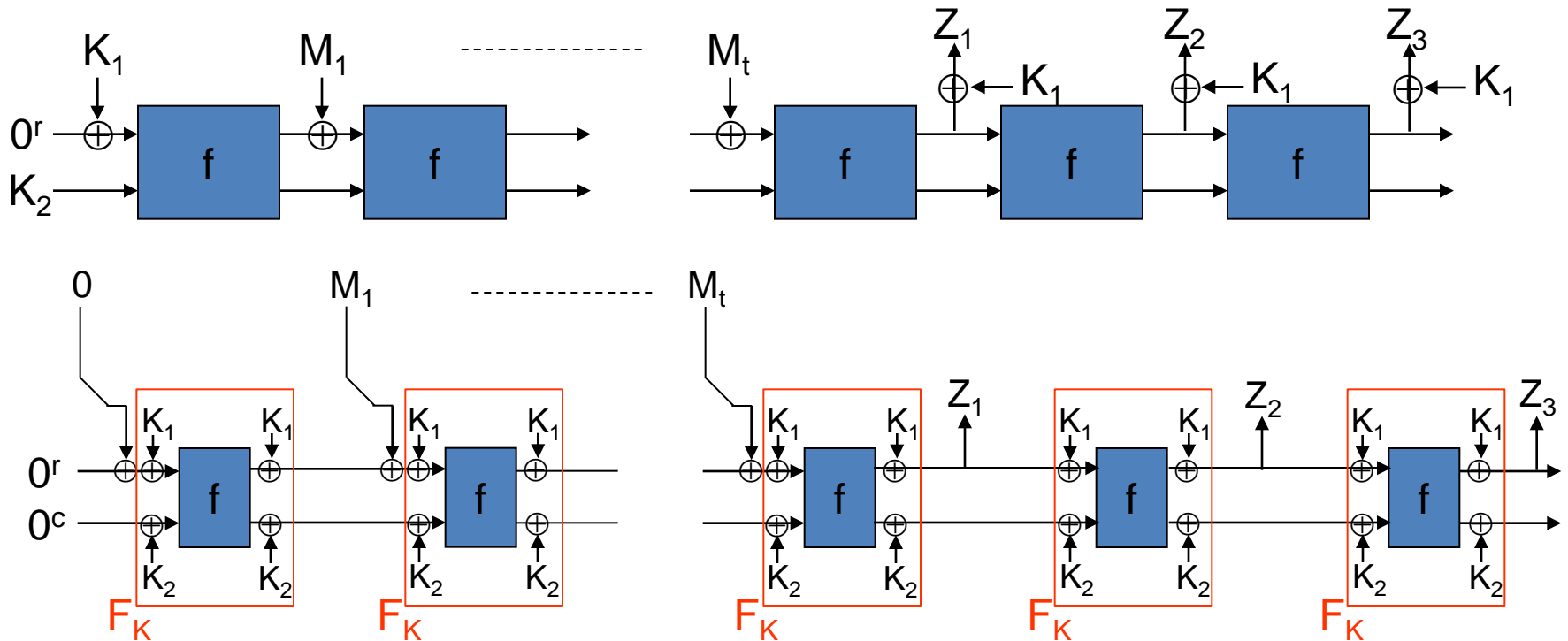


Security Assumption

- The Even-Mansour permutation with a single key is $F_K(\cdot) = f(K \oplus \cdot) \oplus K$.
- Instead of assuming $f(\cdot)$ is an ideal permutation, we assume $F_K(\cdot)$ is pseudorandom.
- If $F_K(\cdot)$ is pseudorandom, then our construction is also pseudorandom.

Underlying Proof Idea

- These two descriptions are the same.



Best-Known Attack on $F_K(\cdot)$

- Due to Dunkelman et al [Eurocrypt '12];
- Assumes $|K|=b$;
- Known plaintext PRP attack on $F_K(\cdot)$;
- Complexity $DT=2^b$, where D and T refer to data and time complexity;
- Generic.

Conclusion

- We showed that a new keyed sponge construction is pseudorandom under the standard model.
- It is an open question whether our technique can be applied to other sponge-like constructions.