First Public DRAFT

# MISPC

# Minimum Interoperability Specification for PKI Components, Version 2 - Second DRAFT

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#### NIST PKI Project Team

For public distribution

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#### **Reviewer's Notes**

This is the first public draft of the MISPC Version 2. Version 2 introduces support for confidentiality services. This support is manifested in extensions to the certificate profile (secs. 3.1.2.3 and 3.1.2.4), new transactions to request key management certificates (secs. 3.5.7 and 3.5.10), combined requests for both signature and key management certificates (sec. 0) and extensions to the revocation request (sec. 3.5.6.) Less drastic changes are sprinkled throughout.

Version 2 realigns the MISPC with the most recent release of the ISO X.509 specification. New certificate and CRL extensions have been added to this specification. The new extensions appear in sections 3.1.3, 3.2.2, and 3.2.3.

Version 2 also realigns the MISPC with the current IETF specifications for the Certificate Management Protocol (CMP) and Certificate Request Message Format (CRMF), ongoing developments in the FIPS approved cryptographic algorithms, and X9 cryptographic algorithms. (The text surrounding RSA and rDSA signatures and keys may be of special interest.)

This document also attempts to incorporate clarifications gleaned from the ongoing CMP interoperability testing. Interim results of this testing has been published as an Internet Draft. "CMP Interoperability Testing: Results and Agreements", <draft-moskowitz-cmpinterop-00.txt> provides a summary of these results. The most notable modification resulting from the workshop is probably the substitution of SHA-1 HMAC for the DES-MAC. As this testing is a work in progress, there may be additional changes in the future.

### Minimum Interoperability Specification for PKI Components

#### 1. Introduction

#### 1.1 Purpose

The Minimum Interoperability Specification for PKI Components (MISPC) provides a basis for interoperation between public key infrastructure (PKI) components from different vendors. This specification will be available to companies interested in offering interoperable PKI components, to Federal agencies developing procurement specifications, and to other interested parties. It will be the basis for a NIST reference implementation. A test suite for conformance to this specification is also planned.

#### 1.2 Scope

This specification supports interoperability for a large scale PKI that issues, revokes and manages digital signature and key management public key certificates. Digital signature certificates support the use of those signatures to replace handwritten signatures in government services, commerce, and legal proceedings, and to allow distant parties, who have no previous relationship, to reliably authenticate each other and conduct business. Key management certificates support the use of key transport and key agreement algorithms to establish or protect symmetric keys for confidentiality in session or store and forward applications. Such a PKI, and the certificates it requires, may be excessive for some applications, and other more streamlined certificates and protocols may be more appropriate for more specialized and restricted applications.

The MISPC addresses:

- public key certificate generation, renewal, and revocation;
- signature generation and verification;
- certificate and certification path validation.

The specification consists primarily of a profile of certificate and CRL extensions and a set of transactions. The transactions include: certification requests, certificate renewal, certificate revocation, and retrieval of certificates and CRLs from repositories.

The MISPC focuses primarily on the aspects of PKI interoperation most apparent to end users, that is how to request and be issued a certificate, how to sign documents, how to retrieve the certificates of others, and how to validate signatures. Some aspects of the "internal" operation of a PKI, as outlined below, have not reached sufficient stability at this point, and are therefore not specified.

In this specification a PKI is broken into five components:

- *Certification Authorities (CAs)* that issue and revoke certificates;
- *Registration Authorities (RAs)* that vouch for the binding between public keys and certificate holder identities and other information;
- *Certificate holders* that are issued certificates and can sign digital documents;
- *Clients* that validate digital signatures and perform key management protocols, and

validate the corresponding certification paths from a known public key of a trusted CA;

• *Repositories* that store and make available certificates and Certificate Revocation Lists (CRLs).

Many entities will include certificate holder and client functionality. CAs and RAs will include both certificate holder and client functionality. End-entity certificate holders will generally also have client functionality. There may be some clients, however, that are not also certificate holders.

Repositories are not necessarily certificate holders and may not include client functionality. This interoperability specification addresses only one aspect of repositories, the protocol used by clients to request certificates and CRLs from the repository. This is because the precise concept, role and business model of repositories is unsettled. The X.509 certificate standard [ISO94-8] itself assumes the existence of an X.500 directory, to satisfy repository requirements, however X.500 directories, while available for some time, have not been, and do not appear to be going to be widely used.

The MISPC specifies the Lightweight Directory Access Protocol (LDAP) version 2 as the vehicle for client access of repositories, primarily because it appears to be the most generally accepted and broadly implemented alternative. This choice does not address, for example, standardized protocols for CAs to use to update repositories, nor does it address protocols for repositories to automatically shadow one another, both of which may be desirable. The former can be addressed on a case by case basis between CAs and their repositories, and the latter may not be necessary.

In the conventional approach to certificate status confirmation (which the MISPC follows), repositories are not trusted entities, rather it is the CA's signature on a CRL that validates the revocation status of certificates. On-line mechanisms for real-time certificate status confirmation would require that repositories themselves be trusted entities and that they authenticate themselves to clients. Protocols for such certificate status confirmation are not yet widely deployed, and their utility seems to depend upon the application. Therefore such protocols are outside the scope of this specification, but, since real-time certificate status confirmation may be needed for some applications, this subject may be addressed in a later revision.

The MISPC does not include a protocol for repositories to authenticate users, which would be needed to implement access by access billing for repository use. Although that may become an important business model for repositories, there does not currently appear to be enough agreement on such a business model and the supporting protocol to make this subject ripe for inclusion in a minimum interoperability specification. This subject may also be addressed in a later revision.

In some cases, out-of-band exchanges must be performed as part of the transactions defined by this specification. The format and contents of such out-of-band transactions are generally outside of the scope of this specification.<sup>1</sup>

#### 1.3 Approach

The MISPC is based on X.509 version 3 certificates and version 2 CRLs. To the extent possible,

<sup>&</sup>lt;sup>1</sup> The format and content of the electronic data provided to an RA when requesting a certificate "in person" is the exception to this rule. See section 3.5.1, *RA-Generated Registration Requests* 

this document adopts data formats and transaction sets defined in existing and evolving standards such as ITU-T X.509 [ISO94-8], ANSI [X9.31], [X9.42], [X9.44], [X9.55], [X9.57], [X9.62], and [X9.63] and the IETF's PKIX documents [RFC2459], [RFC2510], [RFC2511]. In drafting this document, whenever the stability of an evolving standard used in this document has come into question, NIST has made an educated guess regarding the direction to be followed. These issues were reviewed by industry collaborators prior to the release of this specification and represented vigorously within the appropriate standards groups to minimize departure from the stable version of the standards.

#### 1.4 Assumptions

The MISPC assumes that CAs, RAs, and certificate holders are physically separated. Where these entities are physically collocated, support for specified interfaces is not required. In particular, a PKI component that includes both RA and CA functionality is not required to support the MISPC message formats for transactions between these components. However, if that system includes a CA that supports remote RAs in addition to the local RA function, it must support the MISPC transactions for the remote RAs.

The MISPC considers CAs and RAs as functional entities in a PKI. The internal design of these entities is outside the scope of this specification.

The MISPC assumes that, at a minimum, certificate holders will have a signature key and certificate. Certificate holders may optionally maintain a key management key and certificate as well. Certificate holders that wish to request or revoke a key management certificate will use their signature key for authentication to the CA.

The MISPC does not directly support systems that do not require non-repudiation and do not maintain a signature key pair. However, the majority of such entities are computing systems (e.g., routers or link encryptors) and are maintained by administrators. If that administrator has a signature key pair for the purpose of system administration, the MISPC transaction set could be used to support certificate requests and revocation for such an entity.

The MISPC identifies three important digital signature algorithms for which suitable approved or mature draft standards exist. New algorithms could easily be incorporated as they are introduced in standards.

The MISPC also identifies three important key management algorithms for which suitable approved or mature draft standards exist. As with signature algorithms, new algorithms could easily be incorporated as they are introduced in standards.

The MISPC supports both hierarchical and networked trust models [CONOPS]. In hierarchical models, trust is delegated by a CA when it certifies a subordinate CA. Trust delegation starts at a root CA that is trusted by every node in the infrastructure. In network models, trust is established between any two CAs in peer relationships (cross-certification), thus allowing the possibility of multiple trust paths between any two CAs. The MISPC assumes that X.509 v3 extensions, such as **basicConstraints**, **nameConstraints**, **keyUsage**, and **certificatePolicies**, will be included in certificates to explicitly manage trust relationships.

The MISPC assumes that certificates and certificate revocation lists (CRLs) will be available in a repository for retrieval without authentication. MISPC clients will perform path validation by

obtaining the necessary certificates and CRLs from the appropriate repositories. The repository may be an X.500 directory or some other type accessible by using Universal Resource Identifier (URI) notation. Repositories are expected to support the Lightweight Directory Access Protocol (LDAP) [RFC 1777], therefore compliant products are required to support this protocol.

These repositories need not be linked together and other protocols may be used to retrieve certificates and CRLs. The specification requires explicit identification of the certificate repositories used and retrieval mechanisms for the issuer's certificate(s) and CRLs within the certificate.<sup>2</sup>

Certificate Revocation Lists (CRLs) are expected to be a widely implemented mechanism for revoking and validating the status of unexpired certificates. While the use of CRLs for this purpose may not be universal, and some CAs may choose to provide an on-line mechanism for validating certificate status in real time, CRL generation will be necessary for interoperability with users of other CAs. In addition to current checks of certificate validity, CRLs provide an important mechanism for documenting the historical revocation status of certificates. That is, a dated signature may be presumed to be valid if the signature date were within the validity period of the certificate, and the current CRL of the issuing CA at that date did not show the certificate to be revoked.<sup>3</sup>

Therefore, the MISPC assumes that CA products will be able to generate CRLs, and that clients will be able to use CRLs when validating certificates.

#### 1.5 Definitions, Terms, and Acronyms

Abstract Syntax Notation 1 (ASN.1): an abstract notation for structuring complex data objects.

*accredit:* recognize an entity or person to perform a specific action; CAs accredit RAs to act as their intermediary (see registration authority below).

CA certificate: a certificate whose certificate holder is trusted to issue certificates

*certificate (or public key certificate):* A digitally signed data structure defined in the X.509 standard [ISO94-8] that binds the identity of a certificate holder (or subject) to a public key.

certificate holder: An entity that is named as the subject of a valid certificate.

*certificate policy:* A named set of rules that indicates the applicability of a certificate to a particular community and/or class of application with common security requirements. For example, a particular certificate policy might indicate applicability of a type of certificate to the authentication of electronic data interchange transactions for the trading of goods within a given price range.

*certificate user*: An entity that uses certificates to know, with certainty, the public key of another entity.

certificate-using system: An implementation of those functions defined in the X.509 standard

 $<sup>^{2}</sup>$  As a consequence of this assumption, the distinguished name of the subject is not sufficient to retrieve a certificate. MISPC clients must obtain the signer's certificate, or distinguished name of the subject and the identity of the repository, from the signer.

<sup>&</sup>lt;sup>3</sup> This assumes you can accept the date attached to the signature on the basis of a trusted archive or notarization, which are outside the scope of this specification.

[ISO94-8] that are used by a certificate user. This term is defined in the Draft Amendments to X.509 [DAM] and equivalent to the term "client" used in this interoperability specification.

*Certification Authority (CA):* A trusted entity that issues certificates to end entities and other CAs. CAs issue CRLs periodically, and post certificates and CRLs to a repository.

*certification path:* An ordered sequence of certificates, leading from a certificate whose public key is known by a client, to a certificate whose public key is to be validated by the client.

*Certification Practice Statement:* A statement of the practices which a Certification Authority employs in issuing certificates.

*CRL distribution point:* A directory entry or other distribution source for CRLs; a CRL distributed through a CRL distribution point may contain revocation entries for only a subset of the full set of certificates issued by one CA or may contain revocation entries for multiple CAs.

certificate revocation list (CRL): a list of revoked but unexpired certificates issued by a CA.

certify: the act of issuing a certificate.

*client (or PKI client):* A function that uses the PKI to obtain certificates and validate certificates and signatures. Client functions are present in CAs and end entities. Client functions may also be present in entities that are not certificate holders. That is, a system or user that verifies signatures and validation paths is a client, even if it does not hold a certificate itself. See section 2.4.

*cross certificate*: a CA certificate that describes a non-hierarchical trust relationship between two CAs

cross certification: the process of establishing non-hierarchical trust relationships between CAs

*delta-CRL*: A partial CRL indicating only changes since a prior CRL issue.

DES: The symmetric encryption algorithm defined by the Data Encryption Standard (FIPS 46-2).

*DES MAC:* An algorithm for generating a message authentication code (mac) using the symmetric encryption algorithm DES.

*Distinguished Encoding Rules (DER):* rules for encoding ASN.1 objects which give a consistent encoding for each ASN.1 value. Implementations conforming to this specification shall encode ASN.1 objects using the DER.

*digital signature:* a data unit that allows a recipient of a message to verify the identity of the signatory and integrity of the message.

Digital Signature Algorithm (DSA): a digital signature algorithm specified in FIPS PUB 186-1.

*directory service (DS):* a distributed database service capable of storing information, such as certificates and CRLs, in various nodes or servers distributed across a network.

*end entity*: A certificate subject which uses its private key for purposes other than signing certificates.

*Elliptic Curve Digital Signature Algorithm (ECDSA):* a digital signature algorithm that is an analog of DSA using elliptic curve mathematics and specified in ANSI draft standard X9.62 [X9.62].

*hash:* a function which maps strings of bits to fixed-length strings of bits, satisfying the following two properties: it is computationally infeasible to find for a given output an input which maps to this output; and it is computationally infeasible to find for a given input a second input which maps to the same output.

hash code: The string of bits which is the output of a hash function

*LDAP:* The Lightweight Directory Access Protocol, or LDAP, is a directory access protocol. In this document, LDAP refers to the protocol defined by RFC 1777, which is also known as LDAP V2. LDAP V2 describes unauthenticated retrieval mechanisms.

*message authentication code:* a data authenticator generated from the message, usually through cryptographic techniques. In general, a cryptographic key is also required as an input.

message digest: the fixed size result of hashing a message.

*out of band:* Some transactions between PKI components will be performed through physical procedures rather than implemented electronically. Such transactions are described as *out of band* transactions.

*policy mapping*: Recognizing that, when a CA in one domain certifies a CA in another domain, a particular certificate policy in the second domain may be considered by the authority of the first domain to be equivalent (but not necessarily identical in all respects) to a particular certificate policy in the first domain.

*Registration Authority (RA)*: an entity that acts an intermediary between the CA and a prospective certificate subject; the CA trusts the RA to verify the subject's identity and that the subject possesses the private key corresponding to the public key to be bound to that identity in a certificate. Note that equivalent functions are referred to as Local Registration Authority (LRAs) or Organizational Registration Authorities (ORAs) in some documents, including the MISPC Version 1.

*repository*: a database service capable of storing information, such as certificates and CRLs, allowing unauthenticated information retrieval. Repositories include, but are not limited to, directory services.

*RSA:* For the purposes of this specification, RSA is a public-key signature algorithm specified by PKCS #1 [PKCS#1]. As a reversible public-key algorithm, it may also be used for encryption.

self-issued certificate: a CA-certificate whose subject and issuer are identical.

*URI:* A uniform resource identifier, or URI, is a short string containing a name or address which refers to an object in the "web."

*URL*: A uniform resource locator, or URL, is a short string containing an address which refers to an object in the "web." URLs are a subset of URIs.

*Well Known X.500 Directory*: In some environments, an X.500 service may be widely available and used throughout an organization. If such a directory service is used to distribute certificates and CRLs issued by that organization, such information need not be included in the certificate.

#### 2. Infrastructure Component Specifications

This section specifies a minimal set of functions and transactions required for the interoperation of PKI components. It includes specifications for CAs, RAs, certificate holders, and PKI Clients.

#### 2.1 Certification Authority (CA)

CAs generate, revoke, publish, and archive certificates. They rely upon a repository to make certificates and CRLs available to all certificate users.

CAs generate their own key pairs and publish their own certificates. As appropriate, CA should be able to generate, and assess the quality of, any parameters required to generate/verify their signatures. To enable CAs to join existing hierarchically managed infrastructures, they shall be able to request certificates from a parent CA. CAs shall also be able to generate cross certificates, to support cross-certification with other CAs as allowed by their operational policies. CAs archive all transactions, including service requests and responses from and to other PKI components.

CAs accredit RAs, which vouch for the identity and other attributes of users requesting certificates. This accreditation is an off-line decision to accept RA-generated certification requests from that RA. CAs identify certificate holders using X.500 distinguished names. Distinguished names uniquely identify certificate holders.

CAs themselves include both a certificate holder function to request, revoke and renew certificates issued by other CAs (see sec. 2.3) and a client function to retrieve certificates and CRLs, and validate certification paths (see sec. 2.4).

#### 2.1.1 Interoperability-Relevant CA Functional Specifications

CAs perform the following functions:

- Issue and deliver certificates to end-entities and CAs;
- Accept revocation requests from certificate holders and RAs for certificates it issued;
- Post certificates and CRLs to the repository; and
- Request CA certificates.

#### **Issuing Digital Signature Certificates**

CAs support three types of certification requests for digital signature certificates: *self-registration*, *RA-generated registration*, and *renewal*.<sup>4</sup> CAs authenticate the identity of the certificate's subject differently for each type of request. The prospective certificate holder supplies an authenticator in a *self-registration request*; the authenticator is derived from a secret obtained from an RA. RAs generate and sign *RA-generated registration requests*, vouching for the identity of the subject, when the subject physically attends the RA. The subjects of currently valid certificates can vouch for their own identity in a *renewal request* by signing with their current private key.

<sup>&</sup>lt;sup>4</sup> CAs may be configurable to reject one or more classes of certification requests if the certificate policy prohibits such transactions.

In an RA-generated registration request, the RA vouches for the prospective certificate holder's identity and the binding to the public key. When CAs receive certification requests from accredited RAs, they shall process the requests and, if accepted, generate new certificates, post the certificates to a repository<sup>5</sup>, and send them to the requesting RAs. CAs may also send the new certificate to the certificate holders. CAs shall reject RA-generated certification requests that do not come from a recognized RA, that have invalid signatures, or that contain unmatched information. If a CA rejects an RA-generated certification request, it shall report the failure to the RA stating the reason.

In a self-registration request, the RA provides a secret message to the prospective certificate holder. The entity generates its own key pair, forms a certification request, signs it with the corresponding private key material, and includes authentication information based on the secret provided by the RA.<sup>6</sup> The CA receives the request, verifies the requester's identity through the authentication information and verifies that the entity holds the corresponding private key material. If accepted, the CA will generate a new certificate, post the certificate to the repository, and send it to the certificate holder. The CA may reject self-registration requests if the authentication information does not verify, the signature is invalid, or fields contain unmatched information. If a CA rejects a self-registration request, it shall report the failure to the requester stating the reason.

In a renewal request, the established identity of the requester is perpetuated with the request. Certificate renewals are initiated by the certificate holder and sent directly to the CA. CAs process the renewals and, if correct, send the new certificates to the certificate holders and post the new certificates to the repository. CAs may reject certificate renewal requests with invalid signatures, requests from entities not currently certified, and renewal requests that are not allowed by the CA's certification practice statement or the certificate policy. If a CA rejects a certificate renewal request, it shall report the failure to the requesting entity stating the reason.

#### Issuing Key Management Certificates

CAs may support certification requests for key management certificates for a requester that possesses a valid signature certificate issued by that CA. The requester's identity is verified by digitally signing the requests. Requesters may generate their own key pairs, or the CA may generate the keys for the requesters. Both transactions are extended versions of the renewal request.<sup>7</sup>

In the Locally-Generated Key Management Certificate Request, the certificate holder generates a key pair for a key transport or key agreement algorithm. The certificate holder generates a certificate request, includes the public key and signs the request with a signature key currently certified by this CA. The CA may, depending upon its policy, initiate a challenge-response protocol to verify the requester's possession of the private key. The transaction is completed by a certification response from the CA containing either the certificate or an error code.

<sup>&</sup>lt;sup>5</sup> Conforming CAs shall be able to post the certificates they issue to a repository. However, it is not necessary to post end-entity certificates, since the certificate holder may provide the certificate with the signed document.

<sup>&</sup>lt;sup>6</sup> Where the CA and RA are not co-located, this also requires an exchange of secrets between the CA and RA. Details of this exchange are outside the scope of this specification.

<sup>&</sup>lt;sup>7</sup> Neither type of transaction requires or prevents implementation of various key escrow schemes.

In a Centrally-Generated Key Management Certificate Request, the certificate holder generates a certificate request, specifies the desired key management algorithm and signs the request with a signature key currently certified by this CA. The CA responds with a certificate and an encrypted private key.

CAs that claim conformance to the MISPC confidentiality enhancements must implement the Locally-Generated Key Management Certificate Request. Implementation of the Centrally-Generated Key Management Certificate Request is optional.

#### Cross Certification

CAs may issue certificates to other CAs with appropriate constraints. The decision to crosscertify is made out-of-band and involves examination of Certification Practice Statements and certificate policies. Each CA determines the appropriate constraints for path validation by their users. After obtaining the other CA's public key, the CA generates the certificate and posts it to the repository.

Optionally, cross-certifying CAs may exchange certificates, construct certificate-pairs, and post them to the repository.<sup>8</sup>

#### **Revoking Certificates**

CAs shall be capable of generating and issuing certificate revocation lists (CRLs). CAs shall be able to issue CRLs that contain all revoked certificates that they issued and have not expired. Optionally, CAs may also issue indirect and delta CRLs. The types of CRLs issued will be determined by the CA's certification practice statement.

In those cases where a CA issues a single CRL for all revoked certificates it has issued:

- When a new CRL is generated, all revoked unexpired certificates from the previous CRL shall be carried over to the new CRL, and any certificates with approved pending certificate revocation requests shall be added to the new CRL. Certificates on the previous CRL with a reason code of **certificateHold** may be carried over to the new CRL with same reason code, listed on the new CRL with a different and permanent reason code, or omitted from the new CRL. Omission from the new CRL indicates the CA will vouch for the binding between the subject and public key. A certificate with an approved pending certificate revocation request shall be included in the next CRL even if it expires before the CRL is issued.
- In this case, CAs shall only revoke certificates they issued.<sup>9</sup> The signer of the revocation request must either be the certificate holder or an authorized entity (such as an accredited RA) acting on behalf of the certificate holder or the certificate holder's organization. CAs shall validate revocation requests prior to including a certificate in a CRL. Validation of a revocation request shall include validation of the signature on the request. Out-of-band validation of revocation requests signed by RAs may optionally be required

<sup>&</sup>lt;sup>8</sup> A CA may issue a certificate to another CA even if the latter refuses to issue a certificate to the former. In this case, the CA could (optionally) construct a cross certificate pair containing only the **reverseCertificate**.

<sup>&</sup>lt;sup>9</sup> Revocation may be initiated by receipt of a signed request, or by the CA's own procedures. This specification does not address revocations initiated by the CA.

#### by the certificate policy.

CAs shall issue X.509 version 2 CRLs.<sup>10</sup> The fields and extensions utilized, and the values assigned to them, shall be in accordance with section 3.2.1. After generating and signing a CRL, CAs shall send it to the repository.

#### Post Certificates, Cross Certificates, and CRLs

CAs shall be capable of posting certificates, cross certificate pairs, and CRLs for retrieval by PKI clients. CAs shall always post CA certificates, cross certificate pairs, CRLs, and end entity key management certificates. Posting of end-entity digital signature certificates is optional. The mechanisms used to update directories is beyond the scope of this specification.

#### Request CA Certificates

CAs shall be capable of requesting certificates from hierarchically superior CAs to support PKIs based on the hierarchical trust model. This request is supported as described in section 3.5.1. The certificate request shall identify the entity as a CA through the **basicConstraints** extension as described in section 3.1.3.3.

#### 2.1.2 Electronic Transaction Set.

Table 2-1 summarizes electronic transactions used in providing certificate management services. These transactions enable:

- processing of certification requests and certificate revocation requests for end entity certificates;
- posting of certificates and CRLs on the repository;
- the retrieval of certificates and CRLs from the repository for signature validation.

CAs shall process RA-generated certification requests in the form of **CertReq** messages.<sup>11</sup> **CertReq** messages are signed by the RA in the **PKIProtection** structure. By signing requests, RAs vouch for the identity of the certificate holder and confirm that requesting certificate holders are in possession of the corresponding private keys. CAs respond to the RAs or certificate holders with **CertRep** messages. If a request was accepted, the **CertRep** message contains the new certificate. If the request was rejected, the message contains the error code (see sec. 3.5.1).

CAs shall also support the self-registration request, where users who are not current certificate holders sign their own certificate request. The CA shall require the entity to generate authentication information based on out-of-band interaction with an RA. This information substitutes for RA signature to vouch for the requester's identity. To request a certificate without appearing before an RA, the entity obtains some information out-of-band from the RA. This information might be a symmetric key for use in generation of a mac or keyed hash. The entity generates a **CertReq** message and signs it with the entity's new private key. This message is then protected with the information obtained out-of-band as directed by the RA. The CA generates a

<sup>&</sup>lt;sup>10</sup> Version 2 CRLs correspond to the Version 3 certificate; the Version 2 certificate definition did not result in creation of a new CRL format.

<sup>&</sup>lt;sup>11</sup> This section refers to **CertReq**, **CertRep**, **RevReq** and **RevRep** messages. The precise structure and content of these messages is defined in section 3.4.

**CertRep** message; if the request was fulfilled the message contains the new certificate. If the request was rejected, the message contains error codes. This transaction is described in detail in section 3.5.3.<sup>12</sup>

CAs shall process certificate renewal requests in the form of **CertReq** messages. These messages are sent to a CA by the entity requesting the certificate. The message shall include the certificate holder's distinguished name, the serial number of their current certificate, the new public key, and proof of possession of the private key. The message may optionally include a proposed validity period and a proposed key id. The message shall be signed with the private key corresponding to the certificate holder's unexpired, unrevoked certificate. CAs shall respond to the requester in the form of an **CertRep** message. This message shall contain either a new certificate or a failure code. If issued, the certificate shall include the certificate holder's distinguished name and the new public key. CAs are free to modify the validity period proposed in the request. CAs shall generate a key identifier if the message did not include one. (see sec. 3.5.2.)

CAs shall receive **RevReq** messages from RAs or certificate holders. The **RevReq** message shall include the certificate serial number or the certificate holder's distinguished name. CAs shall respond with a **RevRep** message. This message shall include status and failure information, and may include additional details about the revoked certificate(s).

CAs may support one or both of the following transactions:

- key management certificate requests for end-entity generated key pairs; or
- key management certificate requests for centrally generated key pairs.

CAs that support key management certificate requests for end-entity generated key pairs shall process these requests in the form of **CertReq** messages. These messages are sent to a CA by the entity requesting the certificate. The message shall include the certificate holder's distinguished name, the serial number of their current certificate, and the new public key. The message may optionally include a proposed validity period and a proposed key id. The message shall be signed with the private key corresponding to the certificate holder's unexpired, unrevoked signature certificate. If a CA desires proof of possession of the corresponding private key, the CA initiates a challenge-response mechanisms by encrypting a random challenge with the new private key.<sup>13</sup> The requester returns the random challenge, and signs the message as above. If the CA is not concerned with proof of possession, the challenge-response mechanism may be omitted. At this point, the CA shall respond to the requester in the form of an **CertRep** message. This message shall contain either a new certificate or a failure code. If issued, the certificate shall include the certificate holder's distinguished name and the new public key. CAs are free to modify the validity period proposed in the request. CAs shall generate a key identifier if the message did not include one.

<sup>&</sup>lt;sup>12</sup> An alternative syntax for this transaction is specified in section 3.5.5.

<sup>&</sup>lt;sup>13</sup> A CA may also wish to perform public key validation of the public key value. This is a policy decision, and does not affect the certificate request protocol. Mechanisms for public key validation are algorithm dependent.

Transaction	Description	From	То
RA-Generated Registration	RA submits a certificate request on behalf of an authenticated entity	RA	CA
Request (sec. 3.5.1)	CA returns signed certificate or error message	CA	RA
Certificate Revocation	RA or certificate holder requests revocation of a certificate	requester	Issuer CA
(sec. 3.5.6)	CA responds with acceptance or rejection of the request	Issuer CA	requester
Self-Registration Request	message signed with new public key encapsulates certificate request with ORA-directed protection value	client	Issuer CA
(secs. 3.5.3 and 3.5.5)	CA returns signed certificate and CA's certificate or an error message	Issuer CA	client
Certificate Renewal Request	certificate request containing new public key with proof of possession and current certificate serial number; signed with current private key	Certificate holder	CA
(sec. 3.5.2)	CA returns signed certificate or error message	CA	certificate holder
End-Entity Generated Key	End-entity generates a key management certificate request including the public key and signs it with its currently certified signature key.	Certificate holder	CA
Management Certificate	The CA challenges the requester to prove possession of private key	CA	Certificate holder
Request (sec. 3.5.7)	If challenged, the requester generates proof of possession as response to challenge message	Certificate holder	CA
	CA returns signed certificate or error message	certificate holder	CA
Centrally Generated Key Management	End-entity generates a key management certificate request specifying the key management algorithm and a NULL public key and signs it with its currently certified signature key.	Certificate holder	CA
Certificate Request (sec. 3.5.10)	CA generates a key pair and returns signed certificate and encrypted private key or an error message	CA	Certificate holder

Table 2-1

CAs that support key management certificate requests for centrally generated key pairs shall process these requests in the form of **CertReq** messages. These messages are sent to a CA by the entity requesting the certificate. The message shall include the certificate holder's distinguished name and the serial number of their current certificate. The message may optionally include a proposed validity period. The message shall be signed with the private key corresponding to the certificate holder's unexpired, unrevoked certificate. CAs shall respond to the requester in the form of an **CertRep** message. This message shall contain either a new certificate and encrypted private key, or a failure code. If issued, the certificate shall include the certificate holder's distinguished name and the new public key. CAs are free to modify the validity period proposed in the request.

CAs shall post CA certificates, cross certificate pairs, CRLs, and end-entity key management certificates that it issues to a repository. CAs may optionally be capable of posting end entity certificates to a repository.<sup>14</sup>

#### 2.2 Registration Authority (RA)

RAs vouch for the identity of entities requesting certification. RAs may verify that identity by requiring the requesting entity to attend the RA physically with a physical token, or through outof-band mechanisms. Where the entity physically attends the RA, the RA also verifies their possession of private key material corresponding to the public key by verifying a signed message (as described in sec. 3.5.1).

The format for a certificate request on behalf of an entity in physical attendance appears in section 3.5.1. RAs shall verify the entity possesses a complete key pair. After the key pair and the entity's identity are verified, an RA signs and sends an electronic certificate request to the appropriate CA.

Certificate requests on behalf of a user who does not physically attend the RA require that the RA provide authentication information to the entity. This information is used by the entity to authenticate itself to the CA in a self-registration request as defined in section 3.5.3. This specification does not define the content or format of the out-of-band exchange(s) required to implement self-registration requests.

RAs may request certificate revocation for end-entity certificates issued by CAs that have accredited them. The format of the **RevReq** is presented in section 3.5.6. The RA function may be collocated with the CA or performed at a separate facility.

RAs themselves include both a certificate holder function to request, revoke and renew certificates (where it is the subject) issued by CAs (see sec. 2.3) and a client function to retrieve certificates and CRLs and validate certification paths (see sec. 2.4).

<sup>&</sup>lt;sup>14</sup> Posting of end entity digital signature certificates is not strictly required, since the originator of a signature can supply their own certificate.

#### 2.2.1 Interoperability-Relevant RA Functional Specifications

RAs shall perform the following functions:

- Accept and validate certification requests;
- Send certification requests to the CA;
  - Retrieve certificates and CRLs from the repository; and
  - Generate certificate revocation requests.

The RA shall be able to pass the newly signed certificate on to the certificate holder, along with the CA's certificate.

RAs shall generate and sign certificate revocation requests on behalf of certificate holders who no longer possess their private key and suspect compromise.<sup>15</sup> If permitted by the CA's certification practice statement, RAs shall also generate and sign certificate revocation requests on behalf of the certificate holder's organization. Revocation requests are signed by the RA which then sends them to the issuing CA.

#### 2.2.2 Transaction Set

Table 2-3 gives the subset of electronic transactions used by RAs. These transactions enable request, delivery, and revocation of end entity certificates, and the retrieval of certificates and CRLs from the repository for signature validation. The following text provides an overview of these transactions; they are described more fully in section 3.5.

RAs receive certification requests from prospective certificate holders in the form of **CertReq** messages. The **CertReq** message is signed by the prospective certificate holder in the **PKIProtection** structure. After reviewing the requester's credentials and confirming that the prospective certificate holder is in possession of the corresponding private key, the RA extracts the public key information, and creates a new **CertReq** message with the RA's name and signature. The RA sends this message to a CA. RAs shall provide certificate holders with the CA's certificate.

RAs may receive **CertRep** messages from the CA. If a certification request is rejected, the RA will review the error code from the CA and may submit a new request. If a certification request is accepted, the RA may provide the new certificate to the certificate holder.

<sup>&</sup>lt;sup>15</sup> Signature keys lost but not believed compromised are not necessarily revoked; this is determined by policy. Note that confidentiality keys which are lost must be revoked regardless, or a sending party may encrypt and transmit messages the receiver could never decrypt.

Transaction	Description	From	То
RA-Generated Registration Request	User (or system administrator) submits digitally signed certificate request to RA with proof of identity	client	RA
(sec. 3.5.1)	RA submits a certificate request on behalf of an authenticated prospective certificate holder	RA	CA
	CA returns signed certificate or error message	CA	RA
Certificate	RA requests revocation of a certificate	RA	Issuer CA
Revocation (sec. 3.5.6)	CA responds with acceptance or rejection of revocation request	Issuer CA	RA

#### Table 2-1 RA Electronic Transaction Set

RAs shall generate revocation requests upon request of certificate holders who no longer possess their private key or the certificate holder's organization. By signing the request, the RA is vouching for the identity of the requester. RAs shall generate **RevReq** messages, including the certificate serial number or the certificate holder's distinguished name. The **RevReq** message shall be signed by an RA. The CA shall respond to the RA with a **RevReq** message.

This message shall include status and failure information, and may include additional details about the revoked certificate. If the certificate is revoked, the RA shall provide this information to the requester. If the request is rejected, the RA will review the error code and may reformulate the request.

#### 2.3 Certificate Holder Specifications

The PKI provides certificate management functions for certificate holders. Certificate holders include CAs, RAs and other end entities. End entities may include persons and computing systems (e.g., routers and firewalls) or applications (in addition to CAs and RAs).

PKI certificate holders generate signatures and support PKI transactions to obtain, revoke and renew their certificates.

#### 2.3.1 Interoperability-Relevant PKI Certificate Holders Functional Specifications

Certificate holders shall be able to:

- generate signatures;
- generate certificate requests;
- request certificate revocation;
- request certificate renewal (optional).

Certificate holders are also PKI clients, and must also meet the specifications defined in section 2.4.

#### 2.3.2 Certificate Holders Transaction Set

Table 2-3 gives the summary of transactions used by certificate holders. These transactions enable certificate holders to request new certificates and request revocation of certificates held by the certificate holder (if any) for whom the client acts. All certificate holder transactions are performed with the CA that issued the certificate or an RA accredited by that CA.

Certificate holders shall be able to request revocation of their own certificates. This transaction is performed with the CA and permits certificate holders to sign their own certificate revocation requests. Certificate holders generate a **RevReq** message for each certificate they wish to revoke and transmit to the issuing CA. The **RevReq** message shall include the reason for revocation. The CA generates a **RevRep** message for each request and transmits it to the certificate holder. This transaction is described in detail in section 3.5.6.

Certificate holders shall be able to generate a **CertReq** message to present to an RA for in person authenticated certificate requests. The certificate holder constructs and signs the **CertReq** message, so the RA can verify the requester holds corresponding private key material.

Certificate holders may also implement the Certificate Renewal Request. This transaction is performed with the CA and permits certificate holders to sign their own certificate requests (i.e., without an RA verification of identity). CAs shall support this transaction, but its use is determined by the certificate policy. To request a new certificate without appearing before an RA, the certificate holder generates a **CertReq** message and signs it with both the new and current private keys. The CA generates a **CertRep** message; if the request was fulfilled the message contains the new certificate. If the request was rejected, the message contains error codes. This transaction is described in detail in section 3.5.2.

Certificate holders may also implement the self-registration request to request a certificate when they are not current certificate holders. This transaction is performed with the CA and permits certificate holders to sign their own certificate requests. The CA shall require the entity to generate or include information based on out-of-band interaction with an RA. This information substitutes for RA verification of identity. CAs shall support this transaction, but its use is determined by the certificate policy. To request a certificate without appearing before an RA, the entity obtains some information out-of-band from the RA. This information might be a secret key for use in mac generation or a signed message that will simply be included in the request. The entity generates a **CertReq** message and signs it with the entity's new private key. The entity attaches appropriate protection information to the signed message as directed by the RA. The CA generates a **CertRep** message; if the request was fulfilled the message contains the new certificate. If the request was rejected, the message contains error codes. This transaction is described in detail in section 3.5.3.<sup>16</sup>

Certificate holders may also implement the Locally-Generated Key Management Certificate Request. In this transaction, the certificate holder generates a key pair for a key transport or key agreement algorithm. The certificate holder generates a certificate request, includes the public key and signs the request with a signature key currently certified by this CA. The CA may, depending upon its policy, initiate a challenge-response protocol to verify the requester's possession of the private key. The transaction is completed by a certification response from the

<sup>&</sup>lt;sup>16</sup> An alternative syntax for this transaction is presented in section 3.5.5.

CA containing either the certificate or an error code.

Certificate holders may also implement the Centrally-Generated Key Management Certificate Request. In this transaction, the certificate holder generates a certificate request, specifies the desired key management algorithm and signs the request with a signature key currently certified by this CA. The CA responds with a certificate and an encrypted private key or an error code.

Transaction	Description	From	То
RA-Generated Registration (see sec. 3.5.1)	User (or system administrator) submits digitally signed certificate request to RA with proof of identity	client	RA
Certificate Revocation	certificate holder requests revocation of a certificate	certificate holder	Issuer CA
(sec. 3.5.6)	CA responds with acceptance or rejection of revocation request	Issuer CA	certificate holder
Self- Registration Request	message signed with new public key encapsulates certificate request with RA-directed protection value	client	Issuer CA
(secs. 3.5.3 and 3.5.5)	CA returns signed certificate and CA's certificate or an error message	Issuer CA	client
Certificate Renewal Request	certificate request containing new public key with proof of possession and current certificate serial number; signed with current private key	certificate holder	Issuer CA
(sec. 3.5.2)	CA returns signed certificate and CA's certificate or an error message	Issuer CA	certificate holder
End-Entity Generated Key	End-entity generates a key management certificate request including the public key and signs it with its currently certified signature key.	Certificate holder	CA
Management Certificate	The CA challenges the requester to prove possession of private key	CA	Certificate holder
Request (sec. 3.5.7)	If challenged, the requester generates proof of possession as response to challenge message	Certificate holder	CA
	CA returns signed certificate or error message	CA	certificate holder
Centrally Generated Key Management	End-entity generates a key management certificate request specifying the key management algorithm and a NULL public key and signs it with its currently certified signature key.	Certificate holder	СА
Certificate Request (sec. 3.5.10)	CA generates a key pair and returns signed certificate and encrypted private key or an error message	CA	Certificate holder

#### Table 2-2 Certificate Holders Electronic Transaction Set

#### 2.4 Client Specifications

PKI Clients use the PKI to provide certificate processing functions for certificate holders and certificate users, including CAs and other end entities. End entities may also include RAs, persons and computing systems (e.g., routers and firewalls).

At a minimum, PKI Clients validate signatures, obtain certificates and CRLs, and validate certification paths. PKI Clients that serve certificate holders also generate signatures and may support PKI transactions to revoke or renew their certificates.

#### 2.4.1 Interoperability-Relevant PKI Client Functional Specifications

At a minimum, clients shall be able to:

- verify signatures;
- obtain certificates and CRLs from a repository; and
- validate certification paths.

#### 2.4.2 PKI Client Transaction Set

Table 2-4 gives the summary of transactions used by clients. These transactions enable clients to obtain certificates and CRLs from the repository. All client transactions are performed with the certificate repository. All clients shall support the following transactions:

- Retrieve certificates this transaction permits a user to bind to the directory service or a specified repository using LDAP and retrieve one or more certificate(s) according to:
  - subject name; or
  - certificate serial number and issuer's name.
- Retrieve a CRL This transaction permits a user to bind to the directory service or a specified repository using LDAP and retrieve the current CRL for a particular CA, or a specifically identified CRL.

At a minimum, retrieval of certificates and CRLs using the Lightweight Directory Access Protocol (LDAP) shall be supported by all compliant clients. These transactions are described further in [RFC1777].

Transaction	Description	From	То
Retrieve Certificate	Query repository or specified repository for an entity's certificate(s)	client	repository
(see sec. 3.5.12)	return certificate or error message to requester	repository	client
Retrieve CRL (sec. 3.5.13)	Query repository or specified repository for latest CRL issued by a particular CA	client	repository
	return CRL to requester	repository	client

Table 2-4	<b>Client Electronic</b>	Transaction	Set
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#### 3. Data Formats

Basic data formats must be defined for interoperability of PKI components. The data formats include certificate, CRL, and transaction formats. These specifications include data formats for all transactions between infrastructure components, and between PKI clients and infrastructure components.

#### 3.1 Certificate Format

The X.509 v3 certificate format shall be used. Although the revision to ITU-T Recommendation X.509 that specifies the version 3 format is not yet published, the version 3 format has been widely adopted and is specified in American National Standards Institute X9.55-1995 [X9.55], and the Internet Engineering Task Force's Internet Public Key Infrastructure working document [RFC2459]. The X.509 version 3 certificate includes the following:

Version Serial Number Issuer Signature Algorithm Issuer Distinguished Name Validity Period Subject Distinguished Name Subject Public Key Information Issuer Unique Identifier (optional) Subject Unique Identifier (optional) Extensions (optional) Issuer's Signature on all the above fields

#### 3.1.1 Certificate Fields

The Abstract Syntax Notation One (ASN.1) definition of the X.509 certificate syntax is stated in Appendix A. For signature calculation, the certificate is encoded under the ASN.1 Distinguished Encoding Rules (DER). ASN.1 DER encoding is a tag, length, value encoding system for each element.[ISO25-1]

The following items specify the use of the X.509 v3 certificate. With the exception of the optional **subjectUniquelD** and the **issuerUniquelD** fields, CAs shall generate these fields and clients shall be capable of processing them in accordance with the X.509 standard. CAs shall not issue certificates containing the optional **subjectUniquelD** and the **issuerUniquelD** fields. Clients are not required to process **subjectUniquelD** and the **issuerUniquelD** fields; however, they shall reject certificates containing these fields if they do not process them.

#### Version

The **version** field describes the version of the encoded certificate. The value of this field shall be 2, signifying a version 3 certificate.

#### Serial number

The **serialNumber** is an integer assigned by the CA to each certificate. It shall be unique for each certificate issued by a given CA (i.e., the issuer name and serial number identify a unique certificate).

#### **Signature**

The **signature** field contains the algorithm identifier for the algorithm used to sign the certificate. The **signature** field includes an **algorithmIdentifier**, which, in principle may be used to pass parameters. Certificates conforming to this interoperability specification shall be signed with either the DSA, RSA or ECDSA algorithms, and the contents of the **algorithmIdentifier** field shall be as specified in section 3.1.2.2. Certificates shall not use the **signature** field to pass parameters (see <u>Subject Public Key Information</u> below) since this field is not protected by the issuer's signature.<sup>17</sup>

#### Issuer Name

The **issuer** field provides a globally unique identifier of the authority signing the certificate. The syntax of the issuer name is an X.500 distinguished name. The distinguished name is composed of AttributeType - AttributeValue pairs. In general, the AttributeType will be defined by the X.500 series of recommendations; AttributeValue will be of type **DirectoryString**.

**DirectoryString** is a choice of **PrintableString**, **TeletexString**, **BMPString**, **UniversalString**, and **Utf8String**. **PrintableString** is a basic Latin character set supporting upper and lowercase letters, digits, and a handful of special characters. **TeletexString** is a superset of **PrintableString**, adding Latin characters with accents and Japanese characters. **BMPString** is a two-octet character set satisfying most European character sets. **UniversalString** is a multi-octet character set including all the major character sets. **Utf8String** is an alternative encoding of **UniversalString**.

When establishing new names, conforming CAs shall always use the most restrictive choice from {PrintableString, BMPString, and Utf8String}when constructing a DirectoryString. That is, an AttributeValue which requires only basic Latin characters shall always be represented as PrintableString. An AttributeValue that includes accented Latin characters shall be represented as BMPString. Utf8String shall only be used if the character set for BMPString is insufficient.

The **TeletexString** and **UniversalString** are included for backward compatibility, and should not be used for certificates for new subjects. However, these types may be used in certificates where the name was previously established.<sup>18</sup> Clients should be prepared to receive certificates with these types.

Alternative names may be supplied in the **issuerAltName** extension and some users of X.509 certificates apparently contemplate a null **issuer** field. However, certificates conforming to this interoperability specification shall contain the X.500 distinguished name of the certificate issuer in this field.

<sup>&</sup>lt;sup>17</sup> See "A Security Flaw in the X.509 Standard," available from http://www.cygnacom.com/docfiles/dsaflaw.zip, for the rationale for excluding parameters from this field.

<sup>&</sup>lt;sup>18</sup> This is especially important for cross certificates. Changing the name form will invalidate certification paths.

#### <u>Validity</u>

The validity field indicates the dates on which the certificate becomes valid (notBefore) and on which the certificate ceases to be valid (notAfter). The validity field may represent dates in UTCTime or GeneralizedTime.

For dates between the year 1950 and the year 2049 (inclusive), the validity field shall always use **UTCTime**. The **UTCTime** (Coordinated Universal Time) values included in this field shall be expressed in Greenwich Mean Time (Zulu) and shall express granularity to the second. Seconds shall be explicitly stated, even if zero. **UTCTime** shall be expressed as YYMMDDHHMMSSZ. The year field shall be interpreted as follows:

- if YY is equal to or greater than 50, the year shall be 19YY; and
- if YY is less than 50, the year shall be 20YY.

For dates in the year 2050 or later, the validity field shall always use **GeneralizedTime**. The **GeneralizedTime** values included in this field shall be expressed in Greenwich Mean Time (Zulu) and shall express granularity to the second. Seconds shall be explicitly stated, even if zero. That is, **GeneralizedTime** shall be expressed as YYYYMMDDHHMMSSZ. **GeneralizedTime** values MUST NOT include fractional seconds.

(Representation of dates before the year 1950 is not required to implement this specification.)

#### Subject Name

The purpose of the **subject** field is to provide a unique identifier of the subject of the certificate. The syntax of the subject name shall be an X.500 distinguished name. As described for issuer names, conforming CAs shall use the most restrictive choice when constructing **DirectoryStrings**. Alternative names may be supplied in the **subjectAltName** extension and some users of X.509 certificates apparently contemplate a null **subject** field. However, certificates conforming to this interoperability specification shall contain the subject's X.500 distinguished name in this field.

#### Subject Public Key Information

The **subjectPublicKeyInfo** field is used to carry the public key and identify the algorithm with which the key is used. It includes the **subjectPublicKey** field and an **algorithmIdentifier** field with **algorithm** and **parameters** subfields. Certificates conforming to this interoperability specification shall use either the DSA, rDSA, ECDSA or RSA algorithms, and the contents of the **algorithmIdentifier** field shall be as specified in section 3.1.2.2. The **parameters** subfield of the **subjectPublicKeyInfo** field shall be the only method used to pass or obtain DSA or ECDSA parameters.

#### Unique Identifiers

The **subjectUniqueIdentifier** and **issuerUniqueIdentifier** fields are present in the certificate to handle the possibility of reuse of subject and/or issuer names over time. Compliant CAs shall not issue certificates that include these unique identifiers. Compliant PKI clients are not required to process certificates that include these unique identifiers. However, if they do not process these fields, they are required to reject certificates that include these fields.

#### **Extension**

The addition of the **extension** field is the principal change introduced to X.509 v3 certificates. Extensions have three components: **extnld**, that names the extension, **critical**, the criticality flag that specifies that the extension is critical or noncritical, and **extnValue**, the extension value. A certificate may contain any number of extensions, including locally defined extensions. If the criticality flag is set, a client shall either be able to process that extension, or shall not validate the certificate.

A set of standardized extensions has been developed in an amendment to the X.509 standard [DAM]. The use of these standardized extensions in conforming implementations is specified in section 3.1.3 below.

#### Issuer's Signature

The actual signature on the certificate is defined by the use of the **SIGNED** parameterized type, which expands to a **SEQUENCE** of the data being signed (i.e., the certificate), an algorithm identifier, and a **BIT STRING** which is the actual signature. The **algorithmIdentifier** that identifies the algorithm used to sign the certificate. Although this **algorithimIdentifier** field includes a **parameters** field that can, in principle, be used to pass the parameters used by the signature algorithm (see sec. 3.1.2.2), it is not itself a signed object. The **parameters** field of the certificate signature shall not be used to pass parameters. When parameters are used to validate a signature, they shall be obtained from the **subjectPublicKeyInfo** field of the issuing CA's certificate.

#### 3.1.2 Cryptographic Algorithms

This document specifies six classes of cryptographic algorithms: secure hash algorithms; digital signature algorithms; key agreement algorithms; key transport algorithms; message authentication algorithms; and symmetric encryption algorithms. Where describing a digital signature, algorithms are always identified with a secure hash algorithm. Where describing a public key, as in a certificate, the hash algorithm is omitted. This permits a certificate to be used even if a hash algorithm is replaced with a stronger algorithm.

At a minimum, a conforming PKI component shall implement one of the identified digital signature algorithms. For simple PKI clients, it is sufficient to verify one of the identified algorithms; other components are required to generate and verify signatures for one of the identified algorithms. PKI components that claim conformance to the MISPC confidentiality extensions shall support one of the identified key management algorithms. At a minimum, a CA that claims conformance to the MISPC confidentiality extensions must recognize the OID and encoding of at least one of the specified algorithms. A client that claims conformance to the MISPC confidentiality extensions must be able to generate a key pair and assemble a certificate request for one of the specified algorithms.

Conforming components are permitted to implement additional algorithms even if they do not implement all the algorithms identified in this specification. For example, a client that supports one of the specified key agreement algorithms may implement any key transport algorithm it chooses, even if it does not support the specified key transport algorithm.

#### 3.1.2.1 Secure Hash Algorithms

Secure hash algorithms are employed as an interim step in the generation of digital signatures (see 3.1.2.2) for certificates and CRLs. Secure hash algorithms are also used with shared secrets to generate hash-based message authentication codes (HMACs).

This specification requires the use of SHA-1 in the generation of all digital signatures. SHA-1 is fully described in [FIPS 180-1]. In general, SHA-1 is indicated through the object identifier associated with the digital signature algorithm. (see 3.1.2.2).

Where necessary to identify the hash algorithm independently, the ASN.1 object identifier used to identify this hash algorithm is:

## sha1 OBJECT IDENTIFIER ::= { iso(1) identified-organization(3) oiw(14) secsig(3) algorithm(2) 26 }

When this OID appears within the ASN.1 type AlgorithmIdentifier, the parameters component of that type shall be the ASN.1 type NULL.

#### 3.1.2.2 Digital Signature Algorithms

X.509 certificates specify both the algorithm used to sign the certificate (in the **signature** field) and the algorithm of the subject's public key (in the **subjectPublicKeyInfo** field). The two algorithms may be different.

This section specifies the encodings for public keys and digital signatures for four distinct algorithms:

- the Rivest-Shamir-Adelman (RSA) algorithm as specified by Public Key Cryptographic Standard #1 (PKCS-1);
- the Reversible Digital Signature Algorithm (rDSA) specified in ANSI X9.31-1998;
- the Digital Signature Algorithm specified in FIPS 186-1; and
- the Elliptic Curve Digital Signature Algorithm (ECDSA) specified in ANSI X9.62-1998.

To conform to this specification through January 1, 2001, a conforming PKI component shall implement one of the identified digital signature algorithms. (After January 1, 2001, a conforming PKI component shall implement at least one FIPS approved digital signature algorithm. Currently, DSA and rDSA are the FIPS approved digital signature algorithms.)

Note that implementation requirements may vary according to the type of component:

- Conforming clients shall be able to validate signatures of at least one of the identified (FIPS approved after 1/1/2001) algorithms. (To achieve maximum interoperability, it is recommended that clients be capable of validating signatures for all four of the algorithms specified below.)
- Conforming CAs shall be able to sign certificates and Certificate Revocation Lists (CRLs) using at least one of the identified (FIPS approved) algorithms.
- Conforming end entities shall be able to sign with at least one of the identified (FIPS approved) algorithms.

#### RSA (PKCS #1)

The RSA signature algorithm is defined in PKCS #1 [PKCS#1]. Although RSA can be used with several hash algorithms, the only variant used to sign certificates and CRLs conforming to this interoperability specification is RSA with the SHA-1 hash algorithm specified in FIPS 180-1 [FIPS 180]. For this specification, the **sha-1WithRSAEncryption** object identifier is used to identify RSA with SHA-1:

#### pkcs-1 OBJECT IDENTIFIER ::= { iso(1) member-body(2) US(840) rsadsi(113549) pkcs(1) 1 }

sha-1WithRSAEncryption OBJECT IDENTIFIER ::= { pkcs-1(1) 5 }

This object identifier shall appear in the parameterized type **SIGNED** and the **signature** field in both certificates and CRLs signed with RSA. Whenever this object identifier appears as the value for **algorithmIdentifier**, the parameters component shall be **NULL**.

When a certificate or CRL is signed with RSA and SHA-1, the signature shall be generated and encoded as follows:

The certificate or CRL is ASN.1 DER encoded, and is used as the input to the SHA-1 hash function. The SHA-1 output value is ASN.1 encoded as an **OCTET STRING** and the result is encrypted with the RSA algorithms to form the signed quantity. When signing, the RSA algorithm generates an integer y. This signature value is then ASN.1 encoded as a **BIT STRING**, such that the most significant bit in y is the first bit in the bit string and the least significant bit in y is the last bit in the bit string, and included in the **Certificate** or **CertificateList** (in the **signature** field).

(In general the conversion to a bit string occurs in two steps. The integer y is converted to an octet string such that the first octet has the most significance and the last octet has the least significance. The octet string is converted into a bit string such that the most significant bit of the first octet shall become the first bit in the bit string, and the least significant bit of the last octet is the last bit in the bit string.)

When a conforming CA issues a certificate whose **subjectPublicKeyInfo** field contains an RSA public key, the object identifier **rsaEncryption** shall appear as the **algorithmIdentifier** in the **subjectPublickeyInfo** field to identify the key as an RSA public key.

```
rsaEncryption OBJECT IDENTIFIER ::= { pkcs-1 1}
```

Whenever the **rsaEncryption** object identifier is used in the algorithm field of a value of type **AlgorithmIdentifier**, the parameters field shall have ASN.1 type **NULL**.

The RSA public key shall be encoded using the ASN.1 type **RSAPublicKey**:

```
RSAPublicKey ::= SEQUENCE {
modulus INTEGER, -- n
publicExponent INTEGER -- e
}
```

where modulus is the modulus n, and **publicExponent** is the public exponent e. The DER encoded **RSAPublicKey** is the value of the **BIT STRING subjectPublicKey**.

This object identifier is used in public key certificates for both RSA signature keys and RSA

encryption keys. The intended application for the key may be indicated in the key usage field (see sec. 4.2.1.3). The use of a single key for both signature and encryption purposes is not recommended, but is not forbidden.

For interoperability with emerging systems, this specification *strongly recommends* that any component supporting [PKCS#1] also support the rDSA key format, signature format, and signature algorithm. The rDSA key and signature formats are described below; the signature algorithm is specified in [X9.31].

#### <u>rDSA</u>

Recently, a new signature standard based on the RSA algorithm was approved as the X9.31 standard, "X9.31-1998 Digital Signatures Using Reversible Public Key Cryptography for the Financial Services Industry (rDSA)" by ANSI.[X9.31] The Reversible Digital Signature Algorithm (rDSA) standard is based on the RSA algorithm, but specifies an alternative key generation and signature algorithm from that found in [PKCS1].

Following ANSI's recent approval of the ANSI X9.31 standard, the Secretary of Commerce approved an interim modification to FIPS 186 (FIPS 186-1) to approve use of the digital signature technique specified in X9.31 in addition to the algorithm currently specified in FIPS 186. As the algorithms are closely related, cryptographic modules that support [PKCS#1] are expected to support [X9.31] as well, and vice versa. This should ease interoperability problems that may be created by these differences.

Although rDSA could theoretically be used with several hash algorithms, the [X9.31] specifies the SHA-1 hash algorithm specified in FIPS 180-1 [FIPS 180]. Conforming CAs that sign certificates and CRLs with rDSA will always use the SHA-1 hash algorithm. For this specification, the **sha-1WithrDSA** object identifier is used to identify rDSA with SHA-1:

```
tC68arc OBJECT IDENTIFIER ::= { iso(1) member-body(2) US(840)
rsadsi(113549) pkcs(1) 1 }
```

### sha-1WithrDSA OBJECT IDENTIFIER ::= { tC68arc ??? }

This object identifier shall appear in the parameterized type **SIGNED** and the **signature** field in both certificates and CRLs signed with RSA. Whenever this object identifier appears as the value for **algorithmIdentifier**, the parameters component shall be **NULL**.

When a certificate or CRL is signed with RSA and SHA-1, the signature shall be generated as specified in X9.31. The result shall be encoded as specified above for [PKCS1].

When a conforming CA issues a certificate whose **subjectPublicKeyInfo** field contains an rDSA public key, the object identifier **rDSAkey** shall appear as the **algorithmIdentifier** in the **subjectPublickeyInfo** field to identify the key as an rDSA public key.

#### rDSAkey OBJECT IDENTIFIER ::= { tC68arc ??? }

Whenever the **rDSA** object identifier is used in the algorithm field of a value of type **AlgorithmIdentifier**, the parameters field shall have ASN.1 type **NULL**.

Where a certificate contains an rDSA public key, the key shall be encoded in the **subjectPublicKey** field as specified in [????].

For interoperability with legacy systems, this specification *strongly recommends* that any component supporting [X9.31] support the [PKCS1] key format, signature format, and signature algorithm.

#### DSA

The Digital Signature Algorithm is defined in FIPS 186-1 [FIPS186]. The ASN.1 object identifier used to identify DSA public keys shall be:

#### id-dsa ID ::= { iso(1) member-body(2) us(840) x9-57(10040) x9cm(4) dsa(1) }

The Digital Signature Standard (DSS) [FIPS186] specifies that DSA shall be used with the SHA-1 hash algorithm. The ASN.1 object identifier used to identify DSA signatures shall be:

#### id-dsa-with-sha1 ID ::= {

#### iso(1) member-body(2) us(840) x9-57(10040) x9cm(4) id-dsa-with-sha1(3) }

The AlgorithmIdentifier within subjectPublicKeyInfo is the only place within a certificate where id-dsa shall be used. The id-dsa algorithm syntax includes optional parameters. These parameters are commonly referred to as p, q, and g. Where omitted, the parameters component shall be omitted entirely. If the DSA algorithm parameters are absent from the subjectPublicKeyInfo AlgorithmIdentifier and the CA signed the subject certificate using DSA, then the certificate issuer's DSA parameters apply to the subject's DSA key. If the DSA algorithm parameters are absent from the CA signed the certificate using a signature algorithm other than DSA, then clients shall not validate the certificate. The parameters are included using the following ASN.1 structure:

<b>DSAParameters</b> :	= SEQUENCE {
prime1	INTEGER, modulus p
prime2	INTEGER, modulus q
base	INTEGER } – base g

The id-dsa-with-sha1 algorithm identifier shall be used in the SIGNED parameterized type (e.g., in the signature on a certificate or CRL) and the signature fields of certificates and CRLs. Where the id-dsa-with-sha1 algorithm identifier appears as the algorithm field in an AlgorithmIdentifier, the encoding shall omit the parameters field. That is, the AlgorithmIdentifier shall be a SEQUENCE of one component - the OBJECT IDENTIFIER id-dsa-with-sha1. The DSA parameters in the certificate of the issuer shall apply to the verification of the signature.

The DSA public key shall be ASN.1 DER encoded as an INTEGER; this encoding shall be used as the contents (i.e., the value) of the subjectPublicKey component (a BIT STRING) of the SubjectPublicKeyInfo data element.

#### DSAPublicKey ::= INTEGER -- public key Y

When signing, the DSA algorithm generates two values. These values are commonly referred to as r and s. To easily transfer these two values as one signature, they shall be ASN.1 encoded using the following ASN.1 structure:

```
Dss-Sig-Value ::= SEQUENCE {
r INTEGER,
s INTEGER }
```

The encoded signature is conveyed as the value of the **BIT STRING** in the **SIGNED** parameterized type in a **certificate** or **CertificateList**.

#### ECDSA

The Elliptic Curve Digital Signature Algorithm (ECDSA) is defined in the draft ANSI X9.62 standard [X9.62]. The ASN.1 object identifier used to identify the ECDSA algorithm shall be:

#### ansi-X9-62 OBJECT IDENTIFIER ::= { iso(1) member-body(2) us(840) 10045 }

When used to sign certificates, CRLs, or PKI messages, the ECDSA shall be used with the SHA-1 hash algorithm. When ECDSA and SHA-1 are used to sign an X.509 certificate, CRL, or PKI message, the signature shall be identified by the value **ecdsa-with-SHA1**, as defined below:

id-ecSigType OBJECT IDENTIFIER ::= { ansi-X9-62 signatures(4) } ecdsa-with-SHA1 OBJECT IDENTIFIER ::= { id-ecSigType 1 }

When the **ecdsa-with-SHA1** algorithm identifier is used in the **SIGNED** parameterized **TYPE** (e.g., in the signature on a certificate or CRL) it shall have **NULL** parameters. The ECDSA parameters in the certificate of the issuer shall apply to the verification of the signature.

When signing, the ECDSA algorithm generates two values. These values are commonly referred to as r and s. To easily transfer these two values as one signature, they shall be ASN.1 encoded using the following ASN.1 structure:

```
Ecdsa-Sig-Value ::= SEQUENCE {
r INTEGER,
s INTEGER }
```

When certificates contain an ECDSA public key, the **id-ecPublicKey** algorithm identifier shall be used. The **id-ecPublicKey** algorithm identifier is defined as follows:

```
id-public-key-type OBJECT IDENTIFIER ::= { ansi-X9.62 2 }
id-ecPublicKey OBJECT IDENTIFIER ::= { id-publicKeyType 1 }
```

The elliptic curve public key (an **ECPoint** which is an **OCTET STRING**) is mapped to a **subjectPublicKey** (a **BIT STRING**) as follows: the most significant bit of the **OCTET STRING** becomes the most significant bit of the **BIT STRING**, etc.; the least significant bit of the **OCTET STRING** becomes the least significant bit of the **BIT STRING**.

ECDSA requires use of certain parameters with the public key. When the parameters are inherited, the **parameters** field shall contain **implictlyCA**, which is the ASN.1 value **NULL**. When parameters are specified by reference, the parameters field shall contain the **namedCurve** choice, which is an an object identifier. When the parameters are explicitly included, they shall be encoded in the ASN.1 structure **ECParameters**:

```
Parameters ::= CHOICE {
ecParameters ECParameters,
namedCurve CURVES.&id({CurveNames}),
implicitlyCA NULL }
```

The parameters may be explicitly included in the certificate using the following ASN.1 structure:

```
ECParameters ::= SEQUENCE {
       version
                      INTEGER { ecpVer1(1) }
                                                  (ecpVer1),
                                     -- version is always 1
       fieldID
                      FieldID
                                  { {FieldTypes} },
                                     -- identifies the finite field over
                                     -- which the curve is defined
       curve
                      Curve,
                                     -- coefficients a and b of the elliptic curve
                      ECPoint,
                                     -- specifies the base point P
       base
                                     -- on the elliptic curve
                                     -- the order n of the base point
       order
                      INTEGER.
       cofactor
                      INTEGER.
       . . .
       }
FieldElement ::= OCTET STRING
Curve ::= SEQUENCE {
              FieldElement,
       а
       b
              FieldElement.
              BIT STRING OPTIONAL
       seed
```

ECPoint ::= OCTET STRING

}

The components of type **ECParameters** have the following meanings:

- version specifies the version number of the elliptic curve parameters. It shall have the value 1 for this version of the Standard. The notation above creates an INTEGER named ecpVer1 and gives it a value of one. It is used to constrain version to a single value.
- **fieldID** identifies the finite field over which the elliptic curve is defined. Finite fields are represented by values of the parameterized type **FieldID**, constrained to the values of the objects defined in the information object set **FieldTypes**. Additional detail regarding **fieldID** is provided below.
- **curve** specifies the coefficients *a* and *b* of the elliptic curve *E*. Each coefficient shall be represented as a value of type **FieldElement**, an **OCTET STRING**. **seed** is an optional parameter used to derive the coefficients of a randomly generated elliptic curve.
- **base** specifies the base point *P* on the elliptic curve. The base point shall be represented as a value of type **ECPoint**, an **OCTET STRING**.
- order specifies the order *n* of the base point.
- cofactor is the integer  $h = \#E(F_q)/n$ .

The AlgorithmIdentifier within subjectPublicKeyInfo is the only place within a certificate where the parameters may be used. If the ECDSA algorithm parameters are encoded as implicitlyCA in the subjectPublicKeyInfo AlgorithmIdentifier and the CA signed the subject certificate using ECDSA, then the certificate issuer's ECDSA parameters apply to the subject's ECDSA key. If the ECDSA algorithm parameters are encoded as implicitlyCA in the subjectPublicKeyInfo AlgorithmIdentifier and the CA signed the certificate using a signature algorithm other than ECDSA, then clients shall not validate the certificate.

#### FieldID { FIELD-ID:IOSet } ::= SEQUENCE {

```
fieldType FIELD-ID.&id({IOSet}),
  parameters FIELD-ID.&Type({IOSet}{@fieldType}) OPTIONAL
}
FieldTypes FIELD-ID ::= {
  { Prime-p IDENTIFIED BY prime-field } |
  { Characteristic-two IDENTIFIED BY characteristic-two-field },
...
}
```

#### FIELD-ID ::= TYPE-IDENTIFIER

**FieldID** is a parameterized type composed of two components, **fieldType** and **parameters**. These components are specified by the fields **&id** and **&Type**, which form a template for defining sets of information objects, instances of the class **FIELD-ID**. This class is based on the useful information object **class TYPE-IDENTIFIER**, described in X.681 Annex A. In an instance of **FieldID**, "**fieldType**" will contain an object identifier value that uniquely identifies the type contained in "**parameters**." The effect of referencing "**fieldType**" in both components of the **fieldID** sequence is to tightly bind the object identifier and its type.

The information object set **FieldTypes** is used as the single parameter in a reference to type **FieldID**. **FieldTypes** contains two objects followed by the extension marker ("..."). Each object, which represents a finite field, contains a unique object identifier and its associated type. The values of these objects define all of the valid values that may appear in an instance of **fieldID**. The extension marker allows backward compatibility with future versions of this standard which may define objects to represent additional kinds of finite fields.

The object identifier **id-fieldType** represents the root of a tree containing the object identifiers of each field type. It has the following value:

id-fieldType OBJECT IDENTIFIER ::= { ansi-X9-62 fieldType(1) }

The object identifiers **prime-field** and **characteristic-two-field** name the two kinds of fields defined for this specification. They have the following values:

```
prime-field OBJECT IDENTIFIER ::= { id-fieldType 1 }
characteristic-two-field OBJECT IDENTIFIER ::= { id-fieldType 2 }
Prime-p ::= INTEGER -- Field size p
Characteristic-two ::= SEQUENCE {
                                      -- Field size 2<sup>m</sup>
 m
               INTEGER,
               CHARACTERISTIC-TWO.&id({BasisTypes}),
 basis
  parameters
               CHARACTERISTIC-TWO.&Type({BasisTypes}{@basis})
}
BasisTypes CHARACTERISTIC-TWO::= {
 { NULL
                     IDENTIFIED BY onBasis } |
 { Trinomial
                     IDENTIFIED BY tpBasis } |
 { Pentanomial
                     IDENTIFIED BY ppBasis },
}
Trinomial ::= INTEGER
Pentanomial ::= SEQUENCE {
```

```
k1 INTEGER,
k2 INTEGER,
k3 INTEGER
```

}

#### CHARACTERISTIC-TWO ::= TYPE-IDENTIFIER

The object identifier **id-characteristic-two-basis** represents the root of a tree containing the object identifiers for each type of basis for the characteristic-two finite fields. It has the following value:

```
id-characteristic-two-basis OBJECT IDENTIFIER ::= {
     characteristic-two-field basisType(1) }
```

The object identifiers onBasis, tpBasis and ppBasis name the three kinds of basis for characteristic-two finite fields defined by [X9.62]. They have the following values:

onBasis OBJECT IDENTIFIER ::= { id-characteristic-two-basis 1 } tpBasis OBJECT IDENTIFIER ::= { id-characteristic-two-basis 2 } ppBasis OBJECT IDENTIFIER ::= { id-characteristic-two-basis 3 }

# 3.1.2.3 Key Agreement Algorithms

Key agreement algorithms are a class of algorithms where both parties contribute to the derivation of a shared key. X.509 certificates for the purpose of key agreement specify the algorithm of the subject's public key (in the **subjectPublicKeyInfo** field) and convey the public key.

Two key agreement algorithms are recognized by this specification. They are the Diffie-Hellman algorithm and its elliptic curve analog. If the **keyUsage** extension is present in a certificate which conveys a DH public key, the value will be **keyAgreement**.

#### Diffie-Hellman

The Diffie-Hellman algorithm (DH) is defined in the draft ANSI X9.42 standard [X9.42] The ASN.1 object identifier supported by this standard is

# dhpublicnumber OBJECT IDENTIFIER ::= { iso(1) member-body(2) us(840) ansi-x942(10046) number-type(2) 1 }

The **dhpublicnumber** object identifier is intended to be used in the algorithm field of a value of type **AlgorithmIdentifier**. The parameters field of that type, which has the algorithm-specific syntax **ANY DEFINED BY** algorithm, would have ASN.1 type **DHParameter** for this algorithm.

```
DHParameter ::= SEQUENCE {
prime INTEGER, -- p
base INTEGER, -- g }
```

The fields of type **DHParameter** have the following meanings:

• **prime** is the prime p.

• **base** is the base g.

This specification requires that the parameters field be present whenever an **AlgorithmIdentifier** field contains the **dhpublicnumber** object identifier.

The Diffie-Hellman public key (an INTEGER) is mapped to a subjectPublicKey (a BIT STRING) as follows: the most significant bit (MSB) of the INTEGER becomes the MSB of the BIT STRING; the least significant bit (LSB) of the INTEGER becomes the LSB of the BIT STRING.

# Elliptic Curve Diffie-Hellman

The Elliptic Curve Diffie-Hellman algorithm (ECDH) is defined in the draft ANSI X9.63 standard [X9.63]; key encoding and algorithm identifiers are defined in the draft ANSI X9.62 standard [X9.62]. When certificates contain an ECDH public key, the **id-ecPublicKey** algorithm identifier shall be used. The **id-ecPublicKey** algorithm identifier is defined as follows:

# id-public-key-type OBJECT IDENTIFIER ::= { ansi-X9.62 2 }

# id-ecPublicKey OBJECT IDENTIFIER ::= { id-publicKeyType 1 }

The elliptic curve public key (an **ECPoint** which is an **OCTET STRING**) is mapped to a **subjectPublicKey** (a **BIT STRING**) as follows: the most significant bit of the **OCTET STRING** becomes the most significant bit of the **BIT STRING**, etc.; the least significant bit of the **OCTET STRING** becomes the least significant bit of the **BIT STRING**.

When parameters are included, they use the **ECParameters** structure, as defined for ECDSA. However, they shall include the **cofactor** parameter, which is not required for signature keys. The **cofactor** parameter is used to validate parameters.

If the ECDH algorithm parameters are absent from the **subjectPublicKeyInfo AlgorithmIdentifier** and the CA signed the subject certificate using ECDSA, then the certificate issuer's ECDSA parameters apply to the subject's ECDH key. If the ECDH algorithm parameters are absent from the **subjectPublicKeyInfo AlgorithmIdentifier** and the CA signed the certificate using a signature algorithm other than ECDSA, or signed with ECDSA but the **cofactor** parameter is not present, then clients shall not validate the certificate.

ECDH certificates are differentiated from ECDSA certificates by the value of the **keyUsage** extension (see Sec. 3.1.3.1.), which shall be **keyAgreement** or **digitalSignature**. As with the RSA algorithm, the use of a single key for both signature and encryption purposes is not recommended, but is not forbidden.

# 3.1.2.4 Key Transport Algorithms

X.509 certificates can convey public key encryption keys for the purpose of key transport. In this scenario, the initiator uses the public key in the receiver's key management certificate to encrypt a symmetric key. The receiver uses their private key to recover the symmetric key.

Two algorithms for key transport are recognized by this specification - the RSA algorithm and the X9.44 Key Establishment Using Factoring-Based Public Key Cryptography.

# <u>RSA</u>

RSA public keys in the **subjectPublicKeyInfo** field of a key management certificate are encoded identically to RSA signature keys and are identified by the same OID (see Sec. 3.1.2.2.) The purpose of the key is differentiated by the value of the **keyUsage** extension (see Sec. 3.1.3.1.), which shall be **keyExchange** or **digitalSignature**. As stated in Sec. 3.1.2.2, the use of a single key for both signature and encryption purposes is not recommended, but is not forbidden.

# X9.44 Key Establishment Using Factoring-Based Public Key Cryptography

X9.44 public keys in the **subjectPublicKeyInfo** field of a key management certificate are encoded identically to RSA signature keys but are identified by the same OID as rDSA keys (see Sec 3.1.2.2.) The purpose of the key is further specified by the value of the **keyUsage** extension (see Sec. 3.1.3.1.), which shall be **keyExchange** or **digitalSignature**. As stated in Sec. 3.1.2.2, the use of a single key for both signature and encryption purposes is not recommended, but is not forbidden.

# 3.1.2.5 Message Authentication Algorithms

This specification recognizes two message authentication algorithms: the DES-MAC and the SHA1-HMAC. The preferred algorithm is the SHA1-HMAC. The DES-MAC is included for backwards compatibility.

# SHA-1 HMAC

This algorithm provides integrity by computing a SHA-1 HMAC (as specified by [RFC2104]) on data. This algorithm is identified by the following object identifier:

# SHA1-HMAC OBJECT IDENTIFIER ::= { 1 3 6 1 5 5 8 12 }

The length of the mac shall be 96 bits for this specification.

# DES MAC

This algorithm provides integrity by computing a DES MAC (as specified by [FIPS-113]) on data. This algorithm is identified by the following object identifier:

```
DES-MAC OBJECT IDENTIFIER ::= {
iso(1) identified-organization(3) oiw(14) secsig(3) algorithm(2) 10
-- carries length in bits of the MAC as
-- an INTEGER parameter, constrained to 32
-- for this specification
}
```

The parameters field contains an **INTEGER** specifying the length of the MAC; this length shall be 32 bits for this specification.

# 3.1.2.6 Symmetric Encryption Algorithms

This specification requires the use of symmetric encryption to implement challenge-response protocols and protect private keys in transit. For this specification, all symmetric encryption is performed using the Triple Data Encryption Algorithm (tDEA) in electronic codebook mode.

The tDEA algorithm is defined in [X9.52], *Triple Data Encryption Algorithm Modes Of Operation*. The tDEA is based upon the DES algorithm and uses three 56 bit keys: K1; K2; and K3. For this specification, the tDEA algorithm is used with the two key option, where K1 = K3. tDEA keys are never conveyed in X.509 certificates. In some cases, it may be necessary to convey the symmetric key material (K1 and K2) in a PKI message. In such a case, the key material must itself be encrypted.

The remainder of this subsection describes how to specify that data is encrypted with tDEA in ecb mode, indicate the two key option, and encode keying material.

An Algorithmldentifier is required to specify the algorithm used for data encryption. The Algorithmldentifier is composed of an OID and parameters. To indicate that data is encrypted under tDEA in ecb mode, the OID tECB is used; the keying option is specified through the parameters structure ECBParms. The ASN.1 structures are provided below.

```
TDEAldentifier ::= AlgorithmIdentifier {{ TDEAModes }}
TDEAModes ALGORITHM-ID ::= {
 { OID tECB PARMS ECBParms } | -- mode 1 --
  { OID tCBC PARMS TDEAParms } | -- mode 2 --
 { OID tCBC-I PARMS TDEAParms } | -- mode 3 --
 { OID tCFB PARMS CFBParms } | -- mode 4 --
 { OID tCFB-P PARMS CFBParms } | -- mode 5 --
  {OID tOFB PARMS TDEAParms } -- mode 6 --
 { OID tOFB-I PARMS TDEAParms }, -- mode 7 --
}
ECBParms ::= TDEAParms (WITH COMPONENTS {
       ..., ivGeneration ABSENT })
-- Note : the syntax for ivGeneration is not given here,
-- as it is not used in this specification
TDEAParms ::= SEQUENCE {
  keyingOptions KeyingOptions OPTIONAL,
 ivGeneration [0] IVGeneration OPTIONAL
}
-- only the 2-key option is used in this specification
KeyingOptions ::= BIT STRING {
 option-1 (0), -- (3-key) K1, K2 and K3 are independent keys
 option-2 (1), -- (2-key) K1 and K2 are independent and K3 = K1
 option-3 (2) -- (1-key) K1 = K2 = K3
}
id-ansi-x952 OBJECT IDENTIFIER ::= {
  iso(1) member-body(2) us(840) ansi-x952(10047) }
```

```
mode OBJECT IDENTIFIER ::= { id-ansi-x952 1 }
```

```
tECB OBJECT IDENTIFIER ::= { mode 1 }
```

Where tDEA key material must be encoded, the keys are simply concatenated. TwoKeys = [K1 | K2], where the most significant bit of TwoKeys is the most significant bit of K1.

# 3.1.3 Certificate Extensions

A set of standardized extensions has been developed and is specified in an amendment to X.509 [DAM]. Extensions have three components: extension name, criticality flag, and extension value. As specified in the amendment to X.509 [DAM], clients shall not validate certificates that contain an extension with the criticality flag set, unless the client can process that extension.

The standardized extensions that have been defined may be divided into four categories: key and policy information; subject and issuer characteristics; certification path constraints; and CRL identification extensions.

# 3.1.3.1 Key and Policy Information

These extensions provide information to identify a particular public key and certificate. They can be used to identify a particular public key/certificate for a CA which has several certificates. This may help a client to find the particular CA certificate needed to establish a certification path. These extensions may restrict the purposes for which a key may be used, and provide information in CA certificates about equivalent policies.

#### Authority Key Identifier

The **authorityKeyldentifier** extension provides a means of identifying the particular private key used to sign a certificate. The identification can be based on either the key identifier or on the issuer name and serial number. The key identifier method shall be used in certificates conforming to this interoperability specification. This extension is used where an issuer has multiple signing keys (either due to multiple concurrent key pairs or due to changeover). CAs shall be capable of generating this extension, and clients shall be capable of finding and validating certification paths where the issuing CA has several digital signature keys. It is recommended that clients be able to process either the key identifier or the certificate issuer plus certificate serial number form of key identifier to help find certification paths.

# Subject Key Identifier

This field enables differentiation of keys held by a subject. This field shall be included in every certificate issued. This extension shall be noncritical.

#### Key Usage

The **keyUsage** extension defines restrictions on the use of the key contained in the certificate based on policy and/or usage (e.g., signature, encryption). CAs shall support the generation of this extension and clients shall be capable of processing it. While **KeyUsage** is defined as a **BIT** 

**STRING**, conforming CAs shall set only one value within this string in end-entity certificates. For example, **KeyUsage** shall not be both **digitalSignature** and **dataEncipherment** in an end-entity certificate. This extension shall be set to critical.

# Private Key Usage Period

The **privateKeyUsagePeriod** extension applies only to digital signature keys. A signature on a document that purports to be dated outside the private key usage period is not valid.<sup>19</sup> CAs may generate certificates containing this extension but conforming clients are not required to process it.

# Extended Key Usage

The **extendedKeyUsage** extension defines application-specific restrictions on the use of keys contained in a certificate. When this extension is used, interoperability is not a factor. Conforming PKI components are not required to support this extension.

# Certificate Policies

The **certificatePolicies** extension contains one or more object identifiers (OIDs). Each OID indicates a policy under which the certificate has been issued. CAs shall be able to generate certificates with one or more instances of **policyldentifier**. CAs can include the special policy OID anyPolicy. The anyPolicy OID may be thought of as a wildcard which matches every policy.Clients shall be capable of processing **policyldentifier** fields against a list of acceptable policies. (The list of policies is dependent upon on application requirements.) Clients shall compare the policy identifier(s) in the certificate to that list. If a list of acceptable policies is provided, the client shall validate the certification path only if at least one of the acceptable policies appears in the **certificatePolicies** field of each certificate or maps to one of the policies in the **certificatePolicies** field. If the special policy OID **anyPolicy** appears in the **certificatePolicies** in the list of acceptable policies in the list of a certificate and any policy is not inhibited (see section 3.1.3.3), then all policies in the list of acceptable policies are considered to be matched.

Conforming components are required to process the **policyQualifiers** subfield of **certificatePolicies** if present, and shall support the policy qualifiers **id-pkix-cps** and **id-pkix-unotice** (see [RFC2459].) Conforming CAs need not be able to generate this subfield.

# Policy Mapping

**3.1.3.2** This noncritical extension is used in CA certificates. It lists pairs of object identifiers; each pair includes an **issuerDomainPolicy** and a **subjectDomainPolicy**. The pairing indicates that the issuing CA considers its **issuerDomainPolicy** equivalent to the subject CA's **subjectDomainPolicy**. CAs shall be capable of generating the **policyMappings** extension. Clients shall be capable of processing this extension.*Certificate Subject and Issuer Characteristics* 

The subjectAltName, issuerAltName, subjectDirectoryAttributes, and authorityInformationAccess are all noncritical extensions. They provide additional information about other names and

<sup>&</sup>lt;sup>19</sup> Note that verification of time associated with a signature implies use of a notary or trusted timestamp. Both are outside the scope of this specification.

characteristics of the subject and issuer.

# Alternative Name

The **subjectAltName** and **issuerAltName** extensions allow additional identities to be bound to the subject and issuer of the certificate. Defined options include an RFC822 [RFC 822] name (electronic mail address), a DNS name, and a uniform resource identifier (URI.) Multiple instances may be included. Whenever such identities are to be bound in a certificate, the **subjectAltName** or **issuerAltName** fields shall be used.<sup>20</sup>

The **subjectAltName** and **issuerAltName** extensions are normally noncritical in certificates conforming to this interoperability specification. An implementation which recognizes these extensions need not be able to process all the alternatives of the choice. If the alternative used is not supported by the implementation, the extension field is ignored.

This specification defines the semantics associated with an **issuerAltName** field containing a URI. The URI specifies the location of the issuer's certificate(s) which contain the public key material corresponding to the private key used to sign the certificate. The semantics associated with other classes of identities, or any **subjectAltName** entries, are not defined in this specification.

If a CA's certificates are not available from a well-known X.500 directory service, the CA shall include URI alternative names specifying the location of the issuer's certificate(s). Clients are required to process the URI alternative name format and must recognize the LDAP URL [RFC1959]. Clients are not required to recognize any other URI formats.

# Subject Directory Attributes

The **subjectDirectoryAttributes** extension may hold any information about the subject where that information has a defined X.500 Directory attribute. This extension is always noncritical. Implementation and use of this extension is optional.

# Authority Information Access

The **authorityInformationAccess** extension may hold any information about how to access CA information and services for the issuer of the certificate in which the extension appears. Information and services may include on-line validation services and CA policy data. (The location of CRLs is not specified in this extension; that information is provided by the cRLDistributionPoints extension.)

# 3.1.3.3 Certification Path Constraints

The **basicConstraints**, **nameConstraints** and **policyConstraints** extensions all apply restrictions to valid certification paths.

# Basic Constraints

The basicConstraints extension tells whether the subject of the certificate is a CA through the cA

<sup>&</sup>lt;sup>20</sup> X.509 allows null certificate **subject** or **issuer** fields accompanied by a critical **subjectAltName** or **issuerAltName** giving the name in an alternative format. Such certificates are not supported by this interoperability specification.

component and the lengths of certification paths through the **pathLenConstraint** component. CAs shall support the generation of the **basicConstraints** extension in certificates and clients shall be capable of processing it. The **pathLenConstraint** component is meaningful only if **cA** is set to TRUE.

The **basicConstraints** extension shall be included in all CA certificates. CA certificates shall contain a **basicConstraints** extension with the **cA** component set to TRUE. The **basicConstraints** extension may be included in end entity certificates. Where a **basicConstraints** extension appears in an end entity certificate, it shall contain an empty **SEQUENCE** value. The **basicConstraints** extension shall be marked as critical in all CA certificates.

# Name Constraints

The **nameConstraints** field applies only to CA certificates. It indicates a name space in which all subsequent certificates in a certification path must be located. CAs shall be capable of including this field in certificates and clients shall be capable of processing it. If used, it shall be critical.

# Policy Constraints

The **policyConstraints** field applies only to CA certificates. It can be used to constrain the interpretation of policy in two ways: it can be used to prohibit policy mapping or require that each certificate in a path contain an acceptable policy identifier.

This extension contains two field: inhibitPolicyMapping and requireExplicitPolicy. If the inhibitPolicyMapping field is present, the value indicates the number of additional certificates that may appear in the path before policy mapping is no longer permitted. If the requireExplicitPolicy field is present, subsequent certificates shall include an acceptable policy identifier. The value of requireExplicitPolicy indicates the number of additional certificates that may appear in the path before an explicitPolicy indicates the number of additional certificates that may appear in the path before an explicit policy is required.

CAs shall be capable of supporting the issuance of certificates with this extension, and clients shall be capable of processing this extension. If used, it shall be critical.

# Inhibit anyPolicy

This extension inhibits the use of **anyPolicy** in future certificates. This critical extension may appear in certificates issued to CAs. The inhibit any-policy indicates that the special **anyPolicy** OID, with the value {2 5 29 32 0}, is not considered an explicit match for other certificate policies. The value is an **INTEGER** and indicates the number of additional certificates that may appear in the path before any-policy is no longer permitted. For example, a value of one indicates that any-policy may be processed in certificates issued by the subject of this certificate, but not in additional certificates in the path. CAs shall be capable of generating this extension. Clients shall be capable of processing this extension.

# 3.1.3.4 CRL Identification Extensions

These extensions include information in a certificate about where to obtain the Certificate Revocation List (CRL) that applies to that certificate. They facilitate the division of a CA's potentially large CRL into several shorter CRLs, by identifying in the certificate which CRL

applies to a certificate and stating the name of the CRL issuer (which may be a CA other than the CA that issued the certificate).

# CRL Distribution Points

The **cRLDistributionPoints** extension identifies the CRL distribution point or points to which clients should refer to ascertain if a certificate has been revoked. This field has three component fields: **distributionPoint**, **reasons** and **cRLIssuer**.

- The distributionPoint component identifies the location from which the CRL can be obtained. If this field is absent, the CRL distribution point name defaults to the issuer name. This extension provides a mechanism to divide the CRL into manageable pieces if the CA has a large constituency.
- The **reasons** component identifies the reasons for revocation covered by the CRL issued by the corresponding **distributionPoint**. If the **reasons** component is absent, the corresponding **distributionPoint** distributes a CRL which will contain an entry for this certificate, if it has been revoked for any reason. Clients are not required to process the **reasons** component.
- The **cRLIssuer** component identifies the authority that issues and signs the CRL. If this component is absent, the CRL issuer name defaults to the certificate issuer name. One use for this component is to allow the construction of consolidated CRLs, that include certificates issued by more than one CA.

CAs shall include the **cRLDistributionPoints** extension with a **distributionPoint** component. If a CA's CRLs are not available from a well-known X.500 directory service, the CA shall include URI alternative names specifying the location of the current CRL for this certificate in the **distributionPoint** component. Clients shall be able to process the **cRLDistributionPoints** extension; they must recognize the URI format and process at a minimum the LDAP URI. Clients shall be able to use distribution point CRLs and validate CRLs where the **cRLIssuer** component is used. See section 3.2.2 below for a further discussion of distribution points.

Extension	Used By	Use	Critical
Key and Policy Information			
keyldentifier	all	l identifies the key used to sign this certificate (the signing CA may have several keys)	
authorityKeyIdentifier	all	unique with respect to authority.	
authorityCertIssuer	all	identifies issuing authority of CA's certificate; alternative to key identifier	
authorityCertSerialNumber	all	used with authorityCertIssuer	
subjectKeyldentifier	all	enables differentiation of different keys for same subject. Must be unique for subject.	No
keyUsage	all	defines allowed purposes for use of key (e.g., digital signature, key agreement)	Yes*
extendedKeyUsage	all	defines application-specific purposes for keys	No*
privateKeyUsagePeriod	all	digital signature keys only. Signatures on that purport to be dated outside the period are invalid.	No*
certificatePolicies	all	policy identifiers and qualifiers that identify and qualify policies applying to the certificate	No*
policyldentifiers	all	the OID of a policy.	
policyQualifiers	all	more information about the policy	
policyMappings	CA	indicates equivalent policies	No
Certificate Subject and Issuer Ch	aracteri	istics	
subjectAltName	all	used to list alternative names (e.g., rfc822 name, X.400 address, IP address)	No*
issuerAltName	all	used to list alternative names	No*
subjectDirectoryAttributes	all	any attributes (e.g., supported algorithms)	No
authorityInfoAccess	all	How to access issuer's information and services	No
Certification Path Constraints			
basicConstraints	all	constraints on subject's role & path lengths	Yes*
cA	all	distinguish CA from end entity cert.	
pathLenConstraint	CA	max. number of following CAs in cert. path; 0 indicates that CA only issues end entity certs.	
nameConstraints	CA	constrains names in certs issued by subsequent CAs	Yes*
permittedSubtrees	CA	names outside indicated subtrees are forbidden	]
excludedSubtrees	CA	indicates disallowed subtrees	1
policyConstraints	all	constrains certificates issued by subsequent CAs	Yes*
requireExplicitPolicy	all	All certificates in the path must contain an acceptable policy identifier	
inhibitPolicyMapping	all	prevent policy mapping in following certs.	]
inhibitAnyPolicy	CA	prevent matching policies with the anyPolicy OID	Yes
CRL Identification			
crlDistributionPoints	all	divides long CRL into shorter lists	No*
distributionPoint	all	location from which CRL can be obtained	1
reasons	all	reasons for cert. inclusion in CRL	1
	1		-

# Table 3-1 Summary of Standardized Certificate Extensions

NOTES:

\*

Standard allows either critical or noncritical. Indication is for use in interoperable implementations.

Extension	Certificate	Client
Key and Policy Information		
authorityKeyldentifier		
authorityKeyIdentifier	to be included in all certs issued: a random number large enough to generally be globally unique	optional - may be used to help find cert. paths where issuer has multiple certs. (1)
authorityCertIssuer	not used	optional - used to find cert. paths
authorityCertSerialNumber	not used	where issuer has multiple certs. (1)
subjectKeyldentifier	to be included in all certs issued: a random number large enough to generally be globally unique	optional: used with CRLs to identify revoked certificates.
keyUsage	supported	supported
extendedKeyUsage	not used	not used
privateKeyUsagePeriod	supported	optional
certificatePolicies		
policyldentifiers	supported	supported; compared during cert. path validation with a list of acceptable policies
policyQualifiers	supported	supported (see 3.1.3.1)
policyMappings	supported	supported
Certificate Subject and Issuer Ch	aracteristics	
subjectAltName	supported	not used
issuerAltName	supported	not used
subjectDirectoryAttributes	not used	not used
authorityInfoAccess	supported	optional
Certification Path Constraints		
basicConstraints		
cA	used in all certificates	supported
pathLenConstraint	supported	supported
nameConstraints		
permittedSubtrees	supported	supported
excludedSubtrees	supported	supported
policyConstraints		· · ·
requireExplicitPolicy	supported	supported
InhibitPolicyMapping	supported	supported
inhibitAnyPolicy	supported	supported
CRL Identification	· · · · · · · · · · · · · · · · · · ·	
cRLDistributionPoints		
distributionPoint	supported	supported
reasons	supported	supported
cRLIssuer	supported	supported

# Table 3-2 Use of Standardized Certificates by the MISPC

#### NOTES:

For Certificates, "supported" means that CAs shall be able to issue certificates that contain this extension. For clients, "supported" means that the client shall be capable of processing this extension.

(1) Clients shall be capable of finding certification paths where CAs have multiple certificates, whether or not they use this extension to do so.

# 3.1.3.5 Summary of Certificate Extension Use

Table 3-1 summarizes the standardized certificate extensions, while Table 3-2 summarizes the use by the MISPC of standardized extensions for certificates and clients.

# 3.2 Certificate Revocation List (CRL)

Certificate Revocation Lists (CRL) are used to list unexpired certificates that have been revoked or placed on "hold." Certificates may be revoked for a variety of reasons, ranging from routine administrative revocations, (when the certificate's subject leaves the issuing organization, or when responsibilities and certificate values change), to situations where the private key is compromised. A "hold" indicates the CA will not vouch for the binding of the certificate subject and public key at this time.

The X.509 v2 certificate revocation list format is augmented by several optional extensions, similar in concept to those defined for certificates. CAs shall be able to generate X.509 v2 CRLs as specified below, and clients shall be capable of processing them when validating certification paths. The CA that issues a CRL is not necessarily the CA that issued the revoked certificate, and some CAs may issue only CRLs. The X.509 v2 CRL includes the following:

Version Issuer Signature Algorithm Issuer Distinguished Name This Update Next Update Revoked Certificates, a sequence of zero or more of the following sequence: Certificate Serial Number Revocation Date CRL Entry Extensions (optional) CRL Extensions (optional) Issuer's Signature on all the above listed fields

# 3.2.1 CRL Fields

The X.509 v2 CRL ASN.1 syntax is given in Appendix A. For signature calculation, the data that is to be signed is ASN.1 DER encoded. ASN.1 DER encoding is a tag, length, value encoding system for each element.

The following items describe the use of the X.509 v2 CRL.

# Version

This field describes the version of the encoded CRL. The value of this field shall be 1, indicating a v2 CRL.

# <u>Signature</u>

The **signature** field contains the algorithm identifier for the algorithm used to sign the CRL. The contents are identical to the contents of the certificate **signature** field. Refer to <u>Signature</u> in section 3.1.1 for information about this field. The CRL may be signed with any of the algorithms

identified in section 3.1.2.2; in general, the CA should sign the CRL with the same algorithm used to sign the certificates. Refer to section 3.1.2.2 for the signature algorithm object identifiers. The **parameters** subfield of the CRL **signature** field shall not be used to pass DSA or ECDSA parameters. DSA or ECDSA parameters shall be obtained from the **subjectPublicKeyInfo** field of the certificate of the issuer.

# Issuer Name

The **issuer** field provides a globally unique identifier of the CA signing the CRL. The issuer name is an X.500 distinguished name. CRL issuer names with empty sequences are not supported by implementations conforming to this interoperability specification.

# This Update

The **thisUpdate** field indicates the date of the CRL. This field may be represented as **UTCTime** or **GeneralizedTime**. For this specification, **thisUpdate** follow the rules for the certificate **validity** field (see sec. 3.1.1 above).

# Next Update

The **nextUpdate** field indicates the date by which the next CRL will be issued. The next CRL could be issued before the indicated date, but it will not be issued any later than the indicated date. This field may be represented as **UTCTime** or **GeneralizedTime**. For this specification, **nextUpdate** shall follow the rules for the certificate **validity** field (see sec. 3.1.1 above).

# **Revoked Certificates**

The **revokedCertificates** field is a list of the certificates that have been revoked. Each revoked certificate listed contains:

- the certificate serial number, stated in the **userCertificate** field. This element contains the value of **serialNumber** of the revoked certificate. This must be used in conjunction with the name of the issuing CA to identify an unexpired certificate that has been revoked.
- the **revocationDate** field that contains the date of the revocation. The value included in this field shall follow the rules for the certificate validity field (see sec. 3.1.1 above).
- optional CRL entry extensions, that are specified in section 3.2.3 below. The CRL entry extensions may give the reason that the certificate was revoked, state the date that the invalidity is believed to have occurred, and may state the name of the CA that issued the revoked certificate, which may be a different CA from the CA issuing the CRL. Note that the CA that issued the CRL is assumed to be the CA that issued the revoked certificate unless the **certificateIssuer** CRL entry extension is included.

# 3.2.2 CRL Extensions

The extensions defined by ISO/ITU for X.509 v2 CRLs provide methods for associating additional information with entire CRLs. Each CRL extension may be designated as critical or noncritical. A CRL validation shall fail if a client encounters a critical extension that it cannot process.

This section describes CRL extensions that shall be supported. A CRL extension is supported when: the CA is able to generate the extensions in a CRL and the clients are able to process the extension.

# Authority Key Identifier

The **authorityKeyldentifier** is a noncritical CRL extension that identifies the CA's key used to sign the CRL. This extension is useful when a CA uses more than one key; it allows distinct keys to be differentiated (e.g., as key updating occurs). The identification can be based on either the key identifier or on the issuer name and serial number. The key identifier method shall be used, and the **keyldentifier** shall be generated for all CRLs. This extension is useful where an issuer has multiple signing keys (either due to multiple concurrent key pairs or due to changeover). This extension shall be included in all CRLs, and clients shall be able to find and validate CRL certification paths where the issuing CA has multiple signing keys. Clients shall be able to process either the key identifier or the certificate issuer plus serial number form of **authorityKeyldentifier** if they use this extension to find certification paths.

# **Issuer Alternative Name**

The **issuerAltName** is a noncritical CRL extension that contains one or more alternative CA names. Whenever such alternative names are present in a CRL, they shall be placed in the issuer alternative name field. Implementations which recognize this extension need not be able to process all the alternative name formats. Unrecognized alternative name formats may be ignored by an implementation. CAs shall be capable of generating this extension in CRLs, however clients are not required to process it.

# CRL Number

The **cRLNumber** field is a noncritical CRL extension which conveys a monotonically increasing sequence number for each CRL issued by a given CA through a specific CA directory entry or CRL distribution point. This extension can be used to alert certificate users to unscheduled issuance of full CRLs, or to easily determine when a particular CRL supersedes another CRL. This extension shall be included in CRLs.

# **Issuing Distribution Point**

The **issuingDistributionPoint** field is a critical CRL extension that identifies the CRL distribution point for this particular CRL. A distribution point is a directory entry that may be used to retrieve a CRL, and that may differ from the directory entry of the issuing CA. The CRL is signed by the CA's key. CRL distribution points do not have their own key pairs.

In addition, the **issuingDistributionPoint** field specifies CRLs that may contain only end entity certificates, or only CA certificates, or only certificates that have been revoked for a particular reason. Finally, this extension can identify an "indirect CRL," that is a CRL that is issued by a different CA than the CA(s) that issued the revoked certificates. It contains the following components:

• **distributionPoint**, which gives the distribution point name. If used, **distributionPoint** shall be an X.500 distinguished name;

- **onlyContainsUserCerts**, a Boolean value that indicates that the CRL contains only end entity certificates;
- **onlyContainsCACerts**, a Boolean value that indicates that the CRL contains only CA certificates;
- **onlySomeReasons**, a **ReasonsFlag** bit string that indicates the reasons for which certificates are listed in the CRL. Only the following reason flags shall be included in CRLs:
  - **keyCompromise** shall be used to indicate compromise or suspected compromise;
  - cACompromise shall be used to indicate that the certificate has been revoked because of a CA key compromise. It shall only be used to revoke CA certificates;
  - affiliationChanged shall be used to indicate that the certificate was revoked because of a change of affiliation of the certificate subject;
  - **superseded** shall be used to indicate that the certificate has been superseded;
  - certificateHold shall be used to indicate that the certificate's status is questionable and may be revoked;
  - cessationOfOperation shall be used to indicate that the certificate is no longer needed for the purpose for which it was issued, but there is no reason to suspect that the private key has been compromised.
- indirectCRL, a Boolean value that indicates that this is an indirect CRL.

Clients shall be able to process this field.

# Delta CRL Indicator

The **deltaCRLIndicator** is a critical CRL extension that identifies a delta-CRL. The use of delta-CRLs can significantly improve processing time for applications which store revocation information in a format other than the CRL structure. This allows changes to be added to the local database while ignoring unchanged information that is already in the local database.

The value of **BaseCRLNumber** identifies the CRL number of the base CRL that was used as the starting point in the generation of this delta-CRL. The delta-CRL contains the changes between the base CRL and the current CRL. It is the decision of a CA as to whether to provide delta-CRLs.

A client constructing a CRL from a locally held CRL and a delta-CRL shall consider the constructed CRL incomplete and unusable if the CRL number of the locally held CRL is less than the BaseCRLNumber in the delta-CRL.<sup>21</sup> If the delta CRL contains a CRL number extension, the CRL number of the constructed CRL is the value of the CRL number extension. Otherwise, the CRL number associated with the locally held CRL is maintained. Support of delta-CRLs by clients and CAs is optional.

# Summary of CRL Extension Use

Table 3-3 summarizes the standardized CRL extensions, while Table 3-4 summarizes the use of

<sup>&</sup>lt;sup>21</sup> Note that use of delta CRLs imposes an additional security requirement on clients; they must be capable of securely maintaining the composite CRL.

Extension	Use	Critical
authorityKeyIdentifier	identifies the CA key used to sign CRL.	No
Keyldentifier	unique key identifier; alternative to <b>certIssuer</b> &	
	authorityCertSerialNumber	
CertIssuer	name of CA's cert. issuer	
authorityCertSerialNumber	used with <b>certIssuer</b> ; combination must be unique	
IssuerAltName	alternate name of CRL issuer	No*
CRLNumber	sequence number for CRL	No
IssuingDistributionPoint	name of CRL distribution point; also gives reasons	Yes
	for revocations contained in CRL.	
DeltaCRLIndicator	indicates delta CRL (lists certificates. revoked since	Yes
	last full CRL) & gives sequence number	

#### Table 3-3 Summary of CRL Extensions

NOTES:

Standard allows either critical or noncritical. Indication is for use in interoperable implementations.

the standardized CRL extensions for the MISPC.

# **3.2.3 CRL Entry Extensions**

The CRL entry extensions defined for X.509 v2 CRLs provide methods for associating additional information with CRL entries. Each extension in a CRL entry is designated as critical or noncritical. A CRL validation shall fail if it encounters a critical CRL entry extension which it does not know how to process. However, an unrecognized noncritical CRL entry extension may be ignored.

Table 3-4 Summary	of CRL Exte	ensions and their	use in the MISPC
Tuble 5 + Dummary	of CILL LAG	motono ana men	use in the must c

Extension	CRL	Clients
authorityKeyldentifier		
keyldentifier	included in all CRLs issued	optional - used to help find correct CA certificate to validate CRL (1)
certissuer	not generated	optional - issuer/serial number pair used to help find correct authority certificate to validate CRL (1)
certSerialNumber	not generated	
issuerAltName	supported	optional
cRLNumber	supported: included in all CRLs	optional
issuingDistributionPoint	supported	supported
deltaCRLIndicator	optional	optional

NOTES:

- For CRLs, "supported" means that the CA is capable of issuing CRLs that contain this extension.
- For Clients, "supported" means that the client is capable of processing this extension in CRLs.

(1) Clients shall be capable of finding the certificate corresponding to the private key used to sign a CRL, when the CA has multiple certificates, and the certificates are accessible in the appropriate directory, whether or not they use this extension to do so, and whether or not the CRL contains this extension.

# Reason Code

The **reasonCode** is a noncritical CRL entry extension that identifies the reason for the certificate revocation. CAs shall be capable of generating this extension in CRL entries. Processing of the **reasonCode** extension by clients is optional, that is clients shall not validate a certificate if any certificate in the certification path is listed in a current CRL, regardless of the **reasonCode**, and need not provide operator information about the reason for failure. The following enumerated **reasonCode** values are defined:

- **unspecified**; this value shall not be used;
- **keyCompromise** indicates compromise or suspected compromise;
- **cACompromise** indicates that the certificate has been revoked because of a CA key compromise. It shall only be used to revoke CA certificates;
- **affiliationChanged** indicates that the certificate was revoked because of a change of affiliation of the certificate subject;
- superseded indicates that the certificate has been replaced by a more recent certificate;
- **cessationOfOperation** indicates that the certificate is no longer needed for the purpose for which it was issued, but there is no reason to suspect that the private key has been compromised.
- **certificateHold** indicates that the certificate shall not be used at this time. When clients process a certificate that is listed in a CRL with a reasonCode of **certificateHold**, they shall fail to validate the certification path.
- **removeFromCRL**, which is used only with delta-CRLs and indicates that an existing CRL entry should be removed.

# Expiration Date

The **expirationDate** is a noncritical CRL entry extension that indicates the expiration of a hold entry in a CRL. This extension shall not be used in CRLs or by clients.

#### Instruction Code

The **instructionCode** is a noncritical CRL entry extension that provides a registered instruction identifier which indicates the action to be taken after encountering a certificate that has been placed on hold. This extension shall not be used in CRLs.

# Invalidity Date

The **invalidityDate** is a noncritical CRL entry extension that provides the date on which it is known or suspected that the private key was compromised or that the certificate otherwise became invalid. This date may be earlier than the revocation date in the CRL entry. The revocation date in the CRL entry specifies the date that the CA revoked the certificate. Whenever this information is available, CAs are encouraged to share it with CRL users. CAs shall be capable of generating this extension in CRLs. This value is represented as **GeneralizedTime**.

#### Certificate Issuer

The certificateIssuer CRL entry extension is used with an indirect CRL (a CRL that has the

**indirectCRL** indicator set in its **issuingDistributionPoint** extension). If this extension is not present in the first entry of an indirect CRL, the certificate issuer defaults to the CRL issuer. In subsequent entries in an indirect CRL, when the **certificateIssuer** extension is not present, the certificate issuer is the same as the issuer of the preceding CRL entry.

## Summary of CRL Entry Extension Use

Table 3-5 summarizes the CRL entry extensions while Table 3-6 summarizes the use of CRL entry extensions for the MISPC.

Extension	Use	Critical
reasonCode	identifies the reason for the revocation of this	No
	certificate	
instructionCode	used with certificateHold reasonCode;	No
	indicates action to be taken when encountering a	
	held certificate	
invalidityDate	date certificate became invalid	No
certificatelssuer	Issuer of revoked certificate in an indirect CRL	Yes

# Table 3-5 Summary of CRL Entry Extensions

# Table 3-6 Summary of CRL Entry Extensions Use for MISPC

Extension	CRL	Clients
reasonCode	supported; included for all	optional - may be used to provide
	entries	information about validation
		failure
instructionCode	not used	optional
invalidityDate	supported	optional - may be used to provide
		information about validation
		failure
certificatelssuer	optional	optional - necessary to support
		processing of indirect CRLs

NOTES

For CRLs, "supported" means that CAs are capable of issuing CRLs that contain this CRL entry extension. For clients, "supported" means that the client is capable of processing this entry extension in CRLs.

#### 3.3 Certification Path Validation

The procedure specified in section 6 of the [RFC2459], *Certification Path Validation*, shall be implemented by clients.

#### 3.4 Transaction Message Formats

This section presents a set of message formats to support the minimal set of PKI transactions. Systems that implement these transactions shall support these message formats, generating and

recognizing them as appropriate. The message formats are specified in ASN.1; messages shall be encoded and transmitted using the Distinguished Encoding Rules (DER).

These message formats are used to implement transactions described in section 3.5.

# 3.4.1 Overall PKI Message Components

#### PKI Message

Each message has four components

```
PKIMessage ::= SEQUENCE {

header PKIHeader,

body PKIBody,

protection [0] PKIProtection OPTIONAL,

extraCerts [1] SEQUENCE OF Certificate OPTIONAL

}
```

The **extraCerts** field is not used within this specification.

#### PKI Message Header

All PKI messages require some header information for addressing and transaction identification. Some of this information will also be present in a transport specific envelope, however, if the PKI message is signed then this information is also protected (i.e., we make no assumption about secure transport).

The following data structure is used to contain this information:

#### PKIHeader ::= SEQUENCE {

pvno	INTEGER	{ ietf-version2 (1) },
sender	GeneralName,	identifies the sender
recipient	GeneralName,	<ul> <li>identifies the intended recipient</li> </ul>
messageTime	[0] GeneralizedTime	OPTIONAL,
-	on of this message (use	
	still be meaningful up	• •
• •	[1] Algorithmldentifier	•
•	or calculation of protec	tion bits
senderKID		OPTIONAL,
recipKID		OPTIONAL,
	fic keys used for protect	
transactionID	[4] OCTET STRING	•
		be the same in corresponding
• • •	se and confirmation me	
senderNonce		
-	[6] OCTET STRING	
•		on, senderNonce is inserted by the creator
	•	previously inserted in a related message by
	ipient of this message	
freeText	L	OPTIONAL,
•	-	ecific instructions (this field is intended for
human consump	,	
generalInfo	``	1MAX) OF InfoTypeAndValue OPTIONAL
this may be used	d to convey context-spe	ecific information

# -- (this field not primarily intended for human consumption) }

#### PKIFreeText ::= UTF8String

The **transactionID** field within the message header allows the recipient of a response message to correlate this with the request. In the case of an RA there may be many requests "outstanding" at a given moment. The value of this field should be unique from the sender's perspective in order to be useful.

The **messageTime** field indicates the time the message was generated. The value included in this field shall be expressed in Greenwich Mean Time (Zulu) and shall include seconds (i.e., times are **YYYYMMDDHHMMSSZ**), even where the number of seconds is zero. The **messageTime** values shall not include fractional seconds.

The **sender** and **recipient** fields within the message header are defined as **GeneralName**. Systems are required to support X.500 distinguished names and RFC 822 (Internet electronic mail) names.

The **freetext** field is defined as **PKIFreeText**, which is defined as the ASN.1 type **UTF8String**. **UTF8String** is a representation of the Universal Character Set whose encoding preserves full US-ASCII range. For this specification, **PKIFreeText** only includes characters in the US-ASCII range.

The **protectionAlg** is required for all signed messages, and for messages protected by a message authentication code. When a message is not protected (i.e., the **protection** field does not appear in the **PKIMessage** containing this **PKIHeader**), **protectionAlg** must be omitted.

The **senderNonce** and **recipNonce** are not required by client and RA implementations; CAs that receive **PKIMessages** with a **senderNonce** must, at a minimum, be able to return that value as the **recipNonce** in the following message. The **senderKID**, **recipKID**, and **generalInfo** fields are not required to implement this specification.

#### PKI Message Body

PKIBody ::= CHOICE {				
message	e-spe	cific body elements		
ir	[0]	CertReqMessages,		
ip	[1]	CertRepMessage,		
cr	[2]	CertReqMessages,		
ср	[3]	CertRepMessage,		
p10cr	[4]	PKCS10CertReqMessages,		
popdecc	[5]	POPODecKeyChallContent,		
popdecr	[6]	POPODecKeyRespContent,		
rr	[11]	RevReqContent,		
rp	[12]	RevRepContent,		
conf	[19]	PKIConfirmContent,		
}				

Additional message-specific body elements are defined by [RFC2510], [RFC2511]. The additional elements are not required to implement this specification, so they were omitted for clarity. The complete list of message-specific body elements appears in Appendix C.

Other sections of this document refer to InitReq, InitRep, CertReq, CertRep, RevReq, and RevRep messages. These terms refer to PKIMessages with body elements ir, ip, cr, cp, rr, and rp, respectively. A PKCS #10 request refers to a message with a p10cr body element. A confirmation message will have body element conf.

## PKI Message Protection

All PKI messages will be protected for integrity using the following structure:<sup>22</sup>

## PKIProtection ::= BIT STRING

The input to the calculation of the **PKIProtection** is the DER encoding of the following data structure:

ProtectedPart ::= SEQUENCE { header PKIHeader, body PKIBody }

In most cases, the **PKIProtection** field will contain a digital signature and the **protectionAlg** field in the **PKIHeader** will contain an **AlgorithmIdentifier** specifying the digital signature algorithm (e.g., **dsaWithSha-1**) used to protect the message.

In some cases, such as key update, it may be necessary to attach multiple signatures. In this case, signed messages are nested - each signed message becomes a **PKIBody** element **nested**; the next signature is applied to this message. This process is repeated until all signatures have been applied.

Where symmetric techniques are needed for message authentication, the algorithm id shall be one of those identified in section 3.1.2.5 and the **PKIProtection** value shall contain the message authentication code using the DER encoded header and body as input (and the shared secret as the key.) The **PKIHeader** will contain an **AlgorithmIdentifier** specifying a message authentication code algorithm (e.g., SHA1-HMAC).

# 3.4.2 Common Data Structures

The following data types are common to several message formats.

# Certificate Templates

In various PKI management messages, the originator may provide certain values to identify an existing certificate or request certain values be used in the generation of a certificate. The **CertTemplate** structure allows entities to indicate those values. **CertTemplate** includes all the same information as a certificate.

#### CertTemplate ::= SEQUENCE { version [0] Version OPTIONAL, -- used to ask for a particular syntax version serial [1] INTEGER OPTIONAL,

 $<sup>^{22}</sup>$  There is one exception to this rule: an end entity that is not a current certificate holder will not protect the initial message in an RA-Generated request (sec. 3.5.1). This message is delivered to the RA out-of-band; no protection mechanism is required. In this case, **protectionAlg** is omitted from the header.

		to ask for a particular serial ne behalf of a previous certificate	•		
۰ ۲	signingAlg	[2] AlgorithmIdentifier	OPTIONAL,		
	subject	[3] Name	OPTIONAL,		
	alidity	[4] OptionalValidity	OPTIONAL,		
	ssuer	[5] Name	OPTIONAL,		
r	oublicKey	[6] SubjectPublicKeyInfo	OPTIONAL,		
•	ssuerUID	[7] Uniqueldentifier	<b>OPTIONAL</b> , not supported		
s	subjectUID	[8] UniqueIdentifier	<b>OPTIONAL</b> , not supported		
e	extensions	[9] Extensions	OPTIONAL,		
contains the extensions which the requester					
	would	d like in the cert.			
}					
-	alValidity ::= S	•			
		[0] Time OPTIONAL,			
r	notAfter	[1] Time OPTIONAL			
}					

#### CertTemplates ::= SEQUENCE OF CertTemplate

If it appears, the validity field contains the requested issuance date (in the notBefore field) and expiration date (notAfter) for the requested certificate. The values in the CertTemplate validity field shall be interpreted as specified for the certificate validity field (see sec. 3.1.1).

Proving Possession of a New Signature Key

Conforming CAs verify that the prospective subject of a certificate request holds the private key corresponding to the public key provided in a certificate request. This is performed with the signature field in **ProofOfPossession**, which is a **POPOSigningKey** structure. The **POPOSigningKey** structure includes input data, an algorithm identifier, and a signature.

The input data is generated as the DER-encoded **popInput**, and is generated from the data in the certificate request. In general, it contains the subject name and the public key from the **CertTemplate**. When a new subject's request is authenticated through use of a public key mac value, or the RA modifies the subject name, the **popInput** must be constructed from optional data in the **POPOSigningKey** structure and the public key from the **CertTemplate**. This permits an RA to pass proof-of-possession to the CA despite changing the CertTemplate.

#### ProofOfPossession ::= CHOICE {

```
raVerified
                     [0] NULL,
  -- used if the RA has already verified that the requester is in
  -- possession of the private key
                     [1] POPOSigningKey,
  signature
  keyEncipherment [2] POPOPrivKey,
  keyAgreement
                     [3] POPOPrivKey }
POPOSigningKey ::= SEQUENCE {
```

[0] POPOSKInput OPTIONAL, poposkInput algorithmIdentifier AlgorithmIdentifier, signature BIT STRING }

-- The signature (using "algorithmIdentifier") is on the

- -- DER-encoded value of popInput. NOTE: If poposkInput is present
- -- in the pop field, popInput is constructed
- -- with otherinput. If poposkInput is not present, subject is the name
- -- from CertTemplate. Note that the encoding of PopInput is intentionally ambiguous.

PoposkInput ::= CHOICE { Subject name, Sender [0] generalName, publicKeyMAC [1] PKMACValue }

-- The pop is calculated upon the structure popInput, which is defined

```
-- as follows:
```

-- If poposkInput is present

- -- in the pop field, popInput is constructed
- -- with otherinput. If poposkInput is not present, subject is the name
- -- from CertTemplate. Note that the encoding of PopInput is
- -- intentionally ambiguous.

#### <u>CertReqMessage</u>

The **CertReqMsg** is the basic structure for certificate requests. **CertReqMsg** is a **SEQUENCE** of three fields, a **certReq**; **pop**; and, optionally, **regInfo**. **certReq** is of type **CertRequest**, **pop** includes information demonstrating proof-of-possession of the private key; and **regInfo** contains information specific to the registration procedures for this PKI. The **regInfo** field is not required to implement this specification.

```
CertReqMsg ::= SEQUENCE {
certReq CertRequest,
pop ProofOfPossession OPTIONAL,
-- content depends upon key type
regInfo SEQUENCE SIZE(1..MAX) of AttributeTypeAndValue OPTIONAL }
```

#### **CertRequest**

The **CertRequest** syntax consists of a request identifier, a template of certificate content, and an optional sequence of control information.

```
CertRequest ::= SEQUENCE {
certReqId INTEGER, -- ID for matching request and reply
certTemplate CertTemplate, -- Selected fields of cert to be issued
controls Controls OPTIONAL } -- Attributes affecting issuance
```

certReqId is an INTEGER, certTemplate is a CertTemplate (see above), and controls are defined below.

# <u>CertReqMessages</u>

CertReqMessages is a sequence of one or more CertReqMessage structures.

## CertReqMessages ::= SEQUENCE SIZE (1..MAX) OF CertReqMessage

For this specification, **CertReqMessages** may be assumed to be a **SEQUENCE** of exactly one **CertReqMessage**. That is, **MAX** may be defined as one.

## **Controls**

The generator of a **CertRequest** may include one or more control values pertaining to the processing of the request.

## Controls ::= SEQUENCE SIZE (1..MAX) OF AttributeTypeAndValue

The following controls were defined in CMP (it is recognized that this list may expand over time): regToken; authenticator; pkiPublicationInfo; pkiArchiveOptions; oldCertld; protocolEncrKey. This specification requires support for protocolEncrKey is needed when issuing key management certificates; these transactions are optional to implement for this specification. OIDs for these controls are provided below. The remaining controls (oldCertld, regToken, authenticator, pkiPublicationInfo, and pkiArchiveOptions) are not required for this specification and are not described here.

id-pkix OBJECT IDENTIFIER ::= { iso(1) identified-organization(3) dod(6) internet(1) security(5) mechanisms(5) pkix(7) }

-- arc for Internet X.509 PKI protocols and their components id-pkip OBJECT IDENTIFIER :: { id-pkix pkip(5) }

-- Registration Controls in CRMF id-regCtrl OBJECT IDENTIFIER ::= { id-pkip regCtrl(1) } id-regCtrl-oldCertId OBJECT IDENTIFIER ::= { id-regCtrl 5 } id-regCtrl-protocolEncrKey OBJECT IDENTIFIER ::= { id-regCtrl 6 }

# Protocol Encryption Key Control

If present, the protocol encryption key control specifies a key the CA is to use in encrypting a response to **CertReqMessages**.

This control can be used when a CA has information to send to the subscriber that needs to be encrypted. Such information includes a private key generated by the CA for use by the subscriber.

The protocolEncrKey control is indicated by the OID id-reg-protocolEncrKey in the attribute type; the syntax of its value is SubjectPublicKeyInfo.

Status codes for PKI messages

All response messages will include some status information. The following values are defined:

```
-- request granted, with modifications; the requester
       -- is responsible for ascertaining the differences
rejection
                       (2),
        -- request rejected
waiting
                       (3),
        -- the request has been received but has not been processed,
       -- an additional response will follow after processing
revocationWarning
                       (4),
        -- this message contains a warning that a revocation has
        -- been requested and is under consideration
revocationNotification (5),
       -- notification that a revocation has occurred
keyUpdateWarning (6)
       --
```

This specification does not use the status code keyUpdateWarning.

#### Failure Information

}

Responders use the following syntax to provide more information about failure cases.

PKIFailureInfo ::= BIT STRI	NG {	since we can fail in more than
one way	' <b>!</b>	
badAlg	(0),	unrecognized or unsupported algorithm identifier
badMessageCheck	(1),	<ul> <li> integrity check failed (e.g., signature did not verify)</li> </ul>
badRequest	(2),	transaction not permitted or supported
badTime	(3),	messageTime field was not sufficiently close
		to the system time, as defined by local policy
badCertId	(4),	no certificate could be identified matching the
		provided criteria
badDataFormat	(5),	the data submitted has the wrong format
wrongAuthority	(6),	the authority indicated in the request is different from the
		one creating the response token
incorrectData	(7),	the requester's data is incorrect (used for notary services)
missingTimeStamp	(8),	when the timestamp is missing but is required by policy
badPoP	(9)	proof of possession field did not verify
need more failure i	nform	ation
}		
-		
PKIStatusInfo ··- SEOUENC	Έſ	

PKIStatusInfo ::= SEQUEN	CE {	
status	PKIStatus,	
statusString	PKIFreeText	OPTIONAL,
failInfo	PKIFailureInfo	OPTIONAL
•		

#### }

#### Protocol Confirmation

Confirmation messages shall carry all the required information in the **PKIHeader**. As a result, this data structure has a NULL content.

#### PKIConfirmContent ::= NULL

# Certificate Identification

In order to identify particular certificates the **CertId** structure is used.

Proof Of Possession for Key Management Keys

Certificate Request transactions for Key Management certificates may include challenge and response messages to verify the requester's possession of the private key. The following structures form the **PKIBody** of these messages.

POPODecKeyChallContent ::= SEQUENCE OF Challenge -- One Challenge per encryption key certification request (in the -- same order as these requests appear in FullCertTemplates). Challenge ::= SEQUENCE { AlgorithmIdentifier OPTIONAL, owf -- must be present in the first Challenge; may be omitted in any -- subsequent Challenge in POPODecKeyChallContent (if omitted, -- then the owf used in the immediately preceding Challenge is -- to be used). witness **OCTET STRING**, -- the result of applying the one-way function (owf) to a -- randomly-generated INTEGER, A. [Note that a different -- INTEGER must be used for each Challenge.] challenge **OCTET STRING** -- the encryption (under the public key for which the cert. -- request is being made) of Rand, where Rand is specified as -- Rand ::= SEQUENCE { int INTEGER, --- the randomly-generated INTEGER A (above) -sender GeneralName

-- - the sender's name (as included in PKIHeader)

```
-- }
}
```

POPODecKeyRespContent ::= SEQUENCE OF INTEGER

```
-- One INTEGER per encryption key certification request (in the
```

-- same order as these requests appear in FullCertTemplates). The

- -- retrieved INTEGER A (above) is returned to the sender of the
- -- corresponding Challenge.

# Centrally Generated Keys

When a CA generates an encryption key pair on behalf of a certificate holder, it uses the **CertifiedKeyPair** structure to return the certificate.

```
CertifiedKeyPair ::= SEQUENCE {
certificate [0] Certificate OPTIONAL,
```

```
encryptedCert [1] EncryptedValue OPTIONAL,
privateKey [2] EncryptedValue OPTIONAL,
publicationInfo [3] PKIPublicationInfo OPTIONAL
```

}

EncryptedValue ::= SEQUENCE { intendedAlg [0] AlgorithmIdentifier OPTIONAL. -- the intended algorithm for which the value will be used [1] AlgorithmIdentifier OPTIONAL, symmAlg -- the symmetric algorithm used to encrypt the value encSymmKey [2] BIT STRING OPTIONAL, -- the (encrypted) symmetric key used to encrypt the value [3] AlgorithmIdentifier OPTIONAL, kevAla -- algorithm used to encrypt the symmetric key valueHint [4] OCTET STRING OPTIONAL, -- a brief description or identifier of the encValue content -- (may be meaningful only to the sending entity, and used only -- if EncryptedValue might be re-examined by the sending entity -- in the future) encValue BIT STRING } -- the encrypted value itself

## Out-of-band Information

To convey a CA's public key out of band, **OOBCert** structure is used. **OOBCert** is simply the CA's certificate.

#### OOBCert ::= Certificate

#### 3.4.3 Operation-Specific Data Structures

#### Registration/Certification Request

Registration/Certification request (**cr**) messages and initialization requests (**ir**) messages contain a **CertReqMessages** data structure which specifies values for one or more requested certificates.

CertReqMessages is a sequence of one or more CertReqMsg structures.

#### CertReqMessages ::= SEQUENCE SIZE (1..MAX) OF CertReqMsg

With one exception, transactions defined in this specification assume **CertReqMessages** is a **SEQUENCE** of exactly one **CertReqMsg**. The optional transaction requesting a signature certificate and a key management certificate simultaneously requires that **CertReqMessages** be a **SEQUENCE** of size two.

#### Registration/Certification Response

The registration response message (**cp**) and initialization response message (**ip**) contain a **CertRepMessage** structure which is an optional CA public key and a **response**. With one exception, transactions defined in this specification assume the response is a sequence of exactly one **CertResponse**. The exception is a transaction combining a signature certificate and a key management certificate request; the response to this transaction may be a combined response message, with a sequence of two **CertResponse** fields.

The **CertResponse** includes a request id, status information and optionally a **CertifiedKeyPair**. The **CertifiedKeyPair** is a sequence of four optional fields: a certificate, an encrypted certificate, an encrypted private key, and publication information. In this specification, the certificate field will always appear in **CertifiedKeyPair**. **PrivateKey** is used only for centrally generated key management keys; the other fields are never used in this specification.

```
CertRepMessage ::= SEQUENCE {
     caPub
                   [1] Certificate OPTIONAL.
                   SEQUENCE OF CertResponse
     response
}
CertResponse ::= SEQUENCE {
    certRegId
                          INTEGER,
                                        -- to match this response with corresponding request
     certRepStatus
                          PKIStatusInfo.
     certifiedKeyPair
                          CertifiedKeyPair
                                                              -- present if status is granted
                                                OPTIONAL.
                                                              -- or grantedWithMods
    rspinfo
                          OCTET STRING
                                             OPTIONAL
 -- analogous to the id-regInfo-asciiPairs OCTET STRING defined
 -- for regInfo in CertReqMsg [CRMF]
```

}

If certRepStatus contains a failInfo field, the CertResponse shall not include a certifiedKeyPair and the value in the certRepStatus field shall be rejection on the value of status. For the status value waiting none of the optional fields may be present. The status values revocationWarning and revocationNotification should not appear in this message.

The **caPub** and **rspInfo** fields are not required, and may be ignored if present. This interoperability specification does not use the **encryptedCert** and **publicationInfo** fields in **CertifiedKeyPair**.

# Revocation Request Content

When requesting revocation of a certificate the following data structure is used. The name of the requester is present in the **PKIHeader** structure.

#### RevReqContent ::= SEQUENCE OF RevDetails

```
RevDetails ::= SEQUENCE OF {
	certDetails CertTemplate,
	-- allows requester to specify as much as they can about
	-- the cert. for which revocation is requested
	-- (e.g. for case serialNumber not available)
	revocationReason ReasonFlags,
	-- from the DAM, so that CA knows which Dist. point to use
	badSinceDate GeneralizedTime OPTIONAL,
	-- indicates best knowledge of sender
	crlEntryDetails Extensions}
	-- requested crlEntryExtensions
```

ReasonFlags are defined in Appendix B, but are reproduced here for clarity.

```
ReasonFlags ::= BIT STRING {
```

unused	(0),
keyCompromise	(1),
caCompromise	(2),
affiliationChanged	(3),
superseded	(4),
cessationOfOperation	(5),
certificateHold	(6),
removeFromCRL	(8) }

The same information represented by **badSinceDate** and **revocationReason** can be represented in **crlEntryDetails** using the standard X.509 CRL entry extensions **invalidityDate** and **reasonCode**. The preferred location is **crlEntryDetails**, but this specification recommends that CAs process the information in either location. Where there is a conflict, the earlier date should be used. When revocation reasons don't match, the CA should logically OR the values.

# Revocation Response Content

When responding to a revocation request the following data structure is used. If produced this is sent to the requester of the revocation.

```
RevRepContent ::= SEQUENCE {

status PKIStatusInfo,

revCerts [0] SEQUENCE OF CertId OPTIONAL,

-- identifies the cert for which revocation

-- was requested

crIs [1] SEQUENCE OF CertificateList OPTIONAL}

-- the resulting CRL
```

For the purposes of this specification, **revCerts** shall be a **SEQUENCE** of one or more **CertId** fields, and the **crIs** field does not appear.

# PKCS #10 Certification Request

This alternative certification request syntax is defined in [PKCS#10]. It is reproduced here for clarity.

PKCS1	0CertReqMessages ::=	SEQUENCE {	
	certificationRequestIn	fo CertificationRequestInfo	
	signatureAlgorithm	SignatureAlgorithmIdentifier,	
	signature	Signature	
}			
Signat	uro AlgorithmIdontifior .	- AlgerithmIdentifier	
SignatureAlgorithmIdentifier ::= AlgorithmIdentifier			
Signature ::= BIT STRING			
0			
CertificationRequestInfo::= SEQUENCE {			
	version Version		
	subject	Name,	
		SubjectPublicKeyInfo,	
	attributes	[0] IMPLICIT Attributes	
}			

## Version ::= INTEGER

## Attributes ::= SET OF Attribute

Attributes are specified in [PKCS#9]. Support for attributes is optional for conforming implementations. If present, they may be ignored.

# 3.5 PKI Transactions

This section describes PKI specific functions to request, renew, and revoke certificates. This section also provides a brief description of transactions for accessing the directory service.

Compliant CAs shall implement all of the transactions identified in this section. Compliant RAs shall implement the RA-Generated Registration (sec. 3.5.1) and Revocation (sec. 3.5.6) transactions. Compliant certificate holders shall implement the Revocation (sec. 3.5.6) and RA-Generated Registration (sec. 3.5.1) transactions. Self Registration (secs. 3.5.3 and 3.5.5) and Certificate Renewal (sec. 3.5.2) transactions are optional for certificate holders.

Where the product collocates the CA and RA, and does not support remote RAs, messages between the CA and RA are omitted. In this case, the functionality of the RA-generated revocation request must be supported, but there are no protocol or interface requirements. All other transactions must be supported, accepting requests from clients and generating appropriate responses.

# 3.5.1 RA-Generated Registration Requests

An RA may request that a CA issue a certificate for an end entity. This transaction is performed in five steps. In the first step, the end entity provides a public key to the RA in a signed message in an out-of-band transaction (e.g., by physically presenting a diskette). In the second step, the RA requests a certificate from the CA in a signed message. The CA replies to the RA with a signed message containing either a certificate or an error message. The RA provides the end entity with the CA's public key out-of-band. The end entity may receive the certificate from the RA out-of-band, or from the CA electronically.

The final two messages in the transaction are optional; upon request from the CA, the RA will generate a confirmation message. Where the CA receives a confirmation message, it shall respond in kind. For these specifications, messages are not generated unless the CA generates a certificate.

# Certificate Request from an End Entity to the RA

The end entity creates a **PKIMessage** with **PKIBody** element **cr**. The **PKIHeader** includes the following information:

- pvno is one;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the distinguished name of the end entity, or null;
- recipient is the distinguished name of the RA, or null; and
- if the end entity is a current certificate holder, **protectionAlg** is the algorithm identifier for

the signature algorithm used to protect the message; otherwise **protectionAlg** is omitted.

The message body is **CertReqMessages**, which is a sequence of one or more **CertReqMessage** fields. For this transaction, **CertReqMessages** is a sequence of one **CertReqMessage**. The **CertReqMessage** will include the following information:

- **certReq** contains the information that the requester would like included in the certificate; and
- **pop** provides proof of possession of the private key for the new certificate.

The **pop** field shall be generated using the private key corresponding to the public key in the **publicKey** field.

The certReq is a CertRequest, which is a sequence of a certReqID, a CertTemplate, and controls. For this transaction:

- **certReqID** is any integer;
- **certTemplate** is a **CertTemplate** including, at a minimum, the **publicKey** field which provides the public key for the new certificate.

The following information may be included in the **CertTemplate**:

- **signingAlg** specifies the preferred signature algorithm; and
- **subject** specifies the distinguished name for the prospective certificate holder.

If the end entity is a current certificate holder, the **PKIProtection** field contains the end entity's signature, calculated on the DER encoded sequence of the header and body with private key material corresponding to the current certificate. If the end entity is not a current certificate holder, the **PKIProtection** field shall be omitted.

# Certificate Request from RA to CA

The RA creates a **PKIMessage** with **PKIBody** element **cr**. The **PKIHeader** includes the following information:

- **pvno** is one;
- transactionID is an integer unique to this transaction for this RA;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the distinguished name of the RA;
- **recipient** is the distinguished name of the CA; and
- **protectionAlg** is the algorithm identifier for the signature algorithm used to protect the message.

The message body is **CertReqMessages**, which is a sequence of one or more **CertReqMessage** fields. For this transaction, **CertReqMessages** is a sequence of one **CertReqMessage**. The **CertReqMessage** will include the following information:

- **certReq** contains the information that the requester would like included in the certificate; and
- **pop** provides proof of possession of the private key for the new certificate.

If the RA did not modify the subject name, the **pop** field shall be the same as provided by the

requester. If the RA modified, the subject name, **popoSKInput** shall be present and will contain the original subject name.

The certReq is a CertRequest, which is a sequence of a certReqID, a CertTemplate, and controls. For this transaction certReqID may be any integer.

The **CertTemplate** will include the following information:

- version is v3 (2);
- **publicKey** provides the public key for the new certificate; and
- **extensions** specifies, at a minimum, the certificate policy OID to be associated with the certificate.

The following information may be included in the **CertTemplate**:

- **signingAlg** specifies the preferred signature algorithm; and
- **subject** specifies the distinguished name for the prospective certificate holder.

The request shall not include the following information:

- issuerUID; and
- subjectUID.

The **PKIProtection** field contains the RA's signature, calculated on the DER encoded sequence of the header and body.

#### Certificate Response from CA to RA

The CA will return a **PKIMessage** with **PKIBody** element **cp** to the RA.

The **PKIHeader** includes the following information:

- pvno is one;
- transactionID is the same as the transactionID field in the cr message;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the distinguished name of the CA;
- **recipient** is the distinguished name of the RA; and
- **protectionAlg** is the algorithm identifier for the signature algorithm used to protect the message.

If a **senderNonce** was supplied in the certificate request message, the header of the response shall include it as **recipNonce**.

The **PKIBody** element **cp** is of type **CertRepMessage**. For this transaction, **CertRepMessage** will contain a single **response** field. The **response** field is a **SEQUENCE** of a **certReqId**, **status** and **certifiedKeyPair**. If the CA issued a certificate, the body will contain the following information:

- **certReqId** will match the **certReqId** in the request;
- status will be granted or grantedWithMods; and
- The certifiedKeyPair sequence will contain one field, certificate, which will contain the X.509 version 3 certificate.

The certificate must meet the following properties:

- **version** number shall be v3 (2);
- The **publicKey** field shall be the same as in the certificate request;
- the subject distinguished name shall be the same as in the certificate request;
- the issuer name shall be the CA's distinguished name;
- if **notBefore** was present in the certificate request, the certificate shall be valid from the issuance date or the **notBefore** date, whichever is later; and
- if **notAfter** was present in the certificate request, the certificate shall expire on or before that date.

The certificate shall contain the following **extensions**:

- a subjectKeyldentifier field;
- at least one certificate policy OID in the **certificatePolicies** field; and
- an authority key identifier including a **Keyldentifier** field.

If a specific key identifier was specified in the certificate request message, the certificate shall contain that key identifier as the **subjectKeyldentifier** field. If no key identifier was supplied, the CA shall use the 96-bit SHA-1 hash of the subject public key as the **keyldentifier** in the **subjectKeyldentifier** field. The hash shall be calculated over the value (excluding tag and length) of the subject public key field in the certificate.

The certificate shall include URLs in the **issuerAltName** extension and **distributionPoint** field of the **CRLDistributionPoints** extension if the issuer's certificates or CRLs are not available from a well known X.500 directory.

# The failinfo field may not be present if status is granted or grantedWithMods.

If the CA rejected the request, the body shall include the following information:

- status will be rejected; and
- failinfo will contain the appropriate failure codes:
  - badAlg indicates that the CA cannot validate the signature because the algorithm identifier is unrecognized or unsupported;
  - badMessageCheck indicates that the signature in the PKIProtection field was checked but did not match;
  - badPoP indicates that the signature in the popoSigningKey field was checked but did not match;
  - **badRequest** indicates that the responder does not permit or support the transaction;
  - badTime indicates that the messageTime field in the message header was not sufficiently close to the responder's system time.<sup>23</sup>

The certificate field may not be present if status is rejected.

The **PKIProtection** field contains the CA's signature, calculated on the DER encoded sequence of the header and body.

<sup>&</sup>lt;sup>23</sup> This error code assumes a locally defined window of time for responding to a PKI message. The MISPC does not require such a policy, but defines this error code to support such policies.

# Confirmation Messages

Upon receipt of the **cp**, the RA shall generate a **PKIConfirm** message. **PKIHeader** data shall be identical to the certificate request, with the exception of **messageTime**.

The **PKIProtection** field contains the RA's signature, calculated on the DER encoded sequence of the header and body.

Upon receipt of the **PKIConfirm**, the CA shall generate a **PKIConfirm** message. **PKIHeader** data shall be identical to the **cr**, with the exception of **messageTime**.

The **PKIProtection** field contains the CA's signature, calculated on the DER encoded sequence of the header and body.

# 3.5.2 Certificate Renewal Request

An entity that is a current certificate holder may request issuance of a new certificate directly from the CA that issued the current certificate. The requesting entity creates a PKI **kr** (key update request) message requesting a certificate and includes proof of possession of the private key corresponding to the public key in the certificate request. The entity then signs the message with the private key corresponding to the entity's unexpired, unrevoked certificate.

If the CA's Certificate Practice Statement permits certificate renewal,<sup>24</sup> it will return a kp (key update response) message to the certificate holder. This message will contain the certificate or a reason code for the transaction failure.

The final two messages in the transaction are optional; upon request from the CA, the certificate requester will generate a confirmation message. Where the CA receives a confirmation message, it shall respond in kind. For these specifications, messages are not generated unless the CA generates a certificate.

# Certificate Renewal Request from Certificate Holder to CA

The certificate holder creates a key update request: a **PKIMessage** with **PKIBody** element **kr**. The **PKIHeader** includes the following information:

- **pvno** is one;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the distinguished name of the certificate holder;
- **recipient** is the distinguished name of the CA; and
- **protectionAlg** is the algorithm identifier for the signature algorithm used to protect the message.

The message body is **CertReqMessages**, which is a sequence of one or more **CertReqMessage** fields. For this transaction, **CertReqMessages** is a sequence of one **CertReqMessage**. The **CertReqMessage** will include the following information:

• **certReq** contains the information that the requester would like included in the certificate; and

<sup>&</sup>lt;sup>24</sup> Conforming CA *implementations* shall support certificate renewal. However, a particular CA may choose not to support this transaction as a matter of policy.

• **pop** provides proof of possession of the private key for the new certificate.

The **pop** field shall be generated using the private key corresponding to the public key in the **publicKey** field.

The certReq is a CertRequest, which is a sequence of a certReqID, a CertTemplate, and controls. For this transaction:

- **certReqID** is any integer; and
- certTemplate is a CertTemplate including the public key for the new certificate.

The **CertTemplate** will include the following information:

- version is v3 (2); and
- **publicKey** provides the public key for the new certificate.

The following information may be included in the **CertTemplate**:

- **signingAlg** specifies the preferred signature algorithm;
- **subject** specifies the distinguished name for the prospective certificate holder;

If **signingAlg** does not appear, the CA should sign with the algorithm corresponding to the entity's public key.

The request shall not include the following information:

- issuerUID; and
- subjectUID.

The **PKIProtection** field contains a signature generated using the private key associated with the current unexpired, unrevoked certificate and calculated upon the DER encoded sequence of the header and body.

# Certificate Renewal Response from CA to Certificate Holder

The CA will return a key update response (a **PKIMessage** with **PKIBody** element **kp**) message to the certificate holder.

The **PKIHeader** includes the following information:

- pvno is one;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the distinguished name of the CA;
- **recipient** is the distinguished name of the certificate holder and the sender of the **cr** message; and
- **protectionAlg** is the algorithm identifier for the signature algorithm used to protect the message.

If a **transactionID** was supplied in **cr** message, the header of the response will include the same **transactionID**. If a **senderNonce** was supplied in the **senderNonce** message, the header of the response shall include it as **recipNonce**.

The **PKIBody** is the element **kp** and is of type **CertRepMessage**. If the CA issued a certificate, the body will contain the following information:

- status will be granted or grantedWithMods; and
- certificate will contain the new X.509 version 3 certificate.

The certificate shall contain the following **extensions**:

- a subjectKeyldentifier field;
- at least one certificate policy OID in the **certificatePolicies** field; and
- an authority key identifier including a **Keyldentifier** field.

The **certificatePolicies** extension shall be identical to that found in the currently valid certificate. If a specific key identifier was specified in the **cr** message, the certificate shall contain that key identifier as the **subjectKeyldentifier** field. If no key identifier was supplied, the CA shall use the 96-bit SHA-1 hash of the subject public key as the **keyldentifier** in the **subjectKeyldentifier** field. The hash shall be calculated over the value (excluding tag and length) of the subject public key field in the certificate.

If the **kr** message included **extensions** other than the **subjectKeyldentifier**, the CA may modify or ignore the requested **extensions**.

The certificate shall include URLs in the **issuerAltName** extension and **distributionPoint** field of the **CRLDistributionPoints** extension if the issuer's certificates or CRLs are not available from a well known X.500 directory.

The failinfo field may not be present if status is granted or grantedWithMods.

If the CA rejected the request, the body shall include the following information:

- status will be rejected; and
- failinfo will contain the appropriate failure codes:
  - **badAlg** indicates that the CA cannot validate the signature because the algorithm identifier is unrecognized or unsupported;
  - badPoP indicates the signature in the popoSigningKey field was checked but did not match;
  - badMessageCheck indicates that the signature in the PKIProtection field was checked but did not match;
  - **badRequest** indicates that the responder does not permit or support the transaction;
  - badTime indicates that the messageTime field in the message header was not sufficiently close to the responder's system time; and
  - badCertId indicates that no certificate could be identified matching the nonzero serial field.

The certificate field may not be present if status is rejected.

The **PKIProtection** field contains the CA's signature, calculated on the DER encoded sequence of the header and body.

#### Confirmation Message

Upon receipt of the kp, the certificate holder shall generate a PKIConfirm message. PKIHeader

data shall be identical to the certificate request from the RA to the CA, with the exception of **messageTime**.

If the request was accepted, the **PKIProtection** field contains the certificate holder's signature, calculated on the DER encoded sequence of the header and body using the private key corresponding to the new signature certificate. If the request was rejected, the **PKIProtection** field contains the certificate holder's signature, calculated on the DER encoded sequence of the header and body using the private key corresponding to the currently valid signature certificate.

Upon receipt of the **PKIConfirm**, the CA shall generate a **PKIConfirm** message. **PKIHeader** data shall be identical to the **kr**, with the exception of **messageTime**.

The **PKIProtection** field contains the CA's signature, calculated on the DER encoded sequence of the header and body.

## 3.5.3 Self-Registration Request (New Subject)

An entity that has never obtained a certificate from a particular CA may request issuance of a new certificate directly from that CA. The requesting entity creates a **PKIMessage** with body type **ir** requesting a certificate and includes proof of possession of the private key corresponding to the public key in the certificate request. The entity protects the message with a SHA-1 HMAC using a secret key provided by the RA.

If the CA accepts self-registration requests, it will return an **ip** message to the certificate holder. This message will contain the certificate or a reason code for the transaction failure.

## RA-Entity Out-of-Band Transaction

The self-registration request for a certificate begins with the exchange of a secret known to the RA to the entity requesting a certificate. This information will allow the entity to authenticate themselves to the CA through generation of a message authentication code from the shared secret.

The precise content and format of this out-of-band transaction is not specified. However, it should be noted that both the secret key and the public key material for the trusted CA must be conveyed to the entity in a trusted fashion. So, this transaction should include authentication information for the CA from which the certificate will be requested and the public key material for the trusted CA.

## Self-Registration Request from Certificate Holder to CA

The requester creates a **PKIMessage** with a **PKIBody** element **ir**. The **PKIHeader** includes the following information:

- **pvno** is one;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the (proposed) distinguished name of the requester or an electronic mail address;
- **recipient** is the distinguished name of the CA; and
- protectionAlg is the algorithm identifier for the signature algorithm used to protect the

message.

The message body is **CertReqMessages**, which is a sequence of one or more **CertReqMessage** fields. For this transaction, **CertReqMessages** is a sequence of one **CertReqMessage**. The **CertReqMessage** will include the following information:

- **certReq** contains the information that the requester would like included in the certificate;
- **popoSKInput** contains the public key mac value; and
- **pop** provides proof of possession of the private key for the new certificate.

The **pop** field shall be generated using the private key corresponding to the public key in the **CertTemplate**. The input data (**popInput**) for generating **pop** shall be constructed from the public key mac value in **popSKInput** and the public key in the **CertTemplate**.

The certReq is a CertRequest, which is a sequence of a certReqID, a CertTemplate, and controls. For this transaction:

- certReqID is any integer;
- certTemplate is a CertTemplate including the public key for the new certificate; and
- **controls** is omitted.

The **CertTemplate** will include the following information:

- version is v3 (2); and
- **publicKey** provides the public key for the new certificate.

The following information may be included in the **CertTemplate**:

- **signingAlg** specifies the preferred signature algorithm;
- **subject** is present if and only if serial equals zero, and specifies the distinguished name for the prospective certificate holder; and
- **extensions** requests a particular certificate policy OID be specified in the certificate.

The request shall not include the following information:

- issuerUID; and
- subjectUID.

The **PKIProtection** field contains a value that is generated by the requester using the secret value obtained from the RA. The entity generates a 96 bit SHA1-HMAC using the secret key provided by the RA. The **protectionAlg** field shall be set to SHA1-HMAC, and the value of **PKIProtection** shall be the 96 bit message authentication code. The input to the calculation of the **PKIProtection** is the DER encoding of the following data structure:

```
ProtectedPart ::= SEQUENCE {
PKIHeader,
PKIBody}
```

Self-Registration Request Response from CA to Certificate Requester

The CA will return a **PKIMessage** with a **PKIBody** element **ip** to the certificate holder.

The **PKIHeader** includes the following information:

- pvno is one;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the distinguished name of the CA;
- **recipient** is the value of sender in the certificate request header; and
- **protectionAlg** is the algorithm identifier for the signature algorithm used to protect the message.

If a **transactionID** was supplied in **cr** message, the header of the response will include the same **transactionID**. If a **senderNonce** was supplied in the **senderNonce** message, the header of the response shall include it as **recipNonce**.

The **PKIBody** is an **ip** element and is of type **CertRepMessage**. If the CA issued a certificate, the body will contain the following information:

- status will be granted or grantedWithMods; and
- **certificate** will contain the new X.509 version 3 certificate;

The failinfo field may not be present if status is granted or grantedWithMods.

If the CA rejected the request, the body shall include the following information:

- status will be rejected; and
- failinfo will contain the appropriate failure codes:
  - badAlg indicates that the CA cannot validate the signature because the algorithm identifier is unrecognized or unsupported;
  - badPoP indicates the signature in the popoSigningKey field was checked but did not match;
  - **badMessageCheck** indicates the mac in the **PKIProtection** field was rejected;
  - **badRequest** indicates that the responder does not permit or support the transaction;
  - badTime indicates that the messageTime field in the message header was not sufficiently close to the responder's system time; and
  - badCertId indicates that no certificate could be identified matching the nonzero serial field.

The **certificate** field may not be present if **status** is **rejected**. If present, the certificate shall conform to the profile presented in section 3.1.1.

The certificate shall contain the following **extensions**:

- a subjectKeyIdentifier field;
- at least one certificate policy OID in the **certificatePolicies** field; and
- an authority key identifier including a **Keyldentifier** field.

If a specific key identifier was specified in the **ir** message, the certificate shall contain that key identifier as the **subjectKeyldentifier** field. If no key identifier was supplied, the CA shall use the 96-bit SHA-1 hash of the subject public key as the **keyldentifier** in the **subjectKeyldentifier** field. The hash shall be calculated over the value (excluding tag and length) of the subject public key field in the certificate.

If the ir message included extensions other than the subjectKeyldentifier, the CA may modify or ignore the requested extensions.

The certificate shall include URLs in the **issuerAltName** extension and **distributionPoint** field of the **CRLDistributionPoints** extension if the issuer's certificates or CRLs are not available from a well known X.500 directory.

The **PKIProtection** field contains the CA's signature, calculated on the DER encoded sequence of the header and body.

#### Confirmation Message

Upon receipt of the **ip**, the certificate holder shall generate a **PKIConfirm** message. **PKIHeader** data shall be identical to the certificate request from the requester to the CA, with the exception of **messageTime**.

If the request was accepted, the **PKIProtection** field contains the certificate holder's signature, calculated on the DER encoded sequence of the header and body using the private key corresponding to the new signature certificate. If the request was rejected, the **PKIProtection** field contains the SHA-1 HMAC, calculated on the DER encoded sequence of the header and body using the shared secret that authenticated the certificate request. (If the request was rejected for **badMessageCheck**, the **PKIConfirm** message need not be generated.)

Upon receipt of the **PKIConfirm**, the CA shall generate a **PKIConfirm** message. **PKIHeader** data shall be identical to the **ir**, with the exception of **messageTime**.

The **PKIProtection** field contains the CA's signature, calculated on the DER encoded sequence of the header and body.

## 3.5.4 Self-Registration Request (Known Subject)

An entity that is not a current certificate holder but has previously obtained a certificate from a particular CA may request issuance of a new certificate directly from that CA. The requesting entity creates a **PKIMessage** with body type **cr** requesting a certificate. This message includes proof of possession of the private key corresponding to the public key in the certificate request. The entity protects the message with a mac using a secret key provided by the RA.

If the CA accepts self registration requests, it will return a **cp** message to the certificate holder. This message will contain the certificate or a reason code for the transaction failure.

## RA-Entity Out-of-Band Transaction

The self-registration request for a certificate begins with the delivery of a secret known to the RA to the entity requesting a certificate. This information allows the entity to authenticate themselves to the CA through generation of a message authentication code from the shared secret.

The precise content and format of this out-of-band transaction is not specified. However, it should be noted that both the secret key and the public key material for the trusted CA must be conveyed to the entity in a trusted fashion. So, this transaction should include authentication information for the CA from which the certificate will be requested and the public key material for the trusted CA.

## Self-Registration Request from Certificate Holder to CA

The requester creates a **PKIMessage** with a **PKIBody** element **cr**. The **PKIHeader** includes the following information:

- **pvno** is one;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the (proposed) distinguished name of the requester or an electronic mail address;
- **recipient** is the distinguished name of the CA; and
- **protectionAlg** is the algorithm identifier for the signature algorithm used to protect the message.

The message body is **CertReqMessages**, which is a sequence of one or more **CertReqMessage** fields. For this transaction, **CertReqMessages** is a sequence of one **CertReqMessage**. The **CertReqMessage** will include the following information:

- **certReq** contains the information that the requester would like included in the certificate;
- **popoSKInput** contains the mac value; and
- **pop** provides proof of possession of the private key for the new certificate.

The **pop** field shall be generated using the private key corresponding to the public key in the **CertTemplate**. The input data (**popInput**) for generating **pop** shall be constructed from the public key mac value in **popSKInput** and the public key in the **CertTemplate**.

The certReq is a CertRequest, which is a sequence of a certReqID, a CertTemplate, and controls. For this transaction:

- **certReqID** is any integer; and
- **certTemplate** is a **CertTemplate** including the public key for the new certificate.

The **CertTemplate** will include the following information:

- version is v3 (2); and
- **publicKey** provides the public key for the new certificate.

The following information may be included in the **CertTemplate**:

- **signingAlg** specifies the preferred signature algorithm;
- **subject** is present if and only if serial equals zero, and specifies the distinguished name for the prospective certificate holder; and
- **extensions** requests a particular certificate policy OID be specified in the certificate.

The request shall not include the following information:

- issuerUID; and
- subjectUID.

The **PKIProtection** field contains a value that is generated by the requester using the secret value obtained from the RA. The entity generates a mac using the secret key provided by the RA. The **protectionAlg** field shall be set to the OID of the algorithm used to generate the mac, and the value of **PKIProtection** shall be the message authentication code. The input to the calculation of

the **PKIProtection** is the DER encoding of the following data structure:

```
ProtectedPart ::= SEQUENCE {
PKIHeader,
PKIBody}
```

#### Self-Registration Request Response from CA to Certificate Requester

The CA will return a **PKIMessage** with a **PKIBody** element **cp** to the certificate holder.

The **PKIHeader** includes the following information:

- **pvno** is one;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the distinguished name of the CA;
- **recipient** is the value of **sender** in the certificate request header; and
- **protectionAlg** is the algorithm identifier for the signature algorithm used to protect the message.

If a **transactionID** was supplied in **cr** message, the header of the response will include the same **transactionID**. If a **senderNonce** was supplied in the request, the header of the response shall include it as **recipNonce**.

The **PKIBody** is a **cp** element and is of type **CertRepMessage**. If the CA issued a certificate, the body will contain the following information:

- status will be granted or grantedWithMods; and
- **certificate** will contain the new X.509 version 3 certificate;

The failinfo field may not be present if status is granted or grantedWithMods.

If the CA rejected the request, the body shall include the following information:

- status will be rejected; and
- **failInfo** will contain the appropriate failure codes:
  - badAlg indicates that the CA cannot validate the signature because the algorithm identifier is unrecognized or unsupported;
  - badPoP indicates the signature in the popoSigningKey field was checked but did not match;
  - **badMessageCheck** indicates the mac in the **PKIProtection** field was rejected;
  - **badRequest** indicates that the responder does not permit or support the transaction;
  - badTime indicates that the messageTime field in the message header was not sufficiently close to the responder's system time; and
  - badCertId indicates that no certificate could be identified matching the nonzero serial field.

The **certificate** field may not be present if **status** is **rejected**. If present, the certificate shall conform to the profile presented in section 3.1.1.

The certificate shall contain the following **extensions**:

- a subjectKeyIdentifier field;
- at least one certificate policy OID in the **certificatePolicies** field; and
- an authority key identifier including a Keyldentifier field.

If a specific key identifier was specified in the **cr** message, the certificate shall contain that key identifier as the **subjectKeyldentifier** field. If no key identifier was supplied, the CA shall use the first 96 bits of the SHA-1 hash of the subject public key as the **keyldentifier** in the **subjectKeyldentifier** field. The hash shall be calculated over the value (excluding tag and length) of the subject public key field in the certificate.

If the **cr** message included **extensions** other than the **subjectKeyldentifier**, the CA may modify or ignore the requested **extensions**.

The certificate shall include URLs in the **issuerAltName** extension and **distributionPoint** field of the **CRLDistributionPoints** extension if the issuer's certificates or CRLs are not available from a well known X.500 directory.

The **PKIProtection** field contains the CA's signature, calculated on the DER encoded sequence of the header and body.

#### Confirmation Message

Upon receipt of the **cp**, the certificate holder shall generate a **PKIConfirm** message. **PKIHeader** data shall be identical to the certificate request from the RA to the CA, with the exception of **messageTime**.

If the request was accepted, the **PKIProtection** field contains the certificate holder's signature, calculated on the DER encoded sequence of the header and body using the private key corresponding to the new signature certificate. If the request was rejected, the **PKIProtection** field contains the mac, calculated on the DER encoded sequence of the header and body using the shared secret that authenticated the certificate request. (If the request was rejected for **badMessageCheck**, the **PKIConfirm** message need not be generated.)

Upon receipt of the **PKIConfirm**, the CA shall generate a **PKIConfirm** message. **PKIHeader** data shall be identical to the **cr**, with the exception of **messageTime**.

The **PKIProtection** field contains the CA's signature, calculated on the DER encoded sequence of the header and body.

## 3.5.5 PKCS #10 Self-Registration Request

An entity that is not a current certificate holder may request issuance of a certificate directly from the CA using the certificate request syntax defined in PKCS #10. The requesting entity creates a **PKIMessage** of type **PKCSReq** requesting a certificate and includes proof of possession of the private key corresponding to the public key in the body of the certificate request, and protects the **PKIMessage** using a secret key provided by the RA in an out-of-band transaction.

The CA will return a certificate request response message to the certificate requester. This message will contain the certificate or a reason code for the transaction failure.

The out-of-band transaction with the RA and the CA response are identical to the corresponding steps in the Self-Registration Request defined in section 3.5.3.

#### Self registration Request from Certificate Holder to CA

The requester creates a **PKIMessage** with a **PKIBody** element **p10cr**. The **PKIHeader** includes the following information:

- pvno is one;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the (proposed) distinguished name of the requester or an electronic mail address;
- **recipient** is the distinguished name of the CA; and
- **protectionAlg** is the algorithm identifier for the signature algorithm used to protect the message.

The **PKIBody** is a **PKIBody** element **p10cr** which is of type **PKCS10CertReqMessages**. This type is a sequence of a **certificationRequestInfo**, a **signatureAlgorithm** and a **signature**. The **certificationRequestInfo** will include the following information:

- version is v3 (2);
- **subject** is present if and only if serial equals zero, and specifies the distinguished name for the prospective certificate holder; and
- **subjectPublicKeyInfo** provides the public key and corresponding algorithm identifier for the new certificate.

The **signatureAlgorithm** field contains the algorithm identifier associated with the private key used to generate the **signature** field; the signature is generated using the DER-encoded **certificationRequestInfo** as input.

The **PKIProtection** field contains a value that is generated by the requester using the secret value obtained from the RA. The entity generates a mac using the secret key provided by the RA. The **protectionAlg** field shall be set to the OID of the message authentication algorithm used, and the value of **PKIProtection** shall be the message authentication code. The input to the calculation of the **PKIProtection** is the DER encoding of the following data structure:

ProtectedPart ::= SEQUENCE {		
header	PKIHeader,	
body	PKIBody}	

#### PKCS Certificate Request Response from CA to Certificate Requester

The CA will return a **PKIMessage** with a **PKIBody** element **cp** to the certificate holder.

The **PKIHeader** includes the following information:

- **pvno** is one;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the distinguished name of the CA;
- recipient is the value of sender in the certificate request header; and

• **protectionAlg** is the algorithm identifier for the signature algorithm used to protect the message.

If a **transactionID** was supplied in **PKCSReq** message, the header of the response will include the same **transactionID**.

The **PKIBody** element is a **cp**, which is of type **CertRepMessage**. If the CA issued a certificate, the body will contain the following information:

- status will be granted or grantedWithMods; and
- **certificate** will contain the new X.509 version 3 certificate.

If a specific key identifier was specified in the **p10cr** message, the certificate shall contain that key identifier as the **subjectKeyldentifier** field. If no key identifier was supplied, the CA shall use the 96-bit SHA-1 hash of the subject public key as the **keyldentifier** in the **subjectKeyldentifier** field. The hash shall be calculated over the value (excluding tag and length) of the subject public key field in the certificate.

#### The failinfo field may not be present if status is granted or grantedWithMods.

If the CA rejected the request, the body shall include the following information:

- status will be rejected; and
- failinfo will contain the appropriate failure codes:
  - **badAlg** indicates that the CA cannot validate the signature because the algorithm identifier is unrecognized or unsupported;
  - badPoP indicates the signature in the pk10cr's signature field was checked but did not match;
  - badMessageCheck indicates the mac in the PKIMessage's PKIProtection field was rejected;
  - badRequest indicates that the responder does not permit or support the transaction; and
  - **badTime** indicates that the **messageTime** field in the message header was not sufficiently close to the responder's system time.

The certificate field shall not be present if status is rejected.

The certificate shall contain the following **extensions**:

- a subjectKeyIdentifier field;
- at least one certificate policy OID in the certificatePolicies field; and
- an authority key identifier including a **Keyldentifier** field.

If the **cr** message included **extensions** other than the **subjectKeyldentifier**, the CA may modify or ignore the requested **extensions**.

The certificate shall include URLs in the **issuerAltName** extension and **distributionPoint** field of the **CRLDistributionPoints** extension if the issuer's certificates or CRLs are not available from a well known X.500 directory.

The **PKIProtection** field contains the CA's signature, calculated on the DER encoded sequence of the header and body.

### 3.5.6 Revocation Request

Certificate holders may request revocation of their own certificates. To perform this function the certificate holder generates a **RevReq** message, signs it and sends it to the CA that issued the certificates. The signature must be generated with the private key corresponding to an unexpired, unrevoked signature certificate issued by the same CA. The **RevReq** message shall identify the certificate(s) to be revoked and the reason for the revocation. The CA responds with a **RevRep** message.

RAs may request revocation of a certificate issued to an entity on behalf of the certificate holder or the certificate holder's organization. To perform this function, the RA generates a **RevReq** message, signs it with the RA's private key, and sends it to the CA. The RA shall generate a pseudo-random number and shall place it in the **transactionID** field. The **RevReq** message shall identify the certificate(s) to be revoked and the reason for the revocation.

The CA will respond to the revocation requester with an **rp** (**RevRep**) message. If the **rr** (**RevReq**) message included a **transactionID**, the CA shall include its contents as the **transactionID** in the **rp** message. The **rp** message shall contain, at a minimum, the status of the request in the **status** field and identify the certificate(s) for which revocation is requested in the **revDetails** field.

#### Revocation Request from RA or Certificate Holder to CA

The RA or the certificate holder creates a **PKIMessage** with a **PKIBody** element **rr**. The **PKIHeader** includes the following information:

- pvno is one;
- **transactionID** is an integer unique to this transaction for this RA or any integer for the end entity;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the distinguished name of the RA or the certificate holder;
- recipient is the distinguished name of the CA; and
- **protectionAlg** is the algorithm identifier for the signature algorithm used to protect the message.

The **PKIBody** is **RevReqContent**, which is a sequence of **RevDetails**. **RevDetails** is a sequence of **CertDetails** and three optional fields: reason flags; and date and time of compromise or loss; and **crlEntryDetails** (a sequence of CRL entry **extensions**). **CertDetails** is defined as a **CertTemplate**. For this interoperability specification, **RevReqContent** is a sequence of one **RevDetails**. **CertDetails**, at a minimum, includes the following information:

- serial, which contains the serial number of the certificate; and
- issuer, which contains the distinguished name of the certificate issuer.

or

- subject, which contains the distinguished name of the certificate holder; and
- **issuer**, which contains the distinguished name of the certificate issuer.

#### CertDetails may also include a subjectKeyIdentifier in the extensions field.

(If the requester intends to revoke all certificates issued to a particular subject, **CertDetails** shall include only the **subject** and **issuer**. That is, a revocation request specifying either the serial number or **subjectKeyldentifier** applies to that single certificate.)

The **RevDetails** shall include **crlEntryDetails** with a **reasonCode** extension, and may include the **invalidityDate** extension to specify the time after which the certificate(s) should not be trusted. The reason code may not be **removeFromCRL**.

The **PKIProtection** field contains the requester's signature, calculated on the DER encoded sequence of the header and body. The end-entity shall generate the signature with a private key corresponding to a currently valid signature certificate issued by **recipient**.

#### Revocation Response from CA to Requester

The CA will return a **PKIMessage** with a **PKIBody** element **rp** to the requester.<sup>25</sup>

The **PKIHeader** includes the following information:

- pvno is one;
- transactionID is the same as the transactionID field in the CertReq message;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the distinguished name of the CA;
- **recipient** is the distinguished name of the RA; and
- **protectionAlg** is the algorithm identifier for the signature algorithm used to protect the message.

If a **senderNonce** was supplied in the request message, the header of the response shall include it as **recipNonce**.

The **PKIBody** is **RpContent**. If the CA revoked the certificate, the body will contain the following information:

- status will be granted or grantedWithMods; and
- **revDetails** will contain the **CertId**(s) of the revoked certificate(s);

#### The failinfo field may not be present if status is granted or grantedWithMods.

If the CA rejected the request, the body shall include the following information:

- status will be rejected; and
- failinfo will contain the appropriate failure codes:
  - badAlg indicates that the CA cannot validate the signature because the algorithm identifier is unrecognized or unsupported;
  - badMessageCheck indicates that the signature in the PKIProtection fields was checked but did not match;
  - **badRequest** indicates that the responder does not permit or support the transaction;
  - **badTime** indicates that the **messageTime** field in the message header was not

<sup>&</sup>lt;sup>25</sup> If the requester is an RA, the CA may optionally send the **RevRep** message to the certificate holder as well.

- sufficiently close to the responder's system time; or
- **badCertId** indicates that the information in **CertDetails** did not identify an unexpired, unrevoked certificate.

If the certificate in question can be determined, **revDetails** will contain the **CertId** of the certificate whose revocation was rejected.

The **PKIProtection** field shall contain the CA's signature, calculated on the DER encoded sequence of the header and body.

If the CA generates CRLs, and the revocation request was accepted, the CRL entry shall have the following values:

- the serial number of the revoked certificate in the userCertificate field;
- the **revocationDate** shall be the day and time the revocation request was received;
- the **crIEntryExtensions** shall be present and include:
  - the reasonCode shall be the reasonCode found in the RevDetails field, unless otherwise specified in the CA's policy;
  - optionally, the invalidityDate extension may be the badSinceDate found in the RevDetails field, if provided.

## 3.5.7 Encryption Certificate Request for End-Entity Generated Key Pairs

An entity that is a current signature certificate holder may request issuance of a new certificate for encryption keys directly from the CA that issued the current signature certificate. The requesting entity:

- generates a key management key pair;
- creates a PKI **cr** (certificate request) message requesting a key transport or key agreement certificate and includes the public key;
- signs the message with the private key corresponding to the entity's unexpired, unrevoked signature certificate; and
- transmits it to the CA.

If the CA's Certificate Practice Statement requires proof of possession for encryption key pairs, the CA generates a challenge message.<sup>26</sup> If challenged, the requester must demonstrate that it possesses the private key that corresponds to the public key in the certificate request message.

The CA will generate and return a **cp** (certificate response) message to the certificate holder. This message will contain the certificate or a reason code for the transaction failure.

If the CA's Certificate Practice Statement does not require proof of possession for encryption key pairs, the CA will generate and return the **cp** (certificate response) message to the certificate holder. As above, this message will contain the certificate or a reason code for the transaction failure.

 $<sup>^{26}</sup>$  If the CA can determine from the request that a certificate will not be issued, the CA does not generate a challenge message. Instead, the CA returns a **cp** message with a reason code for the transaction failure.

## Certificate Request from Certificate Holder to CA

The certificate holder creates a certificate request message: a **PKIMessage** with **PKIBody** element **cr**. The **PKIHeader** includes the following information:

- pvno is one;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the distinguished name of the certificate holder;
- **recipient** is the distinguished name of the CA; and
- **protectionAlg** is the algorithm identifier for the signature algorithm used to protect the message.

The message body is **CertReqMessages**, which is a sequence of one or more **CertReqMessage** fields. For this transaction, **CertReqMessages** is a sequence of one **CertReqMessage**. The **CertReqMessage** will include the following information:

• **certReq** contains the information that the requester would like included in the certificate.

The certReq is a CertRequest, which is a sequence of a certReqID, a CertTemplate, and controls. For this transaction:

- **certReqID** is any integer; and
- certTemplate is a CertTemplate.

The **CertTemplate** will include at a minimum, the following information:

- version is v3 (2); and
- **publicKey** provides the public key for the new certificate.

The following information may be included in the **CertTemplate**:

• **signingAlg** specifies the preferred signature algorithm for the CA to sign the certificate.

If **signingAlg** does not appear, the CA should sign with the algorithm identified in the **protectionAlg** field in the **PKIHeader**.

The request shall not include the following information:

- issuerUID; and
- subjectUID.

The **PKIProtection** field contains a signature generated using the private key associated with the current unexpired, unrevoked certificate and calculated upon the DER encoded sequence of the header and body.

#### Challenge and Response (Proof-of-Possession)

The CA may, optionally, generate a challenge message to verify that the requester possesses the decryption key corresponding to the encryption key. The requester must ... [Need to complete this thought, whatever it was.]

The specific details of the challenge-response protocol differ according to the key management algorithm associated with the public key. For Diffie-Hellman and Elliptic Curve Diffie-Hellman keys, the challenge response messages are specified in (sec. 3.5.8). For RSA key transport keys, the challenge response messages are specified in (sec. 3.5.9).

## Certificate Response from CA to Certificate Holder

The CA will return a certificate response message (a **PKIMessage** with **PKIBody** element **cp**) to the certificate holder.

The **PKIHeader** includes the following information:

- **pvno** is one;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the distinguished name of the CA;
- **recipient** is the distinguished name of the certificate holder and the **sender** of the previous (popdecr or **cr**) message; and
- **protectionAlg** is the algorithm identifier for the signature algorithm used to protect the message.

If a **transactionID** was supplied in the previous message, the header of the response will include the same **transactionID**. If a **senderNonce** was supplied in the **cp** or **popdecr** message, the header of the response shall include it as **recipNonce**.

The **PKIBody** is the element **cp** and is of type **CertRepMessage**. If the CA issued a certificate, the body will contain the following information:

- status will be granted or grantedWithMods; and
- certificate will contain the new X.509 version 3 certificate.

The certificate shall contain the following **extensions**:

- a subjectKeyldentifier field;
- at least one certificate policy OID in the **certificatePolicies** field; and
- an authority key identifier including a **Keyldentifier** field.

If a specific key identifier was specified in the **cr** message, the certificate shall contain that key identifier as the **subjectKeyldentifier** field. If no key identifier was supplied, the CA shall use the 96-bit SHA-1 hash of the subject public key as the **keyldentifier** in the **subjectKeyldentifier** field. The hash shall be calculated over the value (excluding tag and length) of the subject public key field in the certificate.

If the **cr** message included **extensions** other than the **subjectKeyldentifier**, the CA may modify or ignore the requested **extensions**.

The certificate shall include URLs in the **issuerAltName** extension and **distributionPoint** field of the **CRLDistributionPoints** extension if the issuer's certificates or CRLs are not available from a well-known X.500 directory.

The failinfo field may not be present if status is granted or grantedWithMods.

If the CA rejected the request, the body shall include the following information:

- status will be rejected; and
- failinfo will contain the appropriate failure codes:
  - **badAlg** indicates that the CA cannot validate the signature because the algorithm

identifier is unrecognized or unsupported;

- badPoP indicates the CA generated a challenge but the end-entity's response message did not demonstrate possession of the private key;
- badMessageCheck indicates that the signature in the PKIProtection field was checked but did not match;
- **badRequest** indicates that the responder does not permit or support the transaction;
- badTime indicates that the messageTime field in the message header was not sufficiently close to the responder's system time; and
- badCertId indicates that no current valid, unrevoked signature certificate could be identified for the requester.

The certificate field may not be present if status is rejected.

The **PKIProtection** field contains the CA's signature, calculated on the DER encoded sequence of the header and body.

#### Confirmation Message

Upon receipt of the **cp** message, the certificate holder shall generate a **PKIConfirm** message. **PKIHeader** data shall be identical to the certificate request from the RA to the CA, with the exception of **messageTime**.

The **PKIProtection** field contains the certificate holder's signature, calculated on the DER encoded sequence of the header and body using the private key corresponding to the currently valid signature certificate.

## 3.5.8 Proof of Possession for Diffie-Hellman and Elliptic Curve Key Agreement Keys

If the CA's Certificate Practice Statement requires proof of possession for key agreement key pairs, the CA generates a challenge message.<sup>27</sup> To generate the challenge, the CA will perform the following functions:

- 1. Generate a random number A;
- 2. Generate a witness, R, from the random number A using a one-way function (e.g., SHA-1);
- 3. Identify or generate a key pair for use with the requested public key;<sup>28</sup>
- 4. Derive a symmetric (e.g., triple DES) key through a key agreement algorithm using the requester's public key and the key pair identified in (4.);
- 5. Construct the challenge **Rand** from the random number A and requester's distinguished name (the **sender** field) in the **cr** message and DER encode the **SEQUENCE** of these values;
- 6. Encrypt the sequence **Rand** under the derived symmetric key. The first byte of **Rand** shall be the first byte of the input stream. (If necessary, pad the input stream with random data to complete the last block);
- 7. Construct an encryptedValue structure containing the symmetric algorithm in the symmalg

 $<sup>^{27}</sup>$  If the CA can determine from the request that a certificate will not be issued, the CA does not generate a challenge message. Instead, the CA returns a **cp** message with a reason code for the transaction failure.

<sup>&</sup>lt;sup>28</sup> If the CA generates a key pair, it will need to issue a certificate containing the public key for inclusion in the **extraCerts** field.

field, and the encrypted value of the sequence **Rand** in the **encValue** field;

8. Generate and return a **POPODecKeyChallContent** (proof of possession challenge) message to the certificate holder. This message contains the OID for the hash function, the witness R, and the **encryptedValue** structure (which contains the encrypted sequence **Rand** containing the challenge A and requester's name). The message body and header are protected with the CA's signature. The CA's certificate with the public key from (4.) is appended in the **extraCerts** field.

9. The CA transmits this message to the requester as a challenge.

The requester must send a response to the challenge to demonstrate possession of the corresponding private key. The requester decrypts the challenge and returns A to the CA in a signed message if the hash of the random number A recovered from the decrypted challenge matches R, the **witness** value.

The requester performs the following steps:

- 1. Derive a symmetric key for the algorithm specified in the **symmAlg** field of the **encryptedValue** structure using its private key and the public key in the CA's certificate from the **extraCerts** field;
- 2. Decrypt the **encValue** field (and recover the sequence **Rand**) using the derived symmetric key and the symmetric algorithm specified in the **symmAlg** field of the **encryptedValue** structure;
- 3. Generate a hash of the random number A using the one-way function specified in **owf** and verify that it matches the witness, R;
- 4. Create a PKI **POPODecKeyRespContent** message containing the random number A;
- 5. Sign the message with the private key corresponding to the entity's unexpired, unrevoked signature certificate; and
- 6. Transmit the message to the CA.

The remainder of this section describes the contents of challenge-response messages for Diffie-Hellman and Elliptic Curve Diffie-Hellman Key Agreement keys in detail. It assumes SHA-1 as the hash function and the Triple Data Encryption Algorithm in ecb mode with two keys for symmetric encryption.

#### Challenge (Proof-of-Possession)

The CA may, optionally, generate a challenge message to verify that the requester possesses the decryption key corresponding to the encryption key. The challenge message is a **PKIMessage** with **PKIBody** element **popdecc**.

The **PKIHeader** includes the following information:

- **pvno** is one;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the distinguished name of the CA;
- **recipient** is the distinguished name of the certificate holder and the **sender** of the **cr** message; and
- **protectionAlg** is the algorithm identifier for the signature algorithm used to protect the message.

If a transactionID was supplied in the cr message, the header of the response will include the

same **transactionID**. If a **senderNonce** was supplied in the **cr** message, the header of the response shall include it as **recipNonce**.

### The PKIBody is the element popdecc and is of type POPODecKeyChallContent.

**POPODecKeyChallContent** shall be a **SEQUENCE** of exactly one **Challenge**, which shall include the following information:

- **owf** will contain the algorithm identifier for the SHA-1 hash function;
- witness will contain the result of applying the one-way function specified by **owf** to an integer value *A* randomly generated for this transaction; and
- challenge, an OCTET STRING, shall contain the ASN.1 DER-encoded structure encryptedValue and contain the following fields:
  - symmAlg shall contain the object identifier tDEA-ecb and the keyingOption shall be option-2 (indicating two independent keys, K1 and K2, with K3 = K1);
  - encValue shall contain the encrypted structure Rand, encrypted with tDEA in ecbmode using the two key option. The result shall be encoded as the value of the BIT STRING where the most significant bit of the BIT STRING is the most significant bit of the encrypted data. The symmetric key used to perform the encryption shall be derived using from the CA's private key and the requester's public key using the specified key agreement algorithm.
    - For Diffie-Hellman keys, the shared secret shall be generated using the dHStatic mode. The tDEA key shall be derived from the shared secret using the "Key Derivation Method Based on ASN.1" specified in [X9.42]. The key derivation input AlgorithmID shall be the tDEA-ecb object identifier. The optional key derivation inputs *PartyUInfo*, *PartyVInfo*, *SuppPrivInfo*, and *SuppPubInfo* shall be omitted.
    - For elliptic curve Diffie-Hellman, the shared secret shall be generated using the "Static Unified Model Scheme", as specified in [X9.63]. The tDEA key shall be derived using the key derivation method specified in [X9.63] for the "Static Unified Model Scheme"; the optional key derivation input *SharedInfo* shall be omitted.

**Rand** shall be a sequence of the following fields:

- the integer value A in the field **int**; and
- the requester's distinguished name in the field **sender**.

The **PKIProtection** field shall contain the CA's signature, calculated on the DER encoded sequence of the header and body.

The **extraCerts** field shall contain the certificate for the public key of the key pair used by the CA to perform key agreement and derive the symmetric key.

Response to Challenge (Proof-of-Possession)

If the requester receives a challenge message from the CA, it is required to generate a response to prove possession of the private key. The response message is a **PKIMessage** with **PKIBody** element **popdecr**.

The **PKIHeader** includes the following information:

- pvno is one;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the distinguished name of the entity requesting the key management certificate;
- **recipient** is the distinguished name of the CA; and
- **protectionAlg** is the algorithm identifier for the signature algorithm used to protect the message.

If a **transactionID** was supplied in the **popdecc** message, the header of the response will include the same **transactionID**. If a **senderNonce** was supplied in the **popdecc** message, the header of the response shall include it as **recipNonce**.

The **PKIBody** is the element **popdecr** and is of type **POPODecKeyRespContent**. The **POPODecKeyRespContent** shall be a sequence of exactly one **INTEGER**; the **INTEGER**'s value shall be the retrieved random number. The symmetric key used to perform the decryption shall be derived using from the requester's private key and the CA's public key (from the **extraCerts** field) using the specified key agreement algorithm as specified above.

The **PKIProtection** field contains a signature generated using the private key associated with a current unexpired, unrevoked certificate and calculated upon the DER encoded sequence of the header and body.

## 3.5.9 Proof of Possession for RSA Key Transport keys

If the CA's Certificate Practice Statement requires proof of possession for encryption key pairs, the CA generates a challenge message.<sup>29</sup> To generate the challenge, the CA will perform the following functions:

- 1. Generate a random number A;
- 2. Generate a witness, R, from the random number A using a one-way function (e.g., SHA-1);
- 3. Generate a symmetric key (e.g., a triple-DES key);
- 4. Construct the challenge **Rand** from the random number A and requester's distinguished name (the **sender** field) in the **cr** message and DER encode the **SEQUENCE** of these values;
- 5. Encrypt the sequence **Rand** under the symmetric key generated by the CA. (If necessary, pad the input stream with random data to complete the last block);
- 6. Encrypt the symmetric key under the supplied RSA key;
- 7. Construct an **encryptedValue** structure containing the symmetric algorithm in the **symmalg** field, the encrypted symmetric key in the **encSymmKey** field, and the encrypted value of the sequence **Rand** in the **encValue** field;
- 8. Generate a **POPODecKeyChallContent** (proof of possession challenge) message to the certificate holder. This message contains the OID for the SHA-1 hash function, the witness R, and the **encryptedValue** structure.
- 9. Transmit the message to the requester.

The requester must send a response to the challenge to demonstrate possession of the corresponding private key. The requester decrypts the challenge and returns A to the CA in a

 $<sup>^{29}</sup>$  If the CA can determine from the request that a certificate will not be issued, the CA does not generate a challenge message. Instead, the CA returns a **cp** message with a reason code for the transaction failure.

signed message if the hash of the random number A recovered from the decrypted challenge matches the **witness** value.

The requester performs the following steps:

- 1. Use their private key to decrypt the symmetric key from the **encSymmKey** field of the **encryptedValue** structure;
- 2. Decrypt the sequence **Rand** the using the symmetric key and the symmetric algorithm specified in the **symmAlg** field of the **encryptedValue** structure;
- 3. Generate a hash of the random number A and verifies that it matches the witness, R;
- 4. Verify that the **sender** field in **Rand** is the requester's name;
- 5. Create a PKI **POPODecKeyRespContent** message containing the random number A; and
- 6. Sign the message with the private key corresponding to the entity's unexpired, unrevoked signature certificate; and
- 7. Transmit it to the CA.

The remainder of this section describes the contents of challenge-response messages for RSA key transport keys in detail. It assumes SHA-1 as the hash function and the Triple Data Encryption Algorithm in ecb mode with two keys for symmetric encryption.

#### Challenge (Proof-of-Possession)

The CA may, optionally, generate a challenge message to verify that the requester possesses the decryption key corresponding to the encryption key. The challenge message is a **PKIMessage** with **PKIBody** element **popdecc**.

The **PKIHeader** includes the following information:

- **pvno** is one;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the distinguished name of the CA;
- **recipient** is the distinguished name of the certificate holder and the **sender** of the **cr** message; and
- **protectionAlg** is the algorithm identifier for the signature algorithm used to protect the message.

If a **transactionID** was supplied in the **cr** message, the header of the response will include the same **transactionID**. If a **senderNonce** was supplied in the **cr** message, the header of the response shall include it as **recipNonce**.

The **PKIBody** is the element **popdecc** and is of type **POPODecKeyChallContent**. **POPODecKeyChallContent** shall be a **SEQUENCE** of exactly one **Challenge**, which shall include the following information:

- **owf** will contain the algorithm identifier for a one-way hash function;
- witness will contain the result of applying the one-way function specified by **owf** to an integer value *A* randomly generated for this transaction; and
- challenge, an OCTET STRING, shall contain the ASN.1 DER-encoded structure encryptedValue and contain the following fields:
  - the symmAlg shall contain the OID tDEA-ecb and the keyingOption shall be option-2

(indicating two independent keys, K1 and K2, with K3 = K1);

- **encSymmKey** shall contain the symmetric key generated by the CA, encrypted with the public key in the certificate request;
  - the symmetric key material (K1 and K2) shall be concatenated into *TwoKeys* as described in (sec. 3.1.2.6)
  - *TwoKeys* shall be encrypted using the RSA encryption algorithm according to [PKCS#1], and
  - the result shall be encoded as the value of the **BIT STRING** where the most significant bit of the **BIT STRING** is the most significant bit of the encrypted data.
- the encValue field, which is a BIT STRING, shall contain the encrypted structure Rand, encrypted with tDEA in the electronic code book mode using the symmetric key in encSymmKey. The result shall be encoded as the value of the BIT STRING where the most significant bit of the BIT STRING is the most significant bit of the encrypted data.

**Rand** shall be a sequence of the following fields:

- the integer value A in the field **int**; and
- the requester's distinguished name in the field **sender**.

The **PKIProtection** field contains the CA's signature, calculated on the DER encoded sequence of the header and body.

#### Response to Challenge (Proof-of-Possession)

If the requester receives a challenge message from the CA, it is required to generate a response to prove possession of the private key. The response message is a **PKIMessage** with **PKIBody** element **popdecr**.

The **PKIHeader** includes the following information:

- pvno is one;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the distinguished name of the entity requesting the key management certificate;
- **recipient** is the distinguished name of the CA; and
- **protectionAlg** is the algorithm identifier for the signature algorithm used to protect the message.

If a **transactionID** was supplied in the **popdecc** message, the header of the response will include the same **transactionID**. If a **senderNonce** was supplied in the **popdecc** message, the header of the response shall include it as **recipNonce**.

The **PKIBody** is the element **popdecr** and is of type **POPODecKeyRespContent**. **POPODecKeyRespContent** shall be a sequence of exactly one **INTEGER**; the **INTEGER**'s value shall be the retrieved random number.

The **PKIProtection** field contains a signature generated using the private key associated with the current unexpired, unrevoked signature certificate and calculated upon the DER encoded sequence of the header and body.

### 3.5.10 Request for Centrally-Generated Key Pair and Key Management Certificate

An entity that is a current signature certificate holder may request generation of an encryption key pair and issuance of a corresponding certificate from the CA that issued the current signature certificate. The requesting entity:

- generates a temporary key management key (e.g., RSA, Diffie-Hellman, or elliptic curve Diffie-Hellman);
- creates a PKI **cr** (certificate request) message requesting a key transport or key agreement certificate and includes the temporary key;
- signs the message with the private key corresponding to the entity's unexpired, unrevoked signature certificate; and
- transmits it to the CA.

If the CA's Certificate Practice Statement supports central generation of encryption key pairs, the CA performs the following functions:

- the CA generates the requested key pair and issues a key management certificate
- if the requester included a temporary Diffie-Hellman or ECDH key:
  - the CA generates or identifies a corresponding key pair;
  - performs key agreement with the requester's temporary public key and derives a symmetric key; and
  - encrypts the private key under the derived symmetric key.
- if the requester included an RSA key:
  - the CA generates a symmetric key;
  - encrypts the private key under the new symmetric key; and
  - encrypts the symmetric key under the supplied temporary RSA public key.
- generates and returns a **cp** (certificate response) message to the certificate holder. This message will contain the certificate and the encrypted private key or a reason code for the transaction failure. If the requester included a temporary DH or ECDH key, the CA's certificate with its public key will be included in the **extraCerts** field.

#### Centrally Generated Key Pair Request

The certificate holder creates a certificate request message: a **PKIMessage** with **PKIBody** element **cr**. The **PKIHeader** includes the following information:

- **pvno** is one;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the distinguished name of the certificate holder;
- recipient is the distinguished name of the CA; and
- **protectionAlg** is the algorithm identifier for the signature algorithm used to protect the message.

The message body is **CertReqMessages**, which is a sequence of one or more **CertReqMessage** fields. For this transaction, **CertReqMessages** is a sequence of one **CertReqMessage**. The **CertReqMessage** will include the following information:

• **certReq** contains the information that the requester would like included in the certificate.

The certReq is a CertReqRequest, which is a sequence of a certReqID, a CertTemplate, and controls. For this transaction:

- **certReqID** is any integer;
- certTemplate is a CertTemplate; and
- **controls** will contain the **protocolEncrKey** control. The value of the **protocolEncrKey** registration control will be the temporary public key with its associated parameters (if any) and the algorithm OID.

The **CertTemplate** will include the following information:

- version is v3 (2); and
- **publicKey** specifies the algorithm and optionally the parameters, for the requested key pair.

The following information may be included in the **CertTemplate**:

• **signingAlg** specifies the preferred signature algorithm.

If signingAlg does not appear, the CA should sign with the algorithm specified in protectionAlg.

The request shall not include the following information:

- issuerUID; and
- subjectUID.

The **PKIProtection** field contains a signature generated using the private key associated with a current unexpired, unrevoked certificate and calculated upon the DER encoded sequence of the header and body.

#### Centrally Generated Key Pair Response

The CA will return a key update response (a **PKIMessage** with **PKIBody** element **cp**) message to the certificate holder.

The **PKIHeader** includes the following information:

- pvno is one;
- **messageTime** is the current time with a granularity of seconds;
- **sender** is the distinguished name of the CA;
- **recipient** is the distinguished name of the certificate holder and the **sender** of the **cr** message; and
- **protectionAlg** is the algorithm identifier for the signature algorithm used to protect the message.

If a **transactionID** was supplied in **cr** message, the header of the response will include the same **transactionID**. If a **senderNonce** was supplied in the **cr** message, the header of the response shall

include it as **recipNonce**.

The **PKIBody** is the element **cp** and is of type **CertRepMessage**. If the CA issued a certificate, the body will contain the following information:

- status will be granted or grantedWithMods;
- certificate will contain the new X.509 version 3 certificate; and
- **certifiedKeyPair** will be present.

The **certifiedKeyPair** will include the certificate in the **certOrEncCert** field and the encrypted private key in the **privateKey** field. Where the requester supplied a temporary Diffie-Hellman or elliptic curve Diffie-Hellman public key, the **privateKey** field shall include **symmAlg** and **encValue**:

- symmAlg shall contain the object identifier tDEA-ecb and the keyingOption (in ECBParms) shall be option-2 (indicating two independent keys, K1 and K2, with K3 = K1);
- encValue shall contain the encrypted private key, encrypted with tDEA in ecb mode using the two key option. The encrypted data shall be encoded as the value of the BIT STRING where the most significant bit of the BIT STRING is the most significant bit of the encrypted data. The symmetric key used to perform the encryption shall be derived from the key agreement algorithm.
  - For Diffie-Hellman keys, the shared secret shall be generated using the dHStatic mode. The tDEA key shall be derived from the shared secret using the "Key Derivation Method Based on ASN.1" specified in [X9.42]. The key derivation input AlgorithmID shall be the tDEA-ecb object identifier. The optional key derivation inputs *PartyUInfo*, *PartyVInfo*, *SuppPrivInfo*, and *SuppPubInfo* shall be omitted.
  - For elliptic curve Diffie-Hellman, the shared secret shall be generated using the "Static Unified Model Scheme", as specified in [X9.63]. The tDEA key shall be derived using the key derivation method specified in [X9.63] for the "Static Unified Model Scheme"; the optional key derivation input *SharedInfo* shall be omitted.

Where the requester supplied a temporary RSA public key, the **privateKey** field shall include **symmAlg**, **encSymmKey**, and **encValue**:

- the symmAlg shall contain the OID tDEA-ecb and the keyingOption shall be option-2 (indicating two independent keys, K1 and K2, with K3 = K1);
- **encSymmKey** shall contain the symmetric key generated by the CA, encrypted with the public key in the certificate request;
  - the symmetric key material (K1 and K2) shall be concatenated into *TwoKeys* as described in (sec. 3.1.3.5)
  - *TwoKeys* shall be encrypted using the RSA encryption algorithm according to [PKCS1], and
  - the result shall be encoded in **encSymmKey**. The encrypted data shall be encoded as the value of the **BIT STRING** where the most significant bit of the **BIT STRING** is the

most significant bit of the encrypted data

• **encValue** shall contain the encrypted private key generated by the CA, encrypted with tDEA in the electronic code book mode using the symmetric key in **encSymmKey**. The encrypted data shall be encoded as the value of the **BIT STRING** where the most significant bit of the **BIT STRING** is the most significant bit of the encrypted data.

The certificate shall contain the following **extensions**:

- a subjectKeyIdentifier field;
- at least one certificate policy OID in the **certificatePolicies** field; and
- an authority key identifier including a **Keyldentifier** field.

The CA shall use the 96-bit SHA-1 hash of the subject public key as the **subjectKeyldentifier**. The hash shall be calculated over the value (excluding tag and length) of the subject public key field in the certificate.

If the **cr** message included **extensions**, the CA may modify or ignore the requested **extensions**.

The certificate shall include URLs in the **issuerAltName** extension and **distributionPoint** field of the **CRLDistributionPoints** extension if the issuer's certificates or CRLs are not available from a well-known X.500 directory.

#### The failinfo field may not be present if status is granted or grantedWithMods.

If the CA rejected the request, the body shall include the following information:

- status will be rejected; and
- failinfo will contain the appropriate failure codes:
  - **badAlg** indicates that the CA cannot validate the signature because the algorithm identifier is unrecognized or unsupported;
  - badMessageCheck indicates that the signature in the PKIProtection field was checked but did not match;
  - **badRequest** indicates that the responder does not permit or support the transaction;
  - badTime indicates that the messageTime field in the message header was not sufficiently close to the responder's system time; and
  - badCertId indicates that no current valid, unrevoked signature certificate could be identified for the requester.

The certificate field may not be present if status is rejected.

The **PKIProtection** field contains the CA's signature, calculated on the DER encoded sequence of the header and body.

• Where the requester supplied a temporary Diffie-Hellman or ECDH public key, the **extraCerts** field shall contain the CA's certificate with the public key from the key pair used to perform key agreement.

#### **Confirmation Message**

Upon receipt of the **cp** message, the certificate holder shall generate a **PKIConfirm** message.

**PKIHeader** data shall be identical to the certificate request from the RA to the CA, with the exception of **messageTime**.

The **PKIProtection** field contains the certificate holder's signature, calculated on the DER encoded sequence of the header and body using the private key corresponding to the currently valid signature certificate.

## 3.5.11 Combined Certificate Requests

Certificate requests for signature and key management keys may be combined into a single transaction. Specifically, the RA-generated and self registration requests (secs. 3.5.1, 3.5.3, and 3.5.4) may be combined with the encryption certificate requests (secs. 3.5.7 and 3.5.10). In these cases, **CertReqMessages** will be a **SEQUENCE** of length two. One of the **CertReqMessage** shall be that described in RA-generated or self registration requests; the other shall conform to the **CertReqMessage** described for the encryption certificate request. The messages shall be protected as described in the signature certificate requests.

If the combined request was a self-registration request, the signature certificate request must be approved or both certificate requests must be rejected. If additional messages are required for proof of possession, the requester signs the challenge response message with the signature key in the self registration request.

Proof of possession for the signature key will be implemented using the **pop** field as described in (secs. 3.5.1, 3.5.3, and 3.5.4). If the end entity generated the key management key pair and proof of possession is required for the private key, it shall be performed as described in 3.5.7. If a centrally generated key pair was requested, it shall be encrypted as described in 3.5.10. The CA generated responses shall be those described in the appropriate sections. These responses may be combined into a single **CertRepMessage**, or be transmitted separately. The requester shall use the **CertReqID** to differentiate the responses.

**PKIFreeText** may be used in the response to supply additional information.

Upon receipt of the **cp** message, the certificate holder shall generate a **PKIConfirm** message. **PKIHeader** data shall be identical to the certificate request from the RA to the CA, with the exception of **messageTime**.

If the signature certificate request was accepted, the **PKIProtection** field contains the certificate holder's signature, calculated on the DER encoded sequence of the header and body using the private key corresponding to the new signature certificate. If the request was rejected, the **PKIProtection** field contains the mac, calculated on the DER encoded sequence of the header and body using the shared secret that authenticated the certificate request. (If the request was rejected for **badMessageCheck**, the **PKIConfirm** message need not be generated.)

## 3.5.12 Request Certificate from a Repository

Entities may request certificates from a repository using LDAP V2 as defined in [RFC 2559]. When using LDAP, the entity may request certificates from a repository service using the LDAP search request, as defined in [RFC 2559] or as specified in a given LDAP URL [RFC1959] (e.g., the **authorityInformationAccess** extension.)

## 3.5.13 Request CRL from a Repository

Entities may request CRLs from a repository using LDAP V2 as defined in [RFC 2559]. Entities may request CRLs from a repository using LDAP [RFC 1777]. When using LDAP, the entity may request CRLs from a repository service using the LDAP V2 search request, as defined in [RFC 2559] and [RFC 1777] or as specified in a given LDAP URL [RFC1959] (e.g., the **distributionPoint** field in the **cRLDistributionPoints** extension.)

## 4. References

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[STAB95]	OIW, Stable Implementation Agreements for Open Systems Interconnection Protocols: Part 12 - OS Security. June 1995.
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# Appendix A. X.509 v3 Certificate ASN.1

#### AuthenticationFramework {joint-iso-ccitt ds(5) modules(1) authenticationFramework(7) 2} DEFINITIONS ::= BEGIN

#### -- EXPORTS All --

-- The types and values defined in this module are exported for use in the other ASN.1

-- modules contained within the Directory Specifications, and for the use of other applications

-- which will use them to access Directory services. Other applications may use them for

-- their own purposes, but this will not constrain extensions and modifications needed to

-- maintain or improve the Directory service.

#### **IMPORTS**

-- types --

Certificate ::=	SIGNE	) {SEQL	JENCE{
version		[0] Version DEFAULT v1,	
serialNumber		CertificateSerialNumber,	
signature		Algorithmldentifier,	
issuer		Name,	
validity		Validity,	
subject		Name,	
subjectPublicKeyInfo	D		SubjectPublicKeyInfo}
issuerUniqueIdentific	er	[1]	IMPLICIT UniqueIdentifier OPTIONAL,
·			if present, version must be v1 or v2
subjectUniqueIdentif	fier	[2]	IMPLICIT UniqueIdentifier OPTIONAL,
			if present, version must be v1 or v2
extensions		[3]	Extensions OPTIONAL
			if present, version must be v3 }
Version	::=	INTEGE	ER {v1(0), v2(1), v3(2) }
CertificateSerialNumber	::=	INTEGE	ER
Algorithmidentifier	::=	SEQUE	NCE{
algorithm	ALGOR	ALGORITHM.&id({SupportedAlgorithms}),	
parameters	ALGOR	ITHM.&	Type ({SupportedAlgorithms}{ @algorithm}) OPTIONAL }
-			

-- Definition of the following information object is deferred, perhaps to standardized

-- profiles of to protocol implementation conformance statements. This set is required to

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- specify a table constraint on the **Parameters** component of **Algorithmidentifier**. **SupportedAlgorithms ALGORITHM** ::= { ... } ---

Validity notBefore notAfter		::= SEQUENCE{ eOfTime, eOfTime }			
ChoiceOfTime ::= CHOICE { utcTime UTCTime, generalTime GeneralizedTime }		ne }			
algorithm Alg			lgorithmldentifier,		
Extensions	::=	SEQUENCE OF	Extension		
Extension extnld critical extnValue		SEQUENCE { NSION.&id ({ExtensionSet}), EAN DEFAULT FALSE, OCTET STRING contains a DER encoding of a value of type &ExtnType for the extension object identified by extnld			
<ul> <li>Definition of the following information object set is deferred, perhaps to</li> <li>standardized profiles or to protocol implementation conformance statements.</li> <li>The set is required to specify a table constraint on the critical component</li> <li>of Extension.</li> <li>ExtensionSet EXTENSION ::= {}</li> </ul>					
ExtensionSet EXTENSION ::= {   } EXTENSION ::= CLASS					
{ &id OBJECT IDENTIFIER UNIQUE, &ExtnType					
} WITH SYNTAX					
{ SYNTAX IDENTIFIED F }	BY	&ExtnType &id			
Certificates certificate certificationF	::= Path	SEQUE Certific Forward	•		
ForwardCertificationPath ::= SEQUENCE OF CrossCertificates		NCE OF CrossCertificates			
CertificationPath ::= userCertificate theCACertificates		Certific	•		

SIGNED { SEQUENCE { CertificateList ::= Version OPTIONAL, -- if present, must be v2 version signature AlgorithmIdentifier, issuer Name, thisUpdate ChoiceOfTime, nextUpdate ChoiceOfTime OPTIONAL, revokedCertificates **SEQUENCE OF SEQUENCE {** userCertificate CertificateSerialNumber, revocationDate ChoiceOfTime, Extensions OPTIONAL } OPTIONAL, crlEntryExtensions crlExtensions Extensions OPTIONAL }} [0] **CertificatePair SEQUENCE {** ::= Certificate OPTIONAL, forward [0] **Certificate OPTIONAL** reverse [1] -- at least one of the pair shall be present -- } -- attribute types-userPassword **ATTRIBUTE** ::= { WITH SYNTAX OCTET STRING (SIZE (0..ub-user-password)) EQUALITY MATCHING RULE octetStringMatch id-at-userPassword } ID userCertificate ATTRIBUTE ::= { WITH SYNTAX Certificate ID id-at-userCertificate } ATTRIBUTE cACertificate ::= { WITH SYNTAX Certificate ID id-at-cACertificate } authorityRevocationList ATTRIBUTE ::= { WITH SYNTAX CertificateList ID id-at-authorityRevocationList } certificateRevocationList ATTRIBUTE ::= { WITH SYNTAX CertificateList ID id-at-certificateRevocationList } crossCertificatePair ATTRIBUTE ::= { WITH SYNTAX CertificatePair ID id-at-crossCertificatePair } -- information object classes --**TYPE-IDENTIFIER** ALGORITHM ::= -- Parameterized Types --HASHED {ToBeHashed} **OCTET STRING ( CONSTRAINED-BY {** ::= --must be the result of applying a hashing procedure to the ----DER-encoded octets of a value of -- ToBeHashed }) ENCRYPTED { To\BeEnciphered} BIT STRING (CONSTRAINED BY { :=

--must be the result of applying an encipherment procedure to the ----BER-encoded octets of a value of -- **ToBeEnciphered })** 

- SIGNED { ToBeSigned } ::= SEQUENCE{ ToBeSigned, COMPONENTS OF SIGNATURE { ToBeSigned }),
- SIGNATURE { OfSignature } ::= SEQUENCE { AlgorithmIdentifier, ENCRYPTED { HASHED { OfSignature }}}

-- object identifier assignments --

id-at-userPassword	<b>OBJECT IDENTIFIER</b>	::=	{id-at 35}
id-at-userCertificate	<b>OBJECT IDENTIFIER</b>	::=	{id-at 36}
id-at-cAcertificate	<b>OBJECT IDENTIFIER</b>	::=	{id-at 37}
id-at-authorityRevocationList	<b>OBJECT IDENTIFIER</b>	::=	{id-at 38}
id-at-certificateRevocationList	<b>OBJECT IDENTIFIER</b>	::=	{id-at 39}
id-at-crossCertificatePair	<b>OBJECT IDENTIFIER</b>	::=	{id-at 40}
id-at-supportedAlgorithms	<b>OBJECT IDENTIFIER</b>	::=	{id-at 52}
id-at-deltaRevocationList	OBJECT IDENTIFIER	::=	{id-at 53}

END

# Appendix B. Certificate and CRL Extensions ASN.1

CertificateExtensions {joint-iso-ccitt ds(5) module(1) certificateExtensions(26) 0} DEFINITIONS IMPLICIT TAGS ::= BEGIN

-- EXPORTS ALL --

IMPORTS				
	ionFramework, authenticationFra	mework		
selectedAttributeTy				
	tions {joint-iso-ccitt ds(5) module	e(1)		
usefulDefinitions(0)				
	edName, ATTRIBUTE, Attribute,			
	ROM InformationFramework info ertificateList, AlgorithmIdentifier,			
EXTENSION				
FROM Authentication	onFramework authenticationFran	nework		
DirectoryString				
	ibuteTypes selectedAttributeType	es		
ub-name	le unner Deun de			
FROM UpperBound ORAddress	is upperbounds			
	tService {joint-iso-ccitt mhs(6) m	ts(3)		
	stract-service(1) version-1994 (0)			
IDENTIFIED BY { id- AuthorityKeyldentifier ::= SEQUEN keyldentifier authorityCertIssuer authorityCertIssuer	F construct in this specification. ons ::= { norityKeyIdentifier ce 35 } } ICE { [0] KeyIdentifier [1] GeneralNames [2] CertificateSerialNumber	OPTIONAL, OPTIONAL, OPTIONAL }		
(WITH COMPONENTS	{, authorityCertIssuer PRES authorityCertSerialNumber PI			
WITH COMPONENTS {, authorityCertIssuer ABSENT, authorityCertSerialNumber ABSENT} )				
Keyldentifier ::= OCTET STRING				
subjectKeyIdentifier EXTENSION ::= {				
SYNTAX SubjectKeyIdentifier				
IDENTIFIED BY { id-	ce 14 } }			
SubjectKeyldentifier ::= Keyldentifi	ier			

keyUsage EXTENSION ::= {

SYNTAX KeyUsage IDENTIFIED BY { id-ce 15 } }				
KeyUsage ::= BIT STRING {digitalSignature(0),nonRepudiation(1),keyEncipherment(2),dataEncipherment(3),keyAgreement(4),keyCertSign(5),cRLSign(6) }				
privateKeyUsagePeriod EXTENSION ::= { SYNTAX PrivateKeyUsagePeriod IDENTIFIED BY { id-ce 16 } }				
PrivateKeyUsagePeriod ::= SEQUENCE { notBefore [0] GeneralizedTime OPTIONAL, notAfter [1] GeneralizedTime OPTIONAL } (WITH COMPONENTS {, notBefore PRESENT}   WITH COMPONENTS {, notAfter PRESENT} )				
certificatePolicies EXTENSION ::= { SYNTAX CertificatePoliciesSyntax IDENTIFIED BY { id-ce 32 } }				
CertificatePoliciesSyntax ::= SEQUENCE SIZE (1MAX) OF PolicyInformation				
PolicyInformation ::= SEQUENCE { policyIdentifier CertPolicyId, policyQualifiers SEQUENCE SIZE (1MAX) OF PolicyQualifierInfo OPTIONAL }				
CertPolicyId ::= OBJECT IDENTIFIER				
PolicyQualifierInfo ::= SEQUENCE { policyQualifierId CERT-POLICY-QUALIFIER.&id				
SupportedPolicyQualifiers CERT-POLICY-QUALIFIER ::= { }				
CERT-POLICY-QUALIFIER ::= CLASS { &id OBJECT IDENTIFIER UNIQUE, &Qualifier OPTIONAL } WITH SYNTAX { POLICY-QUALIFIER-ID&id [QUALIFIER-TYPE &Qualifier] }				
policyMappings EXTENSION ::= { SYNTAX PolicyMappingsSyntax IDENTIFIED BY { id-ce 33 } }				

PolicyMappingsSyntax ::= SEQUENCE SIZE (1..MAX) OF SEQUENCE { issuerDomainPolicy CertPolicyId, subjectDomainPolicy CertPolicyId }

supportedAlgorithms ATTRIBUTE ::= { WITH SYNTAX SupportedAlgorithm EQUALITY MATCHING RULE algorithmIdentifierMatch ID { id-at 52 } }

SupportedAlgorithm ::= SEQUENCE { algorithmIdentifier AlgorithmIdentifier, intendedUsage [0] KeyUsage OPTIONAL, intendedCertificatePolicies [1] CertificatePoliciesSyntax OPTIONAL }

-- Certificate subject and certificate issuer attributes extensions --

subjectAltName EXTENSION ::= { SYNTAX GeneralNames IDENTIFIED BY { id-ce 17 } }

GeneralNames ::= SEQUENCE SIZE (1..MAX) OF GeneralName

GeneralName ::= CHOICE {

otherName rfc822Name dNSName x400Address directoryName ediPartyName uniformResourceIdentifier iPAddress registeredID	[0] [1] [3] [4] [5] [6] [7] [8]	INSTANCE OF OTHER-NAME, IA5String, IA5String, ORAddress, Name, EDIPartyName, IA5String, OCTET STRING, OBJECT IDENTIFIER }		
OTHER-NAME ::= TYPE-IDENTIFIER		-		
EDIPartyName ::= SEQUENCE { nameAssigner partyName	[0] [1]	DirectoryString {ub-name} OPTIONAL, DirectoryString {ub-name} }		
issuerAltName EXTENSION ::= { SYNTAX GeneralNames IDENTIFIED BY { id-ce 18 } }				
subjectDirectoryAttributes EXTENSION ::= { SYNTAX AttributesSyntax IDENTIFIED BY { id-ce 9 } }				
AttributesSyntax ::= SEQUENCE SIZE (1MAX) OF Attribute				
Certification path constraints extension	ons			
basicConstraints EXTENSION ::= {				

SYNTAX **BasicConstraintsSyntax** IDENTIFIED BY { id-ce 19 } } BasicConstraintsSyntax ::= SEQUENCE { **BOOLEAN DEFAULT FALSE,** сА pathLenConstraint INTEGER (0..MAX) OPTIONAL } nameConstraints EXTENSION ::= { NameConstraintsSyntax SYNTAX IDENTIFIED BY { id-ce 30 } } NameConstraintsSyntax ::= SEQUENCE { permittedSubtrees [0] GeneralSubtrees OPTIONAL, excludedSubtrees [1] GeneralSubtrees OPTIONAL } GeneralSubtrees ::= SEQUENCE SIZE (1..MAX) OF GeneralSubtree GeneralSubtree ::= SEQUENCE { base GeneralName, minimum [0] BaseDistance DEFAULT 0. maximum [1] BaseDistance OPTIONAL } BaseDistance ::= INTEGER (0..MAX) policyConstraints EXTENSION ::= { SYNTAX **PolicyConstraintsSyntax** IDENTIFIED BY { id-ce 36 } } PolicyConstraints Syntax ::= SEQUENCE SIZE (1..MAX) OF SEQUENCE { requireExplicitPolicy [0] SkipCerts OPTIONAL, inhibitPolicyMapping [1] SkipCerts OPTIONAL } SkipCerts ::= INTEGER (0..MAX) -- Basic CRL extensions -cRLNumber EXTENSION ::= { SYNTAX CRLNumber IDENTIFIED BY { id-ce 20 } } CRLNumber ::= INTEGER (0..MAX) reasonCode EXTENSION ::= { SYNTAX CRLReason IDENTIFIED BY { id-ce 21 } } CRLReason ::= ENUMERATED { (0), unspecified keyCompromise (1), cACompromise (2), affiliationChanged (3), superseded (4), cessationOfOperation (5), certificateHold (6),

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removeFromCRL (8) }

instructionCode EXTENSION ::= { SYNTAX HoldInstruction IDENTIFIED BY { id-ce 23 } }

HoldInstruction ::= OBJECT IDENTIFIER

invalidityDate EXTENSION ::= { SYNTAX GeneralizedTime IDENTIFIED BY { id-ce 24 } }

-- CRL distribution points and delta-CRL extensions --

cRLDistributionPoints EXTENSION ::= {
 SYNTAX CRLDistPointsSyntax
 IDENTIFIED BY { id-ce 31 } }

CRLDistPointsSyntax ::= SEQUENCE SIZE (1..MAX) OF DistributionPoint

DistributionPoint ::= SEQUEN distributionPoint reasons cRLIssuer	ICE { [0] [1] [2]	DistributionPointName OPTIONAL, ReasonFlags OPTIONAL, GeneralNames OPTIONAL }				
DistributionPointName ::= CHOICE {						
fullName		[0]	GeneralNames,			
nameRelativeToCRLIssuer		[1]	RelativeDistinguishedName }			
ReasonFlags ::= BIT STRING {						
unused		(0),				
keyCompromise			(1),			
caCompromise		(2),				
affiliationChanged		(3),				
superseded		(4),				
cessationOfOperation		(5),				
certificateHold		(6) }				
issuingDistributionPoint EXTENSION ::= {						
SYNTAX IssuingDistPointSyntax						
IDENTIFIED BY { id-ce 28 } }						
IssuingDistPointSyntax ::= SEQUENCE { distributionPoint [0] DistributionPointName OPTIONAL,						
onlyContainsUserCerts		[1] BOOLEAN DEFAULT FALSE,				
onlyContainsCACerts		[2] BOOLEAN DEFAULT FALSE,				
onlySomeReasons		[3] ReasonFlags OPTIONAL,				
indirectCRL		[4] BO	OLEAN DEFAULT FALSE }			
certificatelssuer EXTENSION ::= { SYNTAX GeneralNames IDENTIFIED BY { id-ce 29 } }						

deltaCRLIndicator EXTENSION ::= { SYNTAX BaseCRLNumber IDENTIFIED BY { id-ce 27 } } BaseCRLNumber ::= CRLNumber deltaRevocationList ATTRIBUTE ::= {						
WITH SYNTAX CertificateList EQUALITY MATCHING RULE certificateListExactMatch ID {id-at 53 } }						
Matching rules						
certificateExactMatch MATCHING-RULE ::= { SYNTAX CertificateExactAssertion ID id-mr-certificateExactMatch }						
CertificateExactAssertion ::= SEQUENCE { serialNumber CertificateSerialNumber, issuer Name }						
certificateMatch MATCHING- SYNTAX ID	RULE ::= { CertificateAssertion id-mr-certificateMatch }					
CertificateAssertion ::= SEQU serialNumber issuer subjectKeyldentifier authorityKeyldentifier certificateValid privateKeyValid subjectPublicKeyAlgl keyUsage subjectAltName policy pathToName	[0] CertificateSerialNumber O [1] Name [2] SubjectKeyldentifier [3] AuthorityKeyldentifier [4] UTCTime [5] GeneralizedTime	PTIONAL, OPTIONAL, OPTIONAL, OPTIONAL, OPTIONAL, OPTIONAL, OPTIONAL, OPTIONAL, OPTIONAL, OPTIONAL,				
AltNameType ::= CHOICE { builtinNameForm otherNameForm	ENUMERATED { rfc822Name (1), dNSName (2), x400Address (3), directoryName (4), ediPartyName (5), uniformResourceIdentifier iPAddress (7), registeredId (8) }, OBJECT IDENTIFIER }	(6),				
certificatePairExactMatch MATCHING-RULE ::= {     SYNTAX CertificatePairExactAssertion						

ID id-mr-certificatePairExactMatch } CertificatePairExactAssertion ::= SEQUENCE { [0] CertificateExactAssertion OPTIONAL, forwardAssertion [1] CertificateExactAssertion OPTIONAL } reverseAssertion (WITH COMPONENTS {..., forwardAssertion PRESENT} | WITH COMPONENTS {..., reverseAssertion PRESENT} ) certificatePairMatch MATCHING-RULE ::= { SYNTAX **CertificatePairAssertion** ID id-mr-certificatePairMatch } CertificatePairAssertion ::= SEQUENCE { forwardAssertion [0] CertificateAssertion OPTIONAL, [1] CertificateAssertion OPTIONAL } reverseAssertion (WITH COMPONENTS {..., forwardAssertion PRESENT} | WITH COMPONENTS {..., reverseAssertion PRESENT} ) certificateListExactMatch MATCHING-RULE ::= { SYNTAX CertificateListExactAssertion ID id-mr-certificateListExactMatch } CertificateListExactAssertion ::= SEQUENCE { issuer Name. UTCTime. thisUpdate distributionPoint DistributionPointName OPTIONAL } certificateListMatch MATCHING-RULE ::= { CertificateListAssertion SYNTAX ID id-mr-certificateListMatch } CertificateListAssertion ::= SEQUENCE { Name OPTIONAL, issuer minCRLNumber [0] CRLNumber OPTIONAL, maxCRLNumber CRLNumber OPTIONAL, [1] reasonFlags ReasonFlags OPTIONAL, dateAndTime UTCTime OPTIONAL, distributionPoint DistributionPointName OPTIONAL } [2] algorithmIdentifierMatch MATCHING-RULE ::= { SYNTAX AlgorithmIdentifier ID id-mr-algorithmldentifierMatch } -- Object identifier assignments -id-at-supportedAlgorithms **OBJECT IDENTIFIER ::=** {id-at 52} id-at-deltaRevocationList **OBJECT IDENTIFIER** ::= {id-at 53} id-ce-subjectDirectoryAttributes **OBJECT IDENTIFIER** ::= {id-ce 9} id-ce-subjectKeyIdentifier OBJECT IDENTIFIER {id-ce 14} ::=

id-ce-basicConstraints	OBJECT IDENTIFIER	::=	(id oo 10)
			{id-ce 19}
id-ce-cRLNumber	OBJECT IDENTIFIER	::=	{id-ce 20}
id-ce-reasonCode	OBJECT IDENTIFIER	::=	{id-ce 21}
id-ce-instructionCode	<b>OBJECT IDENTIFIER</b>	::=	{id-ce 23}
id-ce-invalidityDate	OBJECT IDENTIFIER	::=	{id-ce 24}
id-ce-deltaCRLIndicator	<b>OBJECT IDENTIFIER</b>	::=	{id-ce 27}
id-ce-issuingDistributionPoint	<b>OBJECT IDENTIFIER</b>	::=	{id-ce 28}
id-ce-certificatelssuer	<b>OBJECT IDENTIFIER</b>	::=	{id-ce 29}
id-ce-nameConstraints	<b>OBJECT IDENTIFIER</b>	::=	{id-ce 30}
id-ce-cRLDistributionPoints	<b>OBJECT IDENTIFIER</b>	::=	{id-ce 31}
id-ce-certificatePolicies	<b>OBJECT IDENTIFIER</b>	::=	{id-ce 32}
id-ce-policyMappings	<b>OBJECT IDENTIFIER</b>	::=	{id-ce 33}
id-ce-policyConstraints	<b>OBJECT IDENTIFIER</b>	::=	{id-ce 34}
id-ce-authorityKeyIdentifier	<b>OBJECT IDENTIFIER</b>	::=	{id-ce 35}
id-mr-certificateExactMatch	<b>OBJECT IDENTIFIER</b>	::=	{id-mr 34}
id-mr-certificateMatch	<b>OBJECT IDENTIFIER</b>	::=	{id-mr 35}
id-mr-certificatePairExactMatch	<b>OBJECT IDENTIFIER</b>	::=	{id-mr 36}
id-mr-certificatePairMatch	<b>OBJECT IDENTIFIER</b>	::=	{id-mr 37}
id-mr-certificateListExactMatch	<b>OBJECT IDENTIFIER</b>	::=	{id-mr 38}
id-mr-certificateListMatch	<b>OBJECT IDENTIFIER</b>	::=	{id-mr 39}
id-mr-algorithmIdentifierMatch	<b>OBJECT IDENTIFIER</b>	::=	{id-mr 40}

--- The following OBJECT IDENTIFIERS are not used by this specification: -- {id-ce 2}, {id-ce 3}, {id-ce 4}, {id-ce 5}, {id-ce 6}, {id-ce 7}, -- {id-ce 8}, {id-ce 10}, {id-ce 11}, {id-ce 12}, {id-ce 13}, -- {id-ce 22}, {id-ce 25}, {id-ce 26}

END

## Appendix C. ASN.1 Module for transactions

The following section contains the complete ASN.1 module from RFC 2510, the *Certificate Management Protocol*. Only a small subset of the messages defined in [RFC2510] are required to implement this specification. The entire module is provided for completeness. Information about messages defined by this ASN.1 module but not used in the MISPC may be found in [RFC2510].

```
PKIX-CMP DEFINITIONS ::=
BEGIN -- EXPLICT TAGS
IMPORTS
       CertRegMessages, EncryptedValue, EncryptedKey
       FROM CMRF
 PKIMessage ::= SEQUENCE {
   Header
              PKIHeader,
              PKIBody,
   Body
   Protection [0] PKIProtection OPTIONAL,
   ExtraCerts [1] SEQUENCE SIZE (1..MAX) OF Certificate OPTIONAL
}
 PKIHeader ::= SEQUENCE {
   pvno
                 INTEGER { ietf-version2 (1) },
   sender
                 GeneralName,
   -- identifies the sender
   recipient
                 GeneralName.
   -- identifies the intended recipient
   messageTime [0] GeneralizedTime
                                           OPTIONAL,
   -- time of production of this message (used when sender
   -- believes that the transport will be "suitable"; i.e.,
   -- that the time will still be meaningful upon receipt)
                     [1] AlgorithmIdentifier OPTIONAL.
   protectionAlg
   -- algorithm used for calculation of protection bits
                     [2] Keyldentifier
   senderKID
                                           OPTIONAL,
   recipKID
                     [3] Keyldentifier
                                           OPTIONAL,
   -- to identify specific keys used for protection
   transactionID
                     [4] OCTET STRING
                                              OPTIONAL,
   -- identifies the transaction: i.e., this will be the same in
   -- corresponding request, response and confirmation messages
   senderNonce
                     [5] OCTET STRING
                                              OPTIONAL,
   recipNonce
                     [6] OCTET STRING
                                              OPTIONAL,
   -- nonces used to provide replay protection, senderNonce
   -- is inserted by the creator of this message; recipNonce
   -- is a nonce previously inserted in a related message by
   -- the intended recipient of this message
   freeText
                     [7] PKIFreeText
                                            OPTIONAL,
   -- this may be used to indicate context-specific instructions
   -- (this field is intended for human consumption)
   generalInfo
                     [8] SEQUENCE SIZE (1..MAX) OF
                     InfoTypeAndValue
                                          OPTIONAL
   -- this may be used to convey context-specific information
```

-- (this field not primarily intended for human consumption) }

```
PKIFreeText ::= SEQUENCE SIZE (1..MAX) OF UTF8String

-- text encoded as UTF-8 String (note: each UTF8String SHOULD

-- include an RFC 1766 language tag to indicate the language

-- of the contained text)
```

```
PKIBody ::= CHOICE {
                        -- message-specific body elements
              [0] CertRegMessages,
                                       --Initialization Request
   ir
   ip
              [1] CertRepMessage,
                                       --Initialization Response
              [2] CertReqMessages,
                                       --Certification Request
   cr
              [3] CertRepMessage,
                                       --Certification Response
   ср
   p10cr
              [4] CertificationReguest, --imported from [PKCS10]
              [5] POPODecKeyChallContent, --pop Challenge
   popdecc
              [6] POPODecKeyRespContent, --pop Response
   popdecr
                                       --Key Update Request
   kur
              [7] CertReqMessages,
              [8] CertRepMessage,
                                       --Key Update Response
   kup
   krr
              [9] CertReqMessages,
                                       --Key Recovery Request
              [10] KevRecRepContent.
                                         --Kev Recovery Response
   krp
              [11] RevReqContent,
                                       --Revocation Request
   rr
              [12] RevRepContent,
                                       --Revocation Response
   rp
              [13] CertRegMessages,
                                        --Cross-Cert. Request
   ccr
              [14] CertRepMessage,
                                       --Cross-Cert. Response
   CCD
   ckuann
              [15] CAKeyUpdAnnContent,
                                           --CA Key Update Ann.
              [16] CertAnnContent,
                                       --Certificate Ann.
   cann
                                       --Revocation Ann.
              [17] RevAnnContent,
   rann
              [18] CRLAnnContent,
                                       --CRL Announcement
   crlann
   conf
              [19] PKIConfirmContent,
                                        --Confirmation
              [20] NestedMessageContent, --Nested Message
   nested
   genm
              [21] GenMsgContent,
                                       --General Message
              [22] GenRepContent,
                                       --General Response
   genp
              [23] ErrorMsgContent
                                       --Error Message
   error
}
 PKIProtection ::= BIT STRING
 ProtectedPart ::= SEQUENCE {
   header PKIHeader,
   bodv
           PKIBody
}
```

PasswordBasedMac ::= OBJECT IDENTIFIER --{1 2 840 113533 7 66 13}

```
PBMParameter ::= SEQUENCE {
salt OCTET STRING,
owf AlgorithmIdentifier,
-- AlgId for a One-Way Function (SHA-1 recommended)
iterationCount INTEGER,
-- number of times the OWF is applied
mac AlgorithmIdentifier
-- the mac AlgId (e.g., DES-MAC, Triple-DES-MAC [PKCS11],
} -- or HMAC [RFC2104, RFC2202])
```

```
DHBasedMac ::= OBJECT IDENTIFIER --{1 2 840 113533 7 66 30}
DHBMParameter ::= SEQUENCE {
  owf
                Algorithmldentifier,
  -- AlgId for a One-Way Function (SHA-1 recommended)
                AlgorithmIdentifier
  mac
  -- the MAC AlgId (e.g., DES-MAC, Triple-DES-MAC [PKCS11],
} -- or HMAC [RFC2104, RFC2202])
NestedMessageContent ::= PKIMessage
PKIStatus ::= INTEGER {
  granted
                             (0),
  -- you got exactly what you asked for
  grantedWithMods
                             (1),
  -- you got something like what you asked for; the
  -- requester is responsible for ascertaining the differences
  rejection
                             (2),
  -- you don't get it, more information elsewhere in the message
  waiting
                             (3),
  -- the request body part has not yet been processed,
  -- expect to hear more later
  revocationWarning
                             (4),
  -- this message contains a warning that a revocation is
  -- imminent
  revocationNotification
                             (5),
  -- notification that a revocation has occurred
  keyUpdateWarning
                             (6)
  -- update already done for the oldCertId specified in
  -- CertReqMsg
}
PKIFailureInfo ::= BIT STRING {
-- since we can fail in more than one way!
-- More codes may be added in the future if/when required.
  badAlg
                             (0),
  -- unrecognized or unsupported Algorithm Identifier
  badMessageCheck
                             (1),
  -- integrity check failed (e.g., signature did not verify)
  badRequest
                             (2),
  -- transaction not permitted or supported
  badTime
                             (3),
  -- messageTime was not sufficiently close to the system time,
  -- as defined by local policy
  badCertId
                             (4),
  -- no certificate could be found matching the provided criteria
  badDataFormat
                             (5),
  -- the data submitted has the wrong format
  wrongAuthority
                             (6),
  -- the authority indicated in the request is different from the
  -- one creating the response token
  incorrectData
                             (7),
  -- the requester's data is incorrect (for notary services)
```

```
missingTimeStamp
                            (8),
  -- when the timestamp is missing but should be there (by policy)
  badPoP
                            (9)
  -- when proof of possession does not verify
}
PKIStatusInfo ::= SEQUENCE {
  status
            PKIStatus,
  statusString PKIFreeText OPTIONAL,
  failInfo
            PKIFailureInfo OPTIONAL
}
OOBCert ::= Certificate
OOBCertHash ::= SEQUENCE {
           [0] AlgorithmIdentifier
                                     OPTIONAL,
  hashAlg
           [1] CertId
                              OPTIONAL,
  certId
               BIT STRING
  hashVal
  -- hashVal is calculated over DER encoding of the
  -- subjectPublicKey field of the corresponding cert.
}
POPODecKeyChallContent ::= SEQUENCE OF Challenge
-- One Challenge per encryption key certification request (in the
-- same order as these requests appear in CertReqMessages).
Challenge ::= SEQUENCE {
               AlgorithmIdentifier OPTIONAL,
  owf
  -- MUST be present in the first Challenge; MAY be omitted in any
  -- subsequent Challenge in POPODecKeyChallContent (if omitted.
  -- then the owf used in the immediately preceding Challenge is
  -- to be used).
                 OCTET STRING,
  witness
  -- the result of applying the one-way function (owf) to a
  -- randomly-generated INTEGER, A. [Note that a different
  -- INTEGER MUST be used for each Challenge.]
  challenge
                 OCTET STRING
  -- the encryption (under the public key for which the cert.
  -- request is being made) of Rand, where Rand is specified as
  -- Rand ::= SEQUENCE {
       int
             INTEGER,
  --
       - the randomly-generated INTEGER A (above)
  --
       sender GeneralName
  --
       - the sender's name (as included in PKIHeader)
  --
  -- }
}
POPODecKeyRespContent ::= SEQUENCE OF INTEGER
-- One INTEGER per encryption key certification request (in the
```

- -- same order as these requests appear in CertReqMessages). The
- -- retrieved INTEGER A (above) is returned to the sender of the
- -- corresponding Challenge.

```
CertRepMessage ::= SEQUENCE {
```

```
caPubs
             [1] SEQUENCE SIZE (1..MAX) OF Certificate OPTIONAL,
  response
                SEQUENCE OF CertResponse
}
CertResponse ::= SEQUENCE {
  certRegId
                 INTEGER,
  -- to match this response with corresponding request (a value
  -- of -1 is to be used if certRegId is not specified in the
  -- corresponding request)
  status
                PKIStatusInfo,
  certifiedKeyPair CertifiedKeyPair OPTIONAL,
  rspinfo
                OCTET STRING
                                   OPTIONAL
  -- analogous to the id-regInfo-asciiPairs OCTET STRING defined
  -- for reginfo in CertRegMsg [CRMF]
}
CertifiedKeyPair ::= SEQUENCE {
  certOrEncCert
                   CertOrEncCert.
  privateKey
               [0] EncryptedValue
                                     OPTIONAL,
  publicationInfo [1] PKIPublicationInfo OPTIONAL
}
CertOrEncCert ::= CHOICE {
  certificate [0] Certificate,
  encryptedCert [1] EncryptedValue
}
KeyRecRepContent ::= SEQUENCE {
  status
                  PKIStatusInfo,
  newSigCert
                  [0] Certificate
                                          OPTIONAL.
                [1] SEQUENCE SIZE (1..MAX) OF
  caCerts
                      Certificate
                                   OPTIONAL,
                  [2] SEQUENCE SIZE (1..MAX) OF
  keyPairHist
                      CertifiedKeyPair OPTIONAL
}
RevReqContent ::= SEQUENCE OF RevDetails
RevDetails ::= SEQUENCE {
  certDetails
                 CertTemplate,
  -- allows requester to specify as much as they can about
  -- the cert. for which revocation is requested
  -- (e.g., for cases in which serialNumber is not available)
  revocationReason ReasonFlags
                                     OPTIONAL,
  -- the reason that revocation is requested
  badSinceDate
                   GeneralizedTime OPTIONAL,
  -- indicates best knowledge of sender
  crlEntryDetails
                   Extensions
                                 OPTIONAL
  -- requested crlEntryExtensions
}
RevRepContent ::= SEQUENCE {
            SEQUENCE SIZE (1..MAX) OF PKIStatusInfo,
  status
```

```
-- in same order as was sent in RevReqContent
```

```
revCerts [0] SEQUENCE SIZE (1..MAX) OF CertId OPTIONAL,
   -- IDs for which revocation was requested (same order as status)
         [1] SEQUENCE SIZE (1..MAX) OF CertificateList OPTIONAL
   crls
   -- the resulting CRLs (there may be more than one)
 }
 CAKeyUpdAnnContent ::= SEQUENCE {
   oldWithNew
                    Certificate, -- old pub signed with new priv
   newWithOld
                    Certificate, -- new pub signed with old priv
   newWithNew
                     Certificate -- new pub signed with new priv
}
 CertAnnContent ::= Certificate
 RevAnnContent ::= SEQUENCE {
                 PKIStatus,
   status
   certId
                CertId.
   willBeRevokedAt GeneralizedTime,
   badSinceDate
                     GeneralizedTime.
                  Extensions OPTIONAL
   crIDetails
   -- extra CRL details(e.g., crl number, reason, location, etc.)
}
 CRLAnnContent ::= SEQUENCE OF CertificateList
 PKIConfirmContent ::= NULL
 InfoTypeAndValue ::= SEQUENCE {
   infoTvpe
                    OBJECT IDENTIFIER,
   infoValue
                    ANY DEFINED BY infoType OPTIONAL
}
 -- Example InfoTypeAndValue contents include, but are not limited to:
 -- { CAProtEncCert = {id-it 1}, Certificate
                                                     }
 -- { SignKeyPairTypes = {id-it 2}, SEQUENCE OF AlgorithmIdentifier }
 -- { EncKeyPairTypes = {id-it 3}, SEQUENCE OF AlgorithmIdentifier }
 -- { PreferredSymmAlg = {id-it 4}, AlgorithmIdentifier
                                                           }
 -- { CAKeyUpdateInfo = {id-it 5}, CAKeyUpdAnnContent
                                                               }
 -- { CurrentCRL
                    = {id-it 6}, CertificateList
                                                     }
 -- where \{id-it\} = \{id-pkix 4\} = \{1 3 6 1 5 5 7 4\}
 -- This construct MAY also be used to define new PKIX Certificate
 -- Management Protocol request and response messages, or general-
 -- purpose (e.g., announcement) messages for future needs or for
 -- specific environments.
 GenMsqContent ::= SEQUENCE OF InfoTypeAndValue
 -- May be sent by EE, RA, or CA (depending on message content).
 -- The OPTIONAL infoValue parameter of InfoTypeAndValue will typically
 -- be omitted for some of the examples given above. The receiver is
```

- -- free to ignore any contained OBJ. IDs that it does not recognize.
- -- If sent from EE to CA, the empty set indicates that the CA may send
- -- any/all information that it wishes.

GenRepContent ::= SEQUENCE OF InfoTypeAndValue

-- The receiver is free to ignore any contained OBJ. IDs that it does -- not recognize.

ErrorMsgContent ::= SEQUENCE { pKIStatusInfo PKIStatusInfo, OPTIONAL, errorCode INTEGER -- implementation-specific error codes PKIFreeText errorDetails OPTIONAL -- implementation-specific error details }

## Appendix D. Certificate Request Message Format ASN.1 Module

The following section contains the complete ASN.1 module from RFC 2511, the *Certificate Request Message Format*. Only a small subset of the structures defined in [RFC2511] are required to implement this specification. The entire module is provided for completeness. Information about structures defined by this ASN.1 module but not used in the MISPC may be found in [RFC2511].

## CRMF DEFINITIONS IMPLICIT TAGS ::= BEGIN

## **IMPORTS** -- Directory Authentication Framework (X.509) Version, AlgorithmIdentifier, Name, Time, SubjectPublicKeyInfo, Extensions, UniqueIdentifier FROM AuthenticationFramework { joint-iso-itu-t ds(5) module(1) authenticationFramework(7) 3 } -- Certificate Extensions (X.509) GeneralName FROM CertificateExtensions {joint-iso-ccitt ds(5) module(1) certificateExtensions(26) 0} -- Cryptographic Message Syntax **EnvelopedData** FROM CryptographicMessageSyntax { iso(1) member-body(2) us(840) rsadsi(113549) pkcs(1) pkcs-9(9) smime(16) modules(0) cms(1) }; CertRegMessages ::= SEQUENCE SIZE (1..MAX) OF CertRegMsg CertRegMsg ::= SEQUENCE { CertReq CertRequest, ProofOfPossession OPTIONAL, Pop -- content depends upon key type realnfo SEQUENCE SIZE(1..MAX) OF AttributeTypeAndValue OPTIONAL } CertRequest ::= SEQUENCE { CertRegId INTEGER, -- ID for matching request and reply CertTemplate, -- Selected fields of cert to be issued CertTemplate Controls OPTIONAL } -- Attributes affecting issuance Controls CertTemplate ::= SEQUENCE { Version [0] Version **OPTIONAL.** serialNumber [1] INTEGER OPTIONAL, signingAlg [2] AlgorithmIdentifier OPTIONAL, [3] Name issuer OPTIONAL, validitv [4] OptionalValidity OPTIONAL, subject [5] Name OPTIONAL, [6] SubjectPublicKeyInfo OPTIONAL, publicKey issuerUID [7] Uniqueldentifier **OPTIONAL**,

```
subjectUID [8] UniqueIdentifier
                                    OPTIONAL,
  extensions [9] Extensions
                                   OPTIONAL }
OptionalValidity ::= SEQUENCE {
  notBefore [0] Time OPTIONAL,
  notAfter [1] Time OPTIONAL } --at least one MUST be present
Controls ::= SEQUENCE SIZE(1..MAX) OF AttributeTypeAndValue
AttributeTypeAndValue ::= SEQUENCE {
           OBJECT IDENTIFIER,
  type
  value
           ANY DEFINED BY type }
ProofOfPossession ::= CHOICE {
  raVerified
                     [0] NULL.
  -- used if the RA has already verified that the requester is in
  -- possession of the private key
  signature
                     [1] POPOSigningKey,
  keyEncipherment
                     [2] POPOPrivKey,
  keyAgreement
                     [3] POPOPrivKey }
POPOSigningKey ::= SEQUENCE {
                     [0] POPOSKInput OPTIONAL,
  poposkInput
  algorithmIdentifier AlgorithmIdentifier,
  signature
                     BIT STRING }
-- The signature (using "algorithmIdentifier") is on the
-- DER-encoded value of popInput. NOTE: If poposkInput is present
-- in the pop field, popInput is constructed
-- with otherinput. If poposkinput is not present, subject is the name
-- from CertTemplate. Note that the encoding of PopInput is
 intentionally ambiguous.
PoposkInput ::= CHOICE {
             Subject name,
             Sender [0] generalName,
             publicKeyMAC [1] PKMACValue
                          }
-- The pop is calculated upon the structure popInput, which is defined
-- as follows:
 PopInput ::= SEQUENCE {
             CHOICE {
             otherinput popsklnput,
             subiect
                       name }.
             publicKey subjectpublicKey
             }
-- If poposkInput is present
-- in the pop field, popInput is constructed
-- with otherinput. If poposkinput is not present, subject is the name
-- from CertTemplate. Note that the encoding of PopInput is
```

-- intentionally ambiguous.

```
PKMACValue ::= SEQUENCE {
 algld AlgorithmIdentifier,
 -- algorithm value shall be PasswordBasedMac {1 2 840 113533 7 66 13}
 -- parameter value is PBMParameter
 value BIT STRING }
PBMParameter ::= SEQUENCE {
   salt
                      OCTET STRING,
   owf
                      Algorithmldentifier,
   -- AlgId for a One-Way Function (SHA-1 recommended)
   iterationCount
                      INTEGER,
   -- number of times the OWF is applied
                      AlgorithmIdentifier
   mac
   -- the MAC AlgId (e.g., DES-MAC, Triple-DES-MAC [PKCS11],
} -- or HMAC [RFC2104, RFC2202])
POPOPrivKey ::= CHOICE {
  thisMessage
                             [0] BIT STRING,
  -- posession is proven in this message (which contains the private
  -- kev itself (encrypted for the CA))
                             [1] SubsequentMessage,
  subsequentMessage
  -- possession will be proven in a subsequent message
                             [2] BIT STRING }
  dhMAC
  -- for keyAgreement (only), possession is proven in this message
  -- (which contains a MAC (over the DER-encoded value of the
  -- certReg parameter in CertRegMsg, which MUST include both subject
  -- and publicKey) based on a key derived from the end entity's
  -- private DH key and the CA's public DH key);
  -- the dhMAC value MUST be calculated as per the directions given
  -- in Appendix A.
SubsequentMessage ::= INTEGER {
  encrCert (0),
  -- requests that resulting certificate be encrypted for the
  -- end entity (following which, POP will be proven in a
  -- confirmation message)
  challengeResp (1) }
  -- requests that CA engage in challenge-response exchange with
  -- end entity in order to prove private key possession
-- Object identifier assignments --
id-pkix OBJECT IDENTIFIER ::= { iso(1) identified-organization(3)
dod(6) internet(1) security(5) mechanisms(5) 7 }
-- arc for Internet X.509 PKI protocols and their components
id-pkip OBJECT IDENTIFIER ::= { id-pkix 5 }
-- Registration Controls in CRMF
id-regCtrl OBJECT IDENTIFIER ::= { id-pkip 1 }
-- The following definition may be uncommented for use with
-- ASN.1 compilers which do not understand UTF8String.
```

```
-- UTF8String ::= [UNIVERSAL 12] IMPLICIT OCTET STRING
id-regCtrl-regToken OBJECT IDENTIFIER ::= { id-regCtrl 1 }
--with syntax:
RegToken ::= UTF8String
id-regCtrl-authenticator OBJECT IDENTIFIER ::= { id-regCtrl 2 }
--with syntax:
Authenticator ::= UTF8String
id-regCtrl-pkiPublicationInfo OBJECT IDENTIFIER ::= { id-regCtrl 3 }
--with syntax:
PKIPublicationInfo ::= SEQUENCE {
 action
         INTEGER {
        dontPublish (0),
        pleasePublish (1) },
 publnfos SEQUENCE SIZE (1..MAX) OF SinglePublnfo OPTIONAL }
  -- publnfos MUST NOT be present if action is "dontPublish"
  -- (if action is "pleasePublish" and publnfos is omitted.
  -- "dontCare" is assumed)
SinglePubInfo ::= SEQUENCE {
  pubMethod INTEGER {
    dontCare (0),
    x500
             (1),
    web
             (2),
    Idap
             (3) }.
  pubLocation GeneralName OPTIONAL }
id-regCtrl-pkiArchiveOptions OBJECT IDENTIFIER ::= { id-regCtrl 4 }
--with syntax:
PKIArchiveOptions ::= CHOICE {
  encryptedPrivKey [0] EncryptedKey,
  -- the actual value of the private key
  keyGenParameters [1] KeyGenParameters,
  -- parameters which allow the private key to be re-generated
  archiveRemGenPrivKey [2] BOOLEAN }
  -- set to TRUE if sender wishes receiver to archive the private
  -- key of a key pair which the receiver generates in response to
  -- this request; set to FALSE if no archival is desired.
EncryptedKey ::= CHOICE {
  encryptedValue
                     EncryptedValue,
  envelopedData [0] EnvelopedData }
  -- The encrypted private key MUST be placed in the envelopedData
  -- encryptedContentInfo encryptedContent OCTET STRING.
EncryptedValue ::= SEQUENCE {
  intendedAlg [0] AlgorithmIdentifier OPTIONAL,
  -- the intended algorithm for which the value will be used
               [1] AlgorithmIdentifier OPTIONAL,
  symmAlg
```

```
-- the symmetric algorithm used to encrypt the value
```

encSymmKey [2] BIT STRING **OPTIONAL**, -- the (encrypted) symmetric key used to encrypt the value [3] AlgorithmIdentifier OPTIONAL, keyAlq -- algorithm used to encrypt the symmetric key valueHint [4] OCTET STRING OPTIONAL, -- a brief description or identifier of the encValue content -- (may be meaningful only to the sending entity, and used only -- if EncryptedValue might be re-examined by the sending entity -- in the future) BIT STRING } encValue -- the encrypted value itself KeyGenParameters ::= OCTET STRING id-regCtrl-oldCertId **OBJECT IDENTIFIER ::= { id-regCtrl 5 }** --with syntax: OldCertId ::= CertId CertId ::= SEQUENCE { issuer GeneralName. serialNumber INTEGER } id-regCtrl-protocolEncrKey OBJECT IDENTIFIER ::= { id-regCtrl 6 } --with syntax: ProtocolEncrKey ::= SubjectPublicKeyInfo -- Registration Info in CRMF id-regInfo OBJECT IDENTIFIER ::= { id-pkip 2 } id-regInfo-utf8Pairs OBJECT IDENTIFIER ::= { id-regInfo 1 } --with syntax UTF8Pairs ::= UTF8String id-regInfo-certReg **OBJECT IDENTIFIER ::= { id-reginfo 2 }** --with syntax CertReq ::= CertRequest

```
END
```

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